

# **LogiCORE™ 1-Gigabit Ethernet MAC v8.3**

## **User Guide**

UG144 August 8, 2007





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## Revision History

The following table shows the revision history for this document.

Date	Version	Revision
09/30/04	1.0	Initial Xilinx release.
04/28/05	2.0	Updated to 1-Gigabit Ethernet MAC version 6.0, Xilinx tools v7.1i SP1.
01/18/06	3.0	Updated to 1-Gigabit Ethernet MAC version 7.0, Xilinx tools v8.1i.
07/13/06	4.0	Updated to 1-Gigabit Ethernet MAC version 8.0, Xilinx tools v8.2i.
09/21/06	4.1	Updated to 1-Gigabit Ethernet MAC version 8.1, added support for Spartan-3A.
02/15/07	4.2	Updated to 1-Gigabit Ethernet MAC version 8.2, Xilinx tools v9.1i.
08/08/07	5.0	Advanced core version to 8.3, updated various tool versions and trademarks for the IP1 I Jade Minor release.

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# About This Guide

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The LogiCORE™ 1-Gigabit Ethernet MAC v8.3 User Guide provides information about generating the core, customizing and simulating the core utilizing the provided example design, and running the design files through implementation using the Xilinx tools.

## Guide Contents

This guide contains the following chapters:

- [Preface, “About this Guide”](#) introduces the organization and purpose of the guide, a list of additional resources, and the conventions used in this document.
- [Chapter 1, “Introduction”](#) describes the core and related information, including recommended design experience, additional resources, technical support, and submitting feedback to Xilinx.
- [Chapter 2, “Core Architecture”](#) provides an overview of the core and discusses the Physical/Client signal interfaces.
- [Chapter 3, “Generating the Core”](#) describes the graphical user interface options used to generate the core.
- [Chapter 4, “Designing with the Core”](#) through [Chapter 8, “Configuration and Status”](#) describe design parameters, including how to initialize the core, generate and consume core packets, and how to operate the Management Interface.
- [Chapter 9, “Constraining the Core”](#) describes the constraints associated with the core.
- [Chapter 10, “Clocking and Resetting”](#) discusses special design considerations associated with clock management logic, including the Gigabit Media Independent Interface (GMII) and Reduced Gigabit Media Independent Interface (RGMII) options.
- [Chapter 11, “Interfacing to Other Cores”](#) describes how to interface the 1-Gigabit Ethernet MAC core to the Ethernet 1000BASE-X PCS/PMA or SGMII core and the Ethernet Statistics core.
- [Chapter 12, “Implementing Your Design”](#) provides instructions for how to set up synthesis, simulation, and implementation environments and how to generate a bitstream through the design flow.
- [Appendix A, “Using the Client-Side FIFO”](#) describes the FIFO provided in the example design that accompanies the GEMAC Core.
- [Appendix B, “Core Verification, Compliance, and Interoperability”](#) describes how the core was verified and certified for compliance.
- [Appendix C, “Calculating DCM Phase-Shifting”](#) provides information about how to calculate the system timing requirements when using DCMs with the core.
- [Appendix D, “Core Latency”](#) describes the latency of the core.

## Additional Resources

For additional information, go to [www.xilinx.com/support](http://www.xilinx.com/support). The following table lists some of the resources you can access from this website or by using the provided URLs.

Resource	Description/URL
Tutorials	Tutorials covering Xilinx design flows, from design entry to verification and debugging <a href="http://www.xilinx.com/support/techsup/tutorials/index.htm">www.xilinx.com/support/techsup/tutorials/index.htm</a>
Answer Browser	Database of Xilinx solution records <a href="http://www.xilinx.com/xlnx/xil_ans_browser.jsp">www.xilinx.com/xlnx/xil_ans_browser.jsp</a>
Data Sheets	Device-specific information on Xilinx device characteristics, including readback, boundary scan, configuration, length count, and debugging <a href="http://www.xilinx.com/xlnx/xweb/xil_publications_index.jsp">www.xilinx.com/xlnx/xweb/xil_publications_index.jsp</a>
Problem Solvers	Interactive tools that allow you to troubleshoot your design issues <a href="http://www.xilinx.com/support/troubleshoot/psolvers.htm">www.xilinx.com/support/troubleshoot/psolvers.htm</a>
Tech Tips	Latest news, design tips, and patch information for the Xilinx design environment <a href="http://www.xilinx.com/xlnx/xil_tt_home.jsp">www.xilinx.com/xlnx/xil_tt_home.jsp</a>

## Conventions

This document uses the following conventions. An example illustrates each convention.

### Typographical

The following typographical conventions are used in this document:

Convention	Meaning or Use	Example
Courier font	Messages, prompts, and program files that the system displays	speed grade: - 100
<b>Courier bold</b>	Literal commands you enter in a syntactical statement	<b>ngdbuild</b> design_name
<i>Italic font</i>	Variables in a syntax statement for which you must supply values	See the <i>Development System Reference Guide</i> for more information.
	References to other manuals	See the <i>User Guide</i> for details.
	Emphasis in text	If a wire is drawn so that it overlaps the pin of a symbol, the two nets are <i>not</i> connected.
Dark Shading	Items that are not supported or reserved	This feature is not supported

Convention	Meaning or Use	Example
Square brackets [ ]	An optional entry or parameter. However, in bus specifications, such as <b>bus[7:0]</b> , they are required.	<b>ngdbuild</b> [option_name] design_name
Braces { }	A list of items from which you must choose one or more	<b>lowpwr</b> = {on off}
Vertical bar	Separates items in a list of choices	<b>lowpwr</b> = {on off}
Vertical ellipsis . . .	Repetitive material that has been omitted	IOB #1: Name = QOUT' IOB #2: Name = CLKIN' . . .
Horizontal ellipsis ...	Omitted repetitive material	<b>allow block</b> block_name loc1 loc2 ... locn;
Notations	The prefix '0x' or the suffix 'h' indicate hexadecimal notation	A read of address 0x00112975 returned 45524943h.
	An '_n' means the signal is active low	usr_teof_n is active low.

## Online Document

The following linking conventions are used in this document:

Convention	Meaning or Use	Example
Blue text	Cross-reference link to a location in the current document	See the section " <a href="#">Additional Resources</a> " for details. See " <a href="#">Title Formats</a> " in <a href="#">Chapter 1</a> for details.
<a href="#">Blue, underlined text</a>	Hyperlink to a website (URL)	Go to <a href="http://www.xilinx.com">www.xilinx.com</a> for the latest speed files.





## Introduction

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The 1-Gigabit Ethernet MAC (GEMAC) core is a fully verified solution that supports Verilog-HDL and VHDL. In addition, the example design provided with the core is provided in both Verilog and VHDL.

This chapter introduces the GEMAC core and provides other related information, including recommended design experience, additional resources, technical support, and ways to submit feedback to Xilinx.

### About the Core

The GEMAC core is a Xilinx CORE Generator™ IP core, included in the latest IP Update on the Xilinx IP Center. For detailed information about the core, see [www.xilinx.com/systemio/gmac/index.htm](http://www.xilinx.com/systemio/gmac/index.htm). For information about licensing options, see Chapter 2, “Licensing the Core,” in the *1-Gigabit Ethernet MAC Getting Started Guide*.

### Recommended Design Experience

Although the GEMAC core is a fully verified solution, the challenge associated with implementing a complete design varies, depending on the configuration and functionality of the application. For best results, previous experience building high performance, pipelined FPGA designs using Xilinx implementation software and user constraint files (UCFs) is recommended.

Contact your local Xilinx representative for a closer review and estimation for your specific requirements.

### Additional Core Resources

For detailed information and updates about the GEMAC core, see the following documents, located on the GEMAC product page [www.xilinx.com/systemio/gmac/index.htm](http://www.xilinx.com/systemio/gmac/index.htm)

- *1-Gigabit Ethernet MAC Data Sheet*
- *1-Gigabit Ethernet MAC Release Notes*
- *1-Gigabit Ethernet MAC Getting Started Guide*

For updates to this document, see the *1-Gigabit Ethernet MAC User Guide*, also located on the GEMAC product page.

## Related Xilinx Ethernet Products and Services

See the Ethernet Products and Services page at:

[www.xilinx.com/products/design\\_resources/conn\\_central/grouping/ethernet.htm](http://www.xilinx.com/products/design_resources/conn_central/grouping/ethernet.htm)

## Specifications

- IEEE 802.3 2002
- Reduced Gigabit Media Independent Interface (RGMII) version 2.0

## Technical Support

For technical support, see [support.xilinx.com/](http://support.xilinx.com/). Questions are routed to a team of engineers with expertise using the GEMAC core.

Xilinx will provide technical support for use of this product as described in the *1-Gigabit Ethernet MAC User Guide* and the *1-Gigabit Ethernet MAC Getting Started Guide*. Xilinx cannot guarantee timing, functionality, or support of this product for designs that do not follow these guidelines.

## Feedback

Xilinx welcomes comments and suggestions about the GEMAC core and the documentation supplied with the core.

### GEMAC Core

For comments or suggestions about the GEMAC core, please submit a WebCase from [support.xilinx.com/](http://support.xilinx.com/). Be sure to include the following information:

- Product name
- Core version number
- Explanation of your comments

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- Document title
- Document number
- Page number(s) to which your comments refer
- Explanation of your comments

## Core Architecture

This chapter describes the GEMAC core architecture, including the major functional blocks and all interfaces.

### System Overview

Figure 2-1 illustrates a block diagram of the GEMAC core with all the major functional blocks and interfaces. Descriptions of the functional blocks and interfaces are provided in the sections that follow.

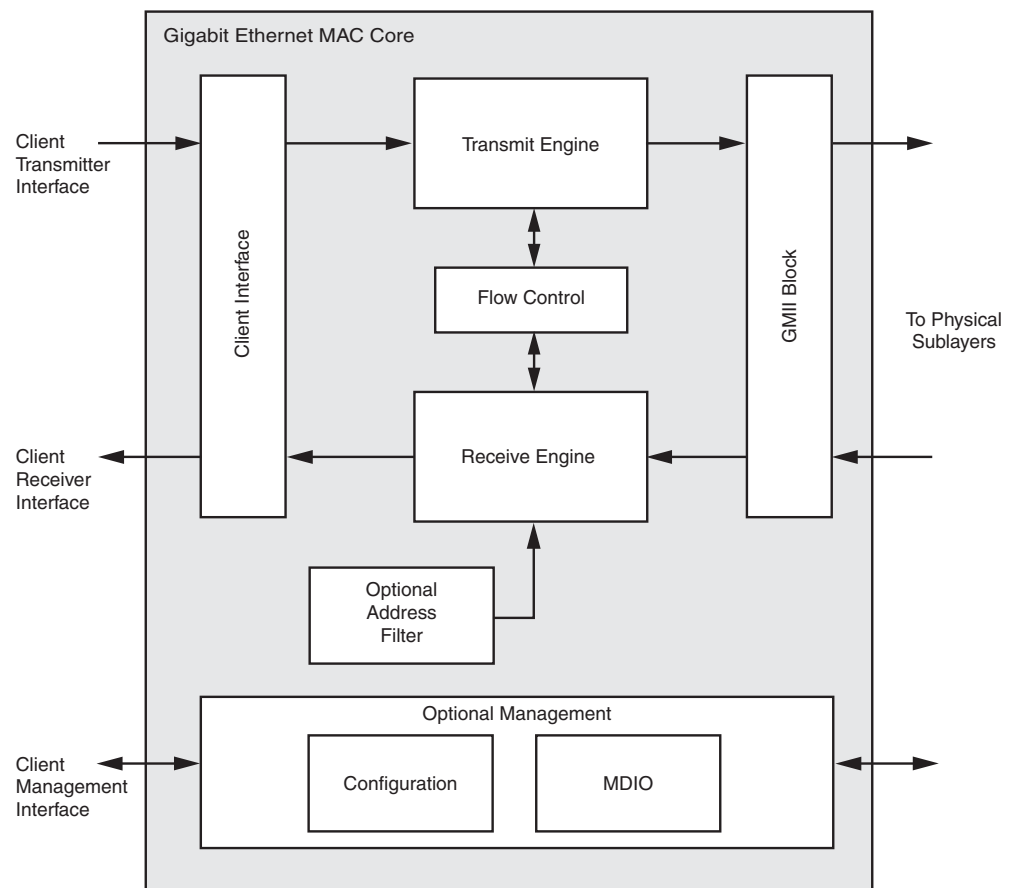


Figure 2-1: Block Diagram

## Core Components

### Transmit Engine

The Transmit Engine accepts Ethernet frame data from the Client Transmitter Interface, adds the preamble field to the start of the frame, adds padding bytes (if required) to ensure that the frame meets the minimum frame length requirements, and adds the frame check sequence (when configured to do so). The transmitter also ensures that the inter-frame spacing between successive frames is at least the minimum specified. The frame is then converted into a format that is compatible with the GMII and sent to the GMII Block.

### Receive Engine

The Receive Engine accepts Ethernet frame data from the GMII Block, removes the preamble field at the start of the frame, removes padding bytes and Frame Check Sequence (if required, and when configured to do so). The receiver also performs error detection on the received frame using information such as the frame check sequence field, received GMII error codes, and legal frame size boundaries.

### Flow Control

The Flow Control block is designed to clause 31 of the *IEEE 802.3-2002* standard. The MAC may be configured to send pause frames and to act upon their reception. These two behaviors can be configured independently.

### Address Filter

The Address Filter checks the address of incoming frames into the receiver. If the Address Filter is enabled, the device will not pass frames that do not contain one of a set of known addresses to the client.

### Management Interface

The optional processor-independent Management Interface has standard address, data, and control signals. It may be used as is, or you can apply a logical shim to interface to common bus architectures. For more information, [Chapter 8, “Configuration and Status.”](#)

This interface is used to access the following blocks.

- **Configuration Register** After power up or reset, the client may reconfigure the core parameters from their defaults. Configuration changes can be written at any time.
- **MDIO Interface** The Management Interface is also used to access the MDIO interface of the GEMAC core; this interface is typically connected to the MDIO port of a physical layer device (PHY) to access its configuration and status registers. The MDIO format is defined in *IEEE802.3* clause 22.

### GMII Block

This implements GMII style signaling for the physical interface of the core and is typically attached to a physical layer device (PHY), either off-chip or internally integrated.

## Core Interfaces

### GMAC Core with Optional Management Interface

Figure 2-2 shows the pinout for the GEMAC core using the optional Management Interface. The interface is unchanged, regardless of whether the optional Address Filter is included.

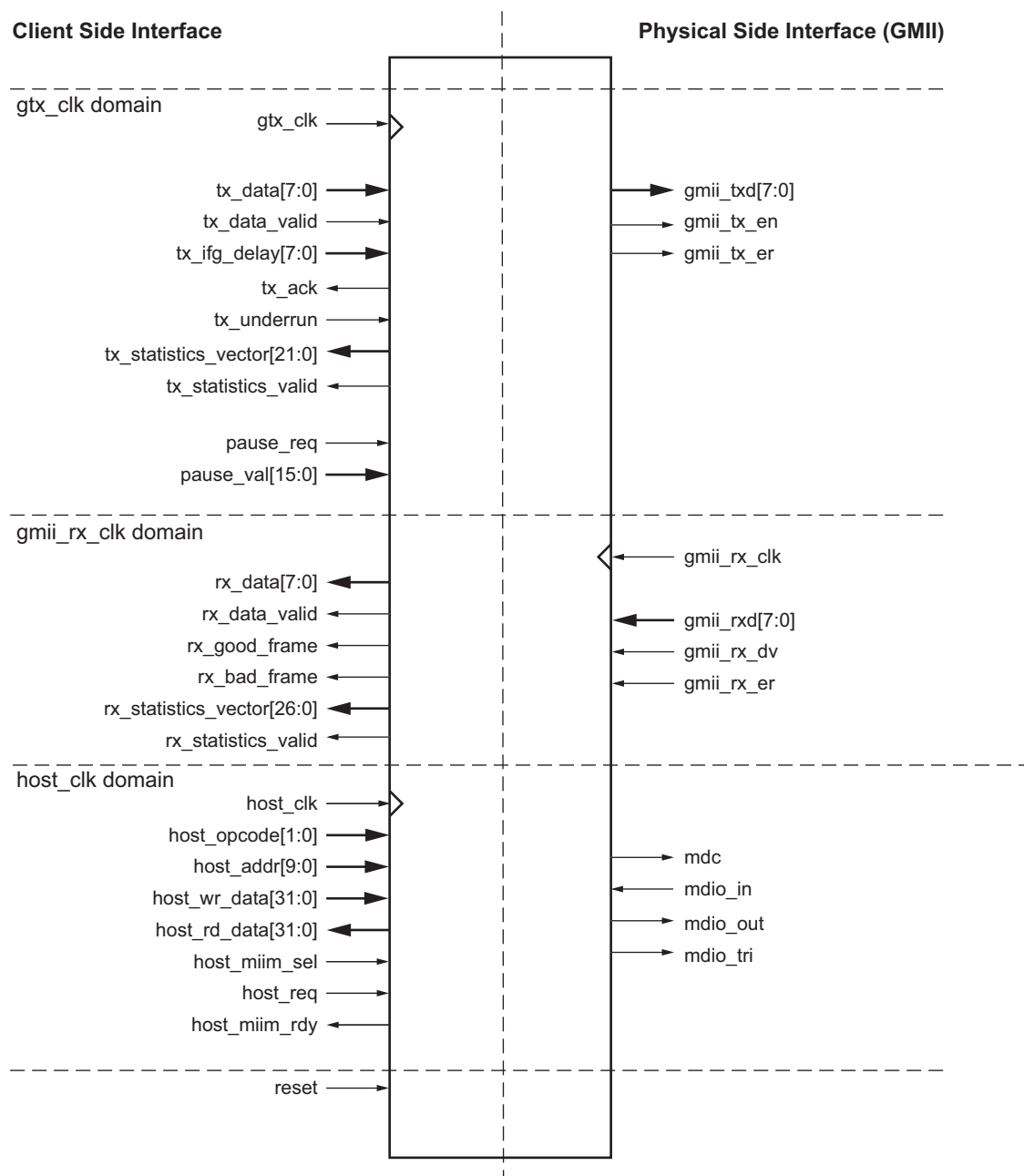


Figure 2-2: Component Pinout for MAC with Optional Management Interface

## GMAC Core Without Management Interface and With Address Filter

Figure 2-3 shows the pinout for the GEMAC core when the optional Management Interface is omitted and the optional Address Filter is included in the core.

The `configuration_vector[64:0]` input provides the method for configuration of the core, and `mac_unicast_address[47:0]` input provides the method of setting the unicast address used by the Address Filter.

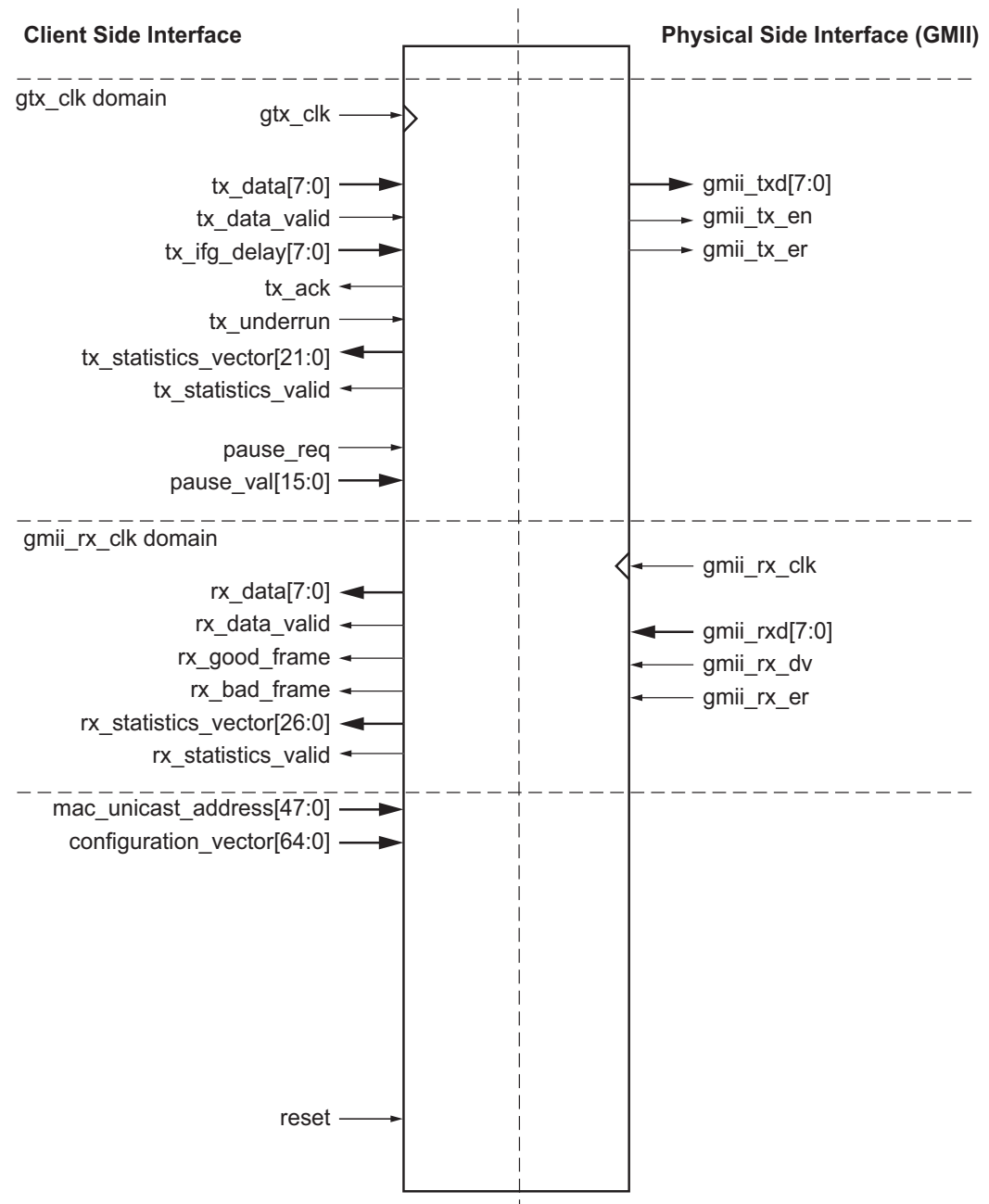


Figure 2-3: Component Pinout for MAC without Optional Management Interface and with Optional Address Filter

## GEMAC Core Without Management Interface and Without Address Filter

Figure 2-4 shows the pinout for the GEMAC core when the optional Management Interface is omitted and the optional Address Filter is omitted.

The `configuration_vector[64:0]` input provides the method for configuration of the core.

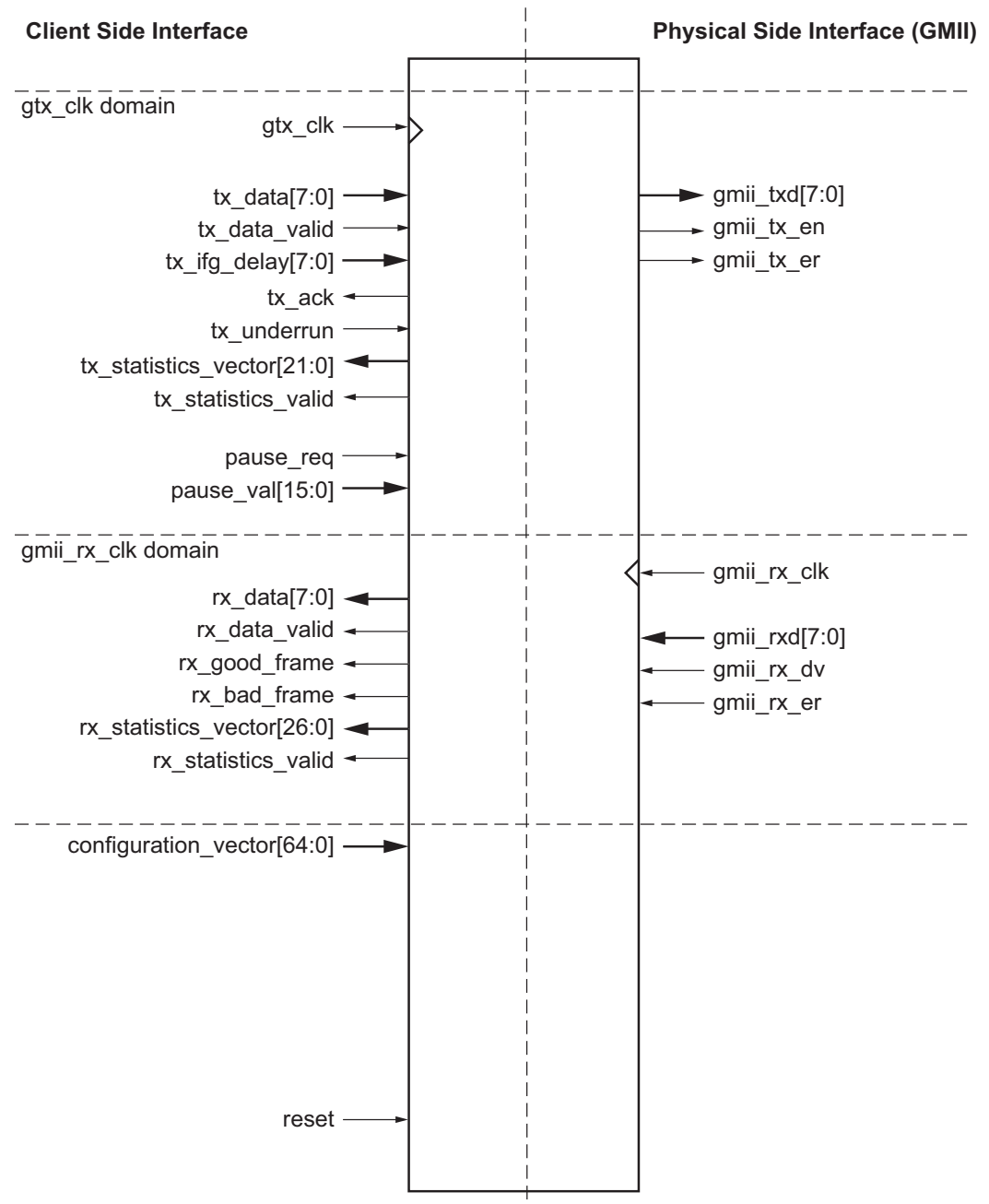


Figure 2-4: Component Pinout for MAC without Optional Management Interface or Optional Address Filter

All ports of the core are internal connections in FPGA fabric. An HDL example design is delivered with the core that will add IBUFs, OBUFs, and IOB flip-flops to the external signals of the Gigabit Media Independent Interface (GMII) or Reduced Gigabit Media Independent Interface (RGMII).

All clock management logic is placed in this example design, which allows for more flexibility in implementation (for example, in designs using multiple cores). This example design is provided in both VHDL and Verilog. For more information about example designs, see the *1-Gigabit Ethernet MAC Getting Started Guide*.

## Client Side Interface

### Transmitter Interface

[Table 2-1](#) describes the client-side transmitter signals of the GEMAC core. These signals are used to transmit data from the client logic into the core. For more information, see [“Transmitting Outbound Frames,” on page 43](#).

The Transmitter Interface is designed to be connected to internal device logic only. Attempting to add external ports to this interface will result in a breakdown of the handshaking protocol used by this interface.

**Table 2-1: Transmitter Client Interface Signal Pins**

Signal	Direction	Clock Domain	Description
gtx_clk	Input	n/a	Clock signal provided to the core at 125 MHz. Tolerance must be within <i>IEEE 802.3-2002</i> specification. This clock signal is used by all of the transmitter logic.
tx_data[7:0]	Input	gtx_clk	Frame data to be transmitted is supplied on this port.
tx_data_valid	Input	gtx_clk	Control signal for tx_data port.
tx_ifg_delay[7:0]	Input	gtx_clk	Control signal for configurable Inter Frame Gap adjustment.
tx_ack	Output	gtx_clk	Handshaking signal asserted when the current data on tx_data has been accepted.
tx_underrun	Input	gtx_clk	Asserted by client to force GEMAC core to corrupt the current frame.
tx_statistics_vector[21:0]	Output	gtx_clk	Provides statistical information about the last frame transmitted.
tx_statistics_valid	Output	gtx_clk	Asserted at end of frame transmission, indicating that the tx_statistics_vector is valid.



## Receiver Interface

[Table 2-2](#) describes the client-side receiver signals of the GEMAC core. These signals are used by to transfer data to the client. For more information, see [“Receiving Inbound Frames,” on page 37](#).

**Table 2-2: Receive Client Interface Signal Pins**

Signal	Direction	Clock Domain	Description
rx_data[7:0]	Output	gmii_rx_clk	Frame data received is supplied on this port.
rx_data_valid	Output	gmii_rx_clk	Control signal for the rx_data port.
rx_good_frame	Output	gmii_rx_clk	Asserted at end of frame reception to indicate that the frame should be processed by the MAC client.
rx_bad_frame	Output	gmii_rx_clk	Asserted at end of frame reception to indicate that the frame should be discarded by the MAC client.
rx_statistics_vector[26:0]	Output	gmii_rx_clk	Provides statistical information about the last frame received.
rx_statistics_valid	Output	gmii_rx_clk	Asserted at end of frame reception, indicating that the rx_statistics_vector is valid.

## Flow Control Interface

[Table 2-3](#) describes the signals used by the client to request a flow control action from the transmit engine. For more information, see [“Using Flow Control,” on page 49](#).

**Table 2-3: Flow Control Interface Signal Pinout**

Signal	Direction	Clock Domain	Description
pause_req	Input	gtx_clk	Pause request. sends a pause frame down the link.
pause_val[15:0]	Input	gtx_clk	Pause value; inserted into the parameter field of the transmitted pause frame.

## Management Interface (Optional)

Table 2-4 describes the optional signals used by the client to access the management features of the GEMAC core. For more information, see “Using the Optional Management Interface,” on page 71.

Table 2-4: Optional Management Interface Signal Pinout

Signal	Direction	Clock Domain	Description
host_clk	Input	n/a	Clock for the Management Interface; must be 10 MHz or above.
host_opcode[1:0]	Input	host_clk	Defines operation to be performed over MDIO interface. Bit 1 is also used as a read/write control signal for configuration register access.
host_addr[9:0]	Input	host_clk	Address of register to be accessed.
host_wr_data[31:0]	Input	host_clk	Data to write to register .
host_rd_data[31:0]	Output	host_clk	Data read from register.
host_miim_sel	Input	host_clk	When asserted, the MDIO interface is accessed. When not asserted, the configuration registers are accessed.
host_req	Input	host_clk	Used to signal a transaction on the MDIO interface.
host_miim_rdy	Output	host_clk	When high, the MDIO interface has completed any pending transaction and is ready for a new transaction.

## MAC Unicast Address (Optional)

Table 2-5 describes the alternative method of access to the unicast address registers when the optional Management Interface is not present.

Table 2-5: Optional MAC Unicast Address Signal Pinout

Signal	Direction	Description
mac_unicast_address[47:0]	Input	Used to assess the MAC unicast address registers when the Management Interface is not used

**Note:** All bits are registered on input but may be treated as asynchronous inputs.

## Configuration Vector (Optional)

Table 2-6 describes the alternative to the optional Management Interface signals. The Configuration Vector uses direct inputs to the core to replace the functionality of the MAC

configuration bits. For more information, see [“Access without the Management Interface,”](#) on page 83.

**Table 2-6: Optional Configuration Vector Signal Pinout**

Signal	Direction	Description
configuration_vector[64:0]	Input	Used to replace the functionality of the MAC Configuration Registers when the Management Interface is not used

**Note:** All bits are registered on input but may be treated as asynchronous inputs.

## Asynchronous Reset

[Table 2-7](#) describes the asynchronous reset signal for the entire core.

**Table 2-7: Reset Signal**

Signal	Direction	Clock Domain	Description
reset	Input	n/a	Asynchronous reset for entire core

## Physical Side Interface

### GMII

[Table 2-8](#) describes the GMII-style interface signals of the core. For more information, see [Chapter 7, “Using the Physical Side Interface.”](#)

**Table 2-8: GMII Interface Signal Pinout**

Signal	Direction	Clock Domain	Description
gmii_txd[7:0]	Output	gtx_clk	Transmit data from MAC
gmii_tx_en	Output	gtx_clk	Transmit control signal from MAC
gmii_tx_er	Output	gtx_clk	Transmit control signal from MAC
gmii_rx_clk	Input	n/a	Receive clock from external PHY (125 MHz)
gmii_rxd[7:0]	Input	gmii_rx_clk	Received data to MAC
gmii_rx_dv	Input	gmii_rx_clk	Received control signal to MAC
gmii_rx_er	Input	gmii_rx_clk	Received control signal to MAC

## MDIO Interface

Table 2-9 describes the MDIO Interface signals. For more information, see [“Using the MDIO interface,” on page 68](#).

**Table 2-9: MDIO Interface Signal Pinout**

Signal	Direction	Clock Domain	Description
mdc	Output	host_clk	Management Clock: programmable frequency derived from host_clk.
mdio_in <sup>1</sup>	Input	host_clk	Input data signal for communication with PHY configuration and status. Tie high if unused.
mdio_out <sup>1</sup>	Output	host_clk	Output data signal for communication with PHY configuration and status.
mdio_tri <sup>1</sup>	Output	host_clk	Tristate control for MDIO signals; 0 signals that the value on mdio_out should be asserted onto the MDIO bus.

1. mdio\_in, mdio\_out and mdio\_tri can be connected to a Tri-state buffer to create a bi-directional mdio signal suitable for connection to an external PHY.

## Generating the Core

The GEMAC core is generated through the Xilinx CORE Generator using a graphical user interface (GUI). This chapter describes the GUI options used to generate and customize the core.

### Graphical User Interface

Figure 3-1 shows the main GEMAC core user GUI screen.

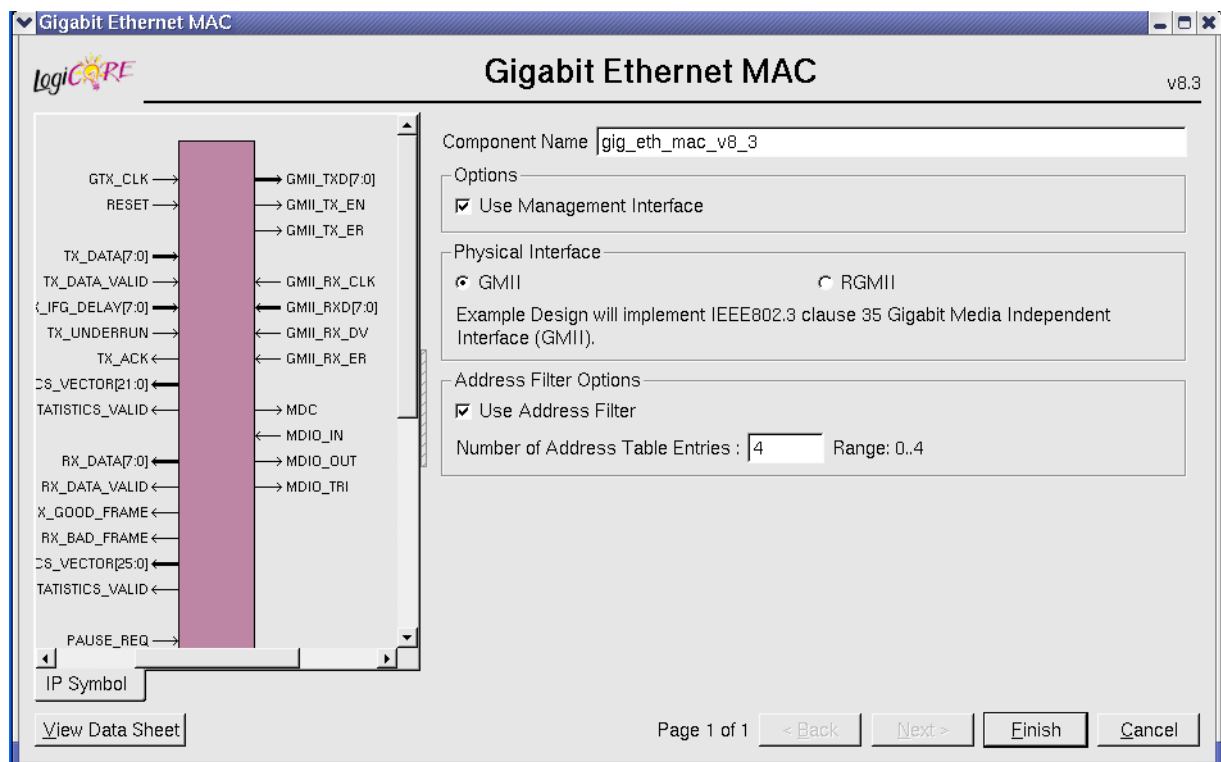


Figure 3-1: 1-Gigabit Ethernet MAC Main Screen

For general help starting and using CORE Generator on your system, see the documentation supplied with Xilinx ISE, including the *CORE Generator Guide* at [www.xilinx.com/support/software\\_manuals.htm](http://www.xilinx.com/support/software_manuals.htm).

## Component Name

The component name is used as the base name of the output files generated for the core. Names must begin with a letter and must be composed from the following characters: a through z, 0 through 9 and “\_”.

## Management Interface

Select this option to include the optional Management Interface (see [“Using the Optional Management Interface,” on page 71](#)). If this option is not selected, the core is generated with a replacement configuration vector (see [“Access without the Management Interface,” on page 83](#)). The default is to use the Management Interface.

## Address Filter

Select this option to include the optional Address Filter. This prevents the reception of frames that are not addressed to this MAC (see [“Address Filter,” on page 41](#)). The default is to use the Address Filter.

## Number of Address Table Entries

The Address Filter can be instantiated with an address table that holds up to 4 additional valid addresses. You may select an integer between 0 and 4 to define the number of addresses that are present in the table.

This option is only available when the Management Interface and Address Filter have been selected. The default is to use 4 address table entries.

## Physical Interface

Depending on the target Xilinx FPGA architecture, it may be possible to select from two different physical interface choices for the core:

- **GMII** See [Chapter 7, “Implementing External GMII”](#)
- **RGMII** See [Chapter 7, “Implementing External RGMII”](#)

The choice of physical interface determines the content of the example design delivered with the core. The external GMII or RGMII is added in the HDL top-level design file. There is no change in the core netlist for this option. The default is the GMII physical interface.

## Parameter Values in the XCO File

XCO file parameter names and their values are identical to the names and values shown in the GUI, except that underscore characters ( \_ ) are used instead of spaces. The text in an XCO file is not case-sensitive.

[Table 3-1](#) defines the XCO file parameters and values and summarizes the GUI defaults. The following is an example of the CSET parameters in an XCO file.

```
CSET component_name = abc123
CSET physical_interface = gmii
CSET management_interface = true
CSET address_filter = true
CSET no_of_address_table_entries = 4
```

Table 3-1: XCO File Values and Default Values

Parameter	XCO File Values	Default GUI Setting
component_name	ASCII text starting with a letter and based upon the following character set: a..z, 0..9 and _	gig_eth_mac_v8_2
physical_interface	One of the following keywords: gmii, rgmii	gmii
management_interface	One of the following keywords: true, false	true
address_filter	One of the following keywords: true, false	true
no_of_address_table_entries	Integer in the range 0 - 4	4

## Output Generation

The output files generated from the CORE Generator tool are placed in the CORE Generator project directory. The list of output files includes the following items.

- Netlist file for the core
- Supporting CORE Generator files
- Release notes and documentation
- Subdirectories containing an HDL example design
- Scripts to run the core through the back-end tools and to simulate the core using the Mentor ModelSim® simulator and Cadence IUS

See the *1-Gigabit Ethernet MAC Getting Started Guide* for more information about the CORE Generator output files and for details on the HDL example design.





# Designing with the Core

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This chapter provides general guidelines for creating designs using the GEMAC core. To work with the example design included with the GEMAC core, see the *1-Gigabit Ethernet MAC Getting Started Guide*.

## General Design Guidelines

This section describes the steps required to turn a GEMAC core into a fully functioning design integrated with user-application logic. Not all implementations require all the design steps described in this chapter. The following sections discuss the design steps required for various implementations. For best results, carefully follow the logic design guidelines.

### Design Steps

Generate the core from the Xilinx CORE Generator. See [Chapter 3, “Generating the Core.”](#)

#### Using the Example Design as a Starting Point

The GEMAC core is delivered through the CORE Generator with an HDL example design built around the core, allowing the functionality of the core to be demonstrated using either a simulation package or in hardware, if placed on a suitable board. [Figure 4-1](#) is a block diagram of the example design. For more information about the example design, see the *1-Gigabit Ethernet MAC Getting Started Guide*.

The example design illustrates how to:

- Instantiate the core from HDL.
- Source and use the client-side interface ports of the core from application logic.
- Connect the physical-side interface of the core (GMII or RGMII) to device IOBs creating an external interface. (See [Chapter 7, “Using the Physical Side Interface.”](#))
- Derive the clock management logic, as described in [Chapter 10, “Clocking and Resetting.”](#)

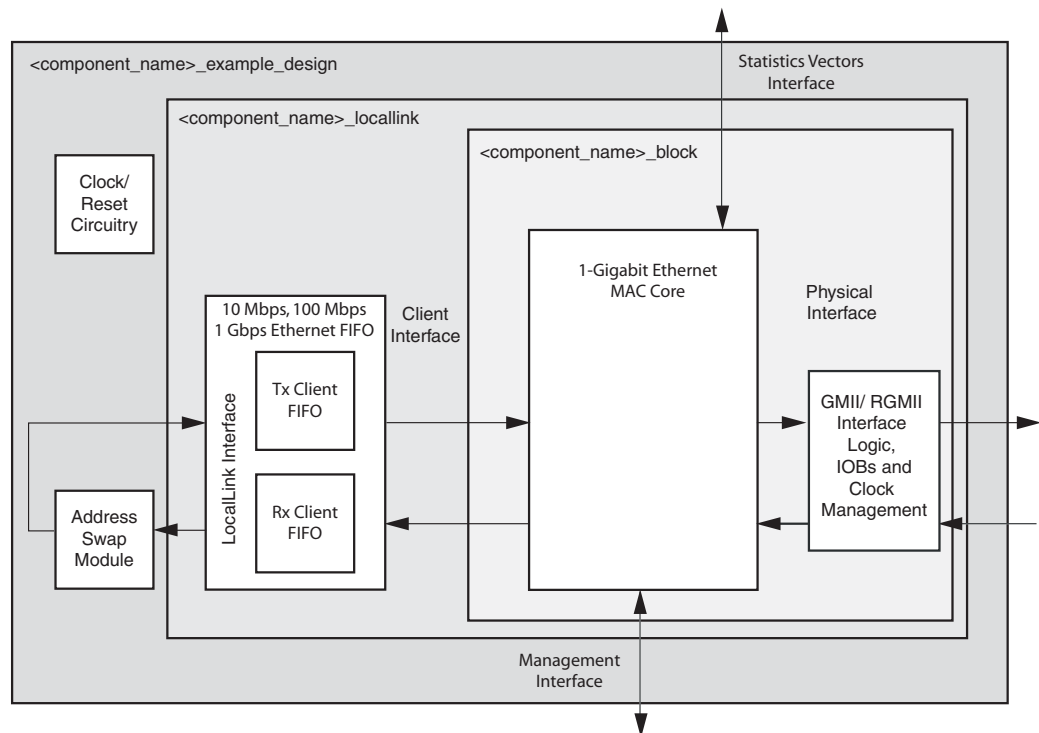


Figure 4-1: 1-Gigabit Ethernet MAC Core Example Design

Using the example design as a starting point, you can do the following:

- Edit the HDL top level of the example design file to:
  - Change the clocking scheme.
  - Add/remove IOBs as required.
  - Replace the client loopback logic with the users specific application logic.
  - Adapt the 10 Mbps, 100 Mbps, 1 Gbps Ethernet FIFO to suit the users specific application (see “Using the Client-Side FIFO”).
- Synthesize the entire design.
 

The Xilinx Synthesis Tool (XST) script and Project file in the `/implement` directory may be adapted to include any HDL files you may want to add.
- Run the `implement` script in the `/implement` directory to create a top-level netlist for the design. The script may also run the Xilinx tools **map**, **par**, and **bitgen**, creating a bitstream that can be downloaded to a Xilinx device. SimPrim-based simulation models for the entire design are also produced by the `implement` scripts.
- Simulate the entire design using the demonstration test bench provided as a template in the `/simulation` directory.
- Download the bitstream to a target device.

## Implementing the 1-Gigabit Ethernet MAC in Your Application

The example design can be studied as an example of how to do the following:

- Instantiate the core from HDL.
- Source and use the client-side interface ports of the core from application logic.

- Connect the physical-side interface of the core (GMII or RGMII) to device IOBs to create an external interface.
- Derive the clock management logic.

After working with the example design, you can write your own HDL application, using single or multiple instances of the GEMAC core. Client-side interfaces and operation of the core are detailed later in this chapter. For more information, see:

- Clock Management Logic in [Chapter 10, “Clocking and Resetting.”](#)
- Using the GEMAC core in conjunction with the Ethernet 1000BASE-X PCS/PMA or SGMII core in [Chapter 11, “Interfacing to Other Cores.”](#)
- Using the GEMAC core in conjunction with the Ethernet Statistics core in [Chapter 11, “Interfacing to Other Cores”](#)
- 10 Mbps, 100 Mbps, 1 Gbps Ethernet FIFO in [Appendix A, “Using the Client-Side FIFO.”](#)

You can synthesize the entire design using any synthesis tool. The GEMAC core is pre-synthesized and is delivered as an NGC netlist (which appears as a black box to synthesis tools).

Run the Xilinx tools **map**, **par**, and **bitgen** to create a bitstream that can be downloaded to a Xilinx device. Care must be taken to constrain the design correctly, and the UCF produced by the CORE Generator should be used as the basis for the your own UCF. See [Chapter 9, “Constraining the Core,”](#) for more information.

You can simulate the entire design and download the bitstream to the target device.

## Know the Degree of Difficulty

A 1-Gigabit Ethernet MAC implementation is challenging to implement in any technology, and all applications require careful attention to system performance requirements. Pipelining, logic mapping, placement constraints, and logic duplication are all methods that help boost system performance.

See [Table 4-1](#) to determine the relative level of difficulty associated with the Spartan™ and Virtex™ device families. These designs relate to meeting the core required system clock frequency of 125 MHz.

See also [Appendix C, “Calculating DCM Phase-Shifting”](#) to meet Spartan-3 and Spartan-3E setup and hold requirements for external GMII.

**Table 4-1: Degree of Difficulty for Various Implementations**

Device Family	Difficulty
Spartan-3A	Difficult
Spartan-3E	Difficult
Spartan-3	Difficult
Virtex-II	Easy
Virtex-II Pro	Easy
Virtex-4	Easy
Virtex-5	Easy

## Keep it Registered

To simplify timing and increase system performance in an FPGA design, keep all inputs and outputs registered between the user application and the core. This means that all inputs and outputs from the user application should come from, or connect to, a flip-flop. While registering signals may not be possible for all paths, it simplifies timing analysis and makes it easier for the Xilinx tools to place-and-route the design.

## Recognize Timing Critical Signals

The UCF provided with the example design identifies the critical signals and timing constraints that should be applied. See [Chapter 9, “Constraining the Core”](#) for more information.

## Use Supported Design Flows

The core is pre-synthesized and delivered as an NGC netlist. The example implementation scripts provided use XST 9.2i as the synthesis tool for the HDL example design. Other synthesis tools may be used for the user application logic. The core is always unknown to the synthesis tool and should appear as a black box. Note that post synthesis, only ISE 9.2i tools are supported.

## Make Only Allowed Modifications

The GEMAC core should not be modified by the user. Any modifications may have adverse effects on system timing and protocol compliance. Supported user configurations of the GEMAC core can only be made by selecting the options in the CORE Generator when the core is generated. For more information, see [Chapter 3, “Generating the Core.”](#)

## Using the Client Side Data Path

This chapter provides general guidelines for creating designs using the GEMAC core, including a detailed description of each client-side data flow interface of the core.

Definitions of the abbreviations used throughout the remainder of this chapter are defined in [Table 5-1](#).

**Table 5-1: Abbreviations Used in Timing Diagrams**

Abbreviation	Definition
DA	Destination address; 6 bytes
SA	Source address; 6 bytes
L/T	Length/type field; 2 bytes
FCS	Frame check sequence; 4 bytes

### Receiving Inbound Frames

Received Ethernet frames are presented to the client logic on the Receiver subset of the Client-Side Interface. For port definition, see [“Receiver Interface,”](#) on page 25.

#### Normal Frame Reception

[Figure 5-1](#) illustrates the timing of a normal inbound frame transfer. The client must be prepared to accept data at any time; there is no buffering within the MAC to allow for latency in the receive client. Once frame reception begins, data is transferred on consecutive clock cycles to the receive client until the frame is complete. The MAC asserts the `rx_good_frame` signal to indicate that the frame was successfully received and that the frame should be analyzed by the client.

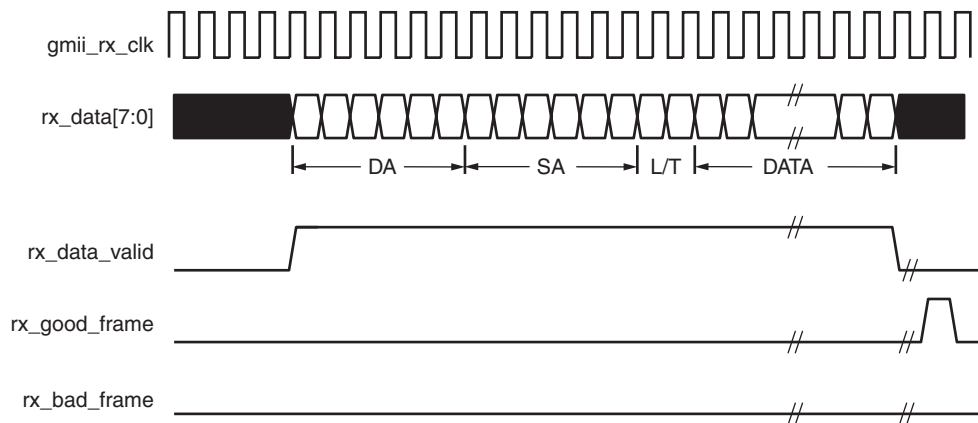


Figure 5-1: Normal Frame Reception

Frame parameters (destination address, source address, length/type and optionally FCS) are supplied on the data bus according to the timing diagram.

If the length/type field in the frame has a length interpretation, indicating that the inbound frame has been padded to meet the Ethernet minimum frame size, the padding is not passed to the client in the data payload. The exception to this is where FCS passing is enabled. See “Client-Supplied FCS Passing.”

When Client-Supplied FCS passing is disabled, `rx_data_valid` is equal to zero between frames for the duration of the padding field (if present), the FCS field, carrier extension (if present), the interframe gap following the frame, and the preamble field of the next frame. When Client-Supplied FCS passing is enabled, `rx_data_valid` is equal to zero between frames for the duration of carrier extension (if present), the interframe gap, and the preamble field of the following frame.

## rx\_good\_frame, rx\_bad\_frame timing

Although the timing diagram (Figure 5-1) shows the `rx_good_frame` signal asserted shortly after the last valid data on `rx_data`, this is not always the case. The `rx_good_frame` or `rx_bad_frame` signals are asserted only after all frame checks are completed. This is after the FCS field has been received (and after reception of carrier extension, if present).

Therefore, either `rx_good_frame` or `rx_bad_frame` is asserted following frame reception at the beginning of the interframe gap.

## Frame Reception with Errors

Figure 5-2 illustrates an unsuccessful frame reception (for example, a fragment frame or a frame with an incorrect FCS). In this case, the `rx_bad_frame` signal is asserted to the client at the end of the frame. It is then the responsibility of the client to drop the data already transferred for this frame.

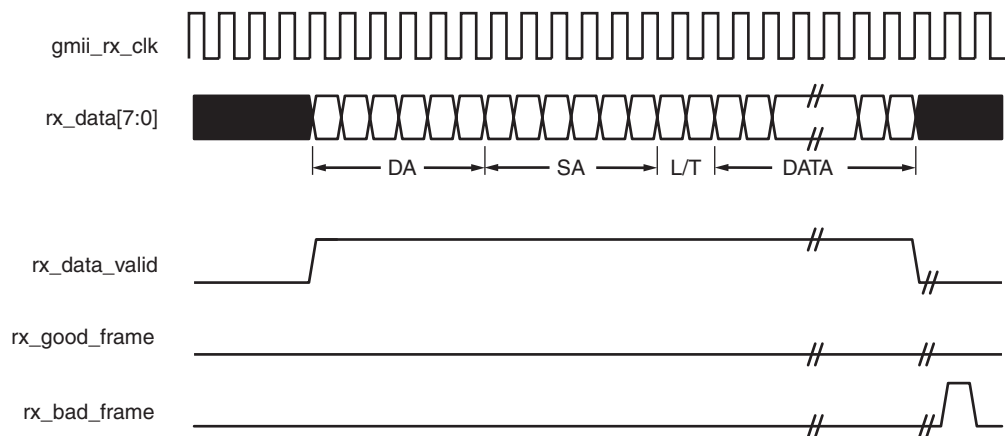


Figure 5-2: Frame Reception with Error

## Client-Supplied FCS Passing

If the GEMAC core is configured to pass the FCS field to the client (see [“Configuration Registers,” on page 71](#)), this is handled as shown in [Figure 5-3](#). In this case, any padding inserted into the frame to meet Ethernet minimum frame length specifications will be left intact and passed to the client.

Note that even though the FCS is passed up to the client, it is also verified by the GEMAC core, and `rx_bad_frame` asserted if the FCS check fails.

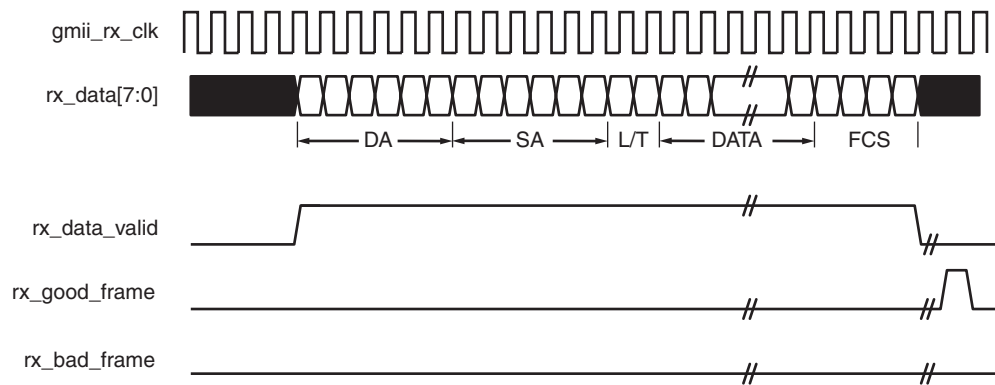


Figure 5-3: Frame Reception with In-Band FCS Field

## VLAN Tagged Frames

[Figure 5-4](#) illustrates the reception of a VLAN tagged frame (if enabled). The VLAN frame is passed to the client so that the frame may be identified as VLAN tagged. This is followed by the Tag Control Information bytes, V1 and V2. More information on the interpretation of these bytes may be found in *IEEE 802.3-2002* standard.

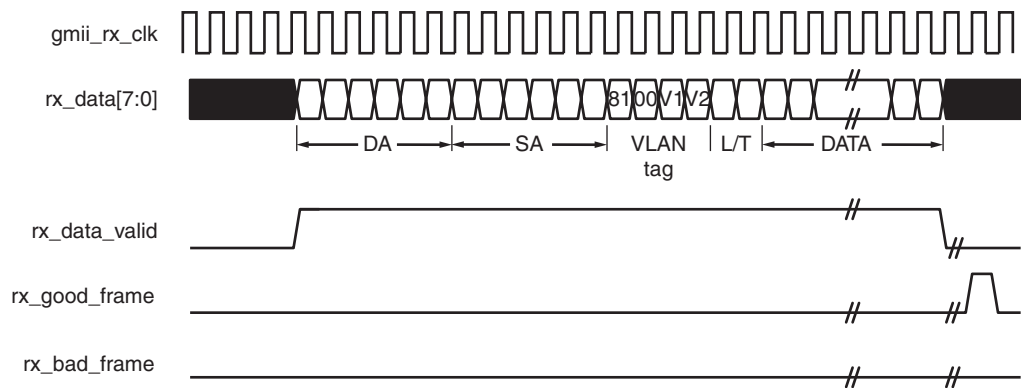


Figure 5-4: Reception of a VLAN Tagged Frame

## Maximum Permitted Frame Length

The maximum legal length of a frame specified in *IEEE 802.3-2002* is 1518 bytes for non-VLAN tagged frames. VLAN tagged frames may be extended to 1522 bytes. When jumbo frame handling is disabled and the core receives a frame which exceeds the maximum legal length, `rx_bad_frame` is asserted. When jumbo frame handling is enabled, frames which are longer than the legal maximum are received in the same way as shorter frames. For more information about enabling and disabling jumbo frame handling, see [“Configuration Registers,”](#) on page 71.

## Length/Type Field Error Checks

### Enabled

Default operation is with the length/type error checking enabled (see [“Receiver Configuration,”](#) on page 72). In this mode, the following checks are made on all frames received. If either of these checks fail, the frame is marked as bad.

A value in the length/type field that is greater than or equal to decimal 46, but less than decimal 1536 (a length interpretation), is checked against the actual data length received.

A value in the length/type field that is less than decimal 46 is checked to see that the data field is padded to exactly 46 bytes (so that the resultant frame is a minimum frame size of 64 bytes total in length).

Furthermore, if padding is indicated (the length/type field is less than decimal 46) and client-supplied FCS passing is disabled, the length value in the length/type field will be used to deassert `rx_data_valid` after the indicated number of data bytes so that the padding bytes are removed from the frame. See [“Client-Supplied FCS Passing.”](#)

### Disabled

When the length/type error checking is disabled and the length/type field has a length interpretation, the MAC does not check the length value against the actual data length received. See [“Receiver Configuration”](#) in Chapter 8. A frame containing only this error is marked as good.

However, if the length/type field is less than decimal 46 then the MAC will mark a frame as bad if it is not the minimum frame size of 64 bytes.



If padding is indicated and client-supplied FCS passing is disabled, then a length value in the length/type field will not be used to deassert `rx_data_valid`. Instead, `rx_data_valid` is deasserted before the start of the FCS field, and any padding is not removed from the frame.

## Address Filter

If the optional Address Filter is included in the core, the MAC is able to reject frames that do not contain a known address in their destination address field. If a frame is rejected, the `rx_data_valid` signal is not asserted for the duration of the frame. In addition, neither `rx_good_frame` or `rx_bad_frame` are asserted at the end of the frame. The statistics vectors are still output with a valid pulse at the end of the rejected frame.

If the Address Filter is not in promiscuous mode, it will reject frames in which the destination address does not meet any of the following criteria:

- It is equal to the broadcast address defined in the *IEEE 802.3-2002* specification.
- It is equal to the pause multicast address defined in the *IEEE 802.3-2002* specification.
- The destination address field contains the pause frame MAC source address specified in the Receiver Configuration Word 0 and Word 1.
- It is equal to the MAC unicast address. When the optional Management Interface is present, this is found in the unicast address configuration registers ([Table 8-8](#) and [Table 8-9, page 75](#)). If the Management Interface is not present the unicast address is input on the `mac_unicast_address` input.
- It matches any of the addresses stored in the MAC address table. The address table is only present when the MAC contains the optional Management Interface and the core was built with one or more address table entries.

## Receiver Statistics Vector

The statistics for the frame received are contained within the `rx_statistics_vector`. The vector is driven synchronously by the receiver clock, `gmii_rx_clk`, following frame reception. The bit field definition for the vector is defined in [Table 5-2](#).

All bit fields, with the exception of byte valid, are valid only when the `rx_statistics_valid` is asserted. This is illustrated in [Figure 5-5](#). Byte valid is significant on every `gmii_rx_clk` cycle.

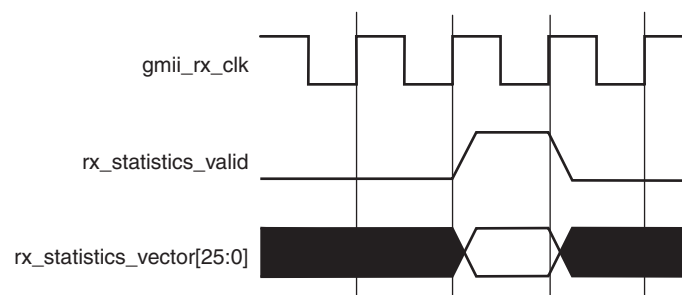


Figure 5-5: Receiver Statistics Vector Timing

Table 5-2: Bit Definition for the Receiver Statistics Vector

rx_statistics_vector bit(s)	Name	Description
26	Address Match	If the optional Address Filter is included in the core, this bit is asserted if the address of the incoming frame matches one of the stored or pre-set addresses in the Address Filter. If the Address Filter is omitted from the core, or is configured in promiscuous mode, this line is held high.
25	Length/Type Out of Range	If the length/type field contained a length value that did not match the number of MAC client data bytes received and the length/type field checks are enabled, then this bit is asserted.  This bit is also asserted if the length/type field is less than 46 and the frame is not padded to exactly 64 bytes. This is independent of whether or not the length/type field checks are enabled.
24	Bad Opcode	Asserted if the previous frame was error-free and contained the special control frame identifier in the length/type field, but contained an opcode that is unsupported by the MAC (any opcode other than Pause).
23	Flow Control Frame	Asserted if the previous frame met all the following conditions: <ul style="list-style-type: none"> <li>• error-free</li> <li>• contained the special control frame identifier in the length/type field</li> <li>• contained a destination address that matched either the MAC Control Multicast Address or the configured source address of the MAC</li> <li>• contained the supported Pause opcode</li> <li>• was acted upon by the MAC</li> </ul>
22	Byte Valid	Asserted if a MAC frame byte (DA to FCS inclusive) is in the process of being received. This is valid on every clock cycle.  Do not use this as an enable signal to indicate that data is present on rx_data.
21	VLAN frame	Asserted if the previous frame contained a VLAN identifier in the length/type field when receiver VLAN operation is enabled.

Table 5-2: Bit Definition for the Receiver Statistics Vector

rx_statistics_vector bit(s)	Name	Description
20	Out of Bounds	Asserted if the previous frame exceeded the specified <i>IEEE802.3-2002</i> maximum legal length (see <a href="#">“Maximum Permitted Frame Length”</a> ). This is only valid if jumbo frames are disabled.
19	Control Frame	Asserted if the previous frame contained the special control frame identifier in the length/type field.
18:5	Frame Length	The length of the previous frame in number of bytes. The count will stick at 16,383 for any jumbo frames larger than this value.
4	Multicast Frame	Asserted if the previous frame contained a multicast address in the destination address field.
3	Broadcast Frame	Asserted if the previous frame contained the broadcast address in the destination address field.
2	FCS Error	Asserted if the previous frame received had an incorrect FCS value or the MAC detected error codes during frame reception.
1	Bad Frame	Asserted if the previous frame received contained errors.
0	Good Frame	Asserted if the previous frame received was error-free.

## Transmitting Outbound Frames

Ethernet frames to be transmitted are presented to the client logic on the Transmitter subset of the Client-Side Interface. For port definition, see [“Transmitter Interface,”](#) on page 24.

### Normal Frame Transmission

[Figure 5-6](#) illustrates the timing of a normal outbound frame transfer. When the client wishes to transmit a frame, it places the first column of data onto the `tx_data` port and asserts a ‘1’ onto `tx_data_valid`.

When the GEMAC core has read this first byte of data, and in accordance with flow control requests and interpacket gap requirements, it will assert the `tx_ack` signal; on the next and subsequent rising clock edges, the client must provide the remainder of the data for the frame.

The end of frame is signalled to the GEMAC core by taking `tx_data_valid` low.

For maximum flexibility in switching and routing applications, the Ethernet frame parameters (destination address, source address, length/type and optionally FCS) are encoded within the same data stream that the frame payload is transferred upon, rather

than on separate ports. This is illustrated in the timing diagrams. Definitions of the abbreviations used in the timing diagrams are defined in [Table 5-1](#).

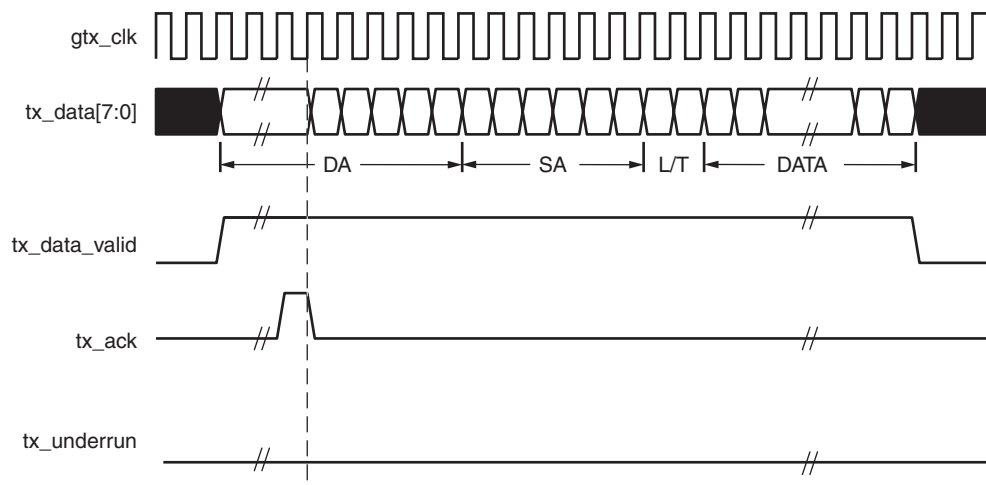


Figure 5-6: Normal Frame Transmission

## Padding

When fewer than 46 bytes of data are supplied by the client to the GEMAC core, the transmitter module will add padding up to the minimum frame length. The exception to this is when the GEMAC core is configured for client-passed FCS; in this case, the client must also supply the padding to maintain the minimum frame length. See [“Client-Supplied FCS Passing”](#) for more information.

## Client-Supplied FCS Passing

The transmission timing depicted in [Figure 5-7](#) shows the GEMAC core configured to have the FCS field passed in by the client. In this case, it is the responsibility of the client to ensure that the frame meets the Ethernet minimum frame length requirements as the GEMAC core will not perform any padding of the payload. See [“Configuration Registers,”](#) on page 71 for more information.

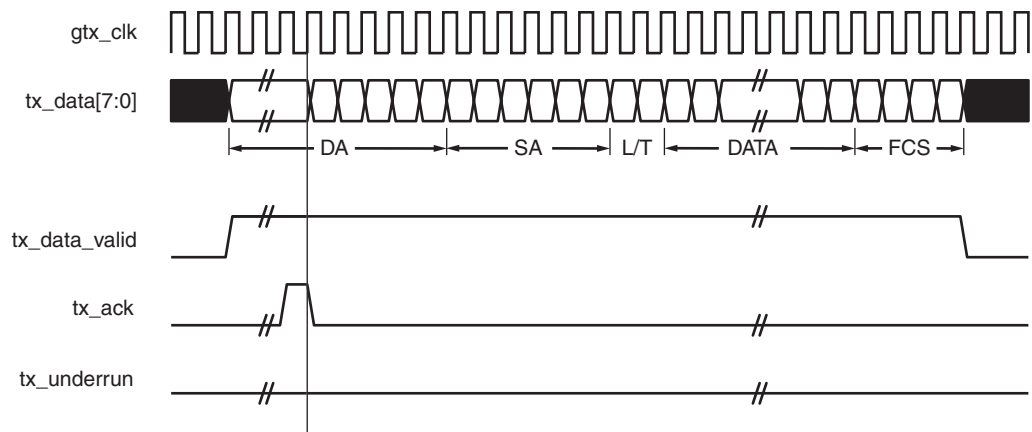


Figure 5-7: Frame Transmission with Client-supplied FCS

## Client Underrun

Figure 5-8 illustrates the timing of an aborted transfer. An example of this situation is a FIFO connected to the client interface that empties before a frame transfer is complete. When the client asserts `tx_underrun` during a frame transmission, the GEMAC core inserts an error code to corrupt the current frame, and then falls back to idle transmission. It is the responsibility of the client to re-queue the aborted frame for transmission.

When an underrun occurs, `tx_data_valid` may be asserted on the clock cycle after the `tx_underrun` assertion to request a new transmission.

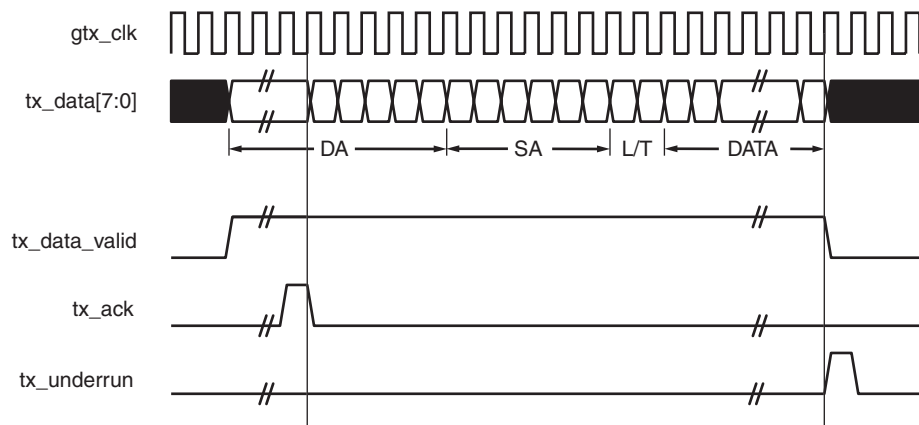


Figure 5-8: Frame Transmission with Underrun

## VLAN Tagged Frames

Figure 5-9 illustrates transmission of a VLAN tagged frame (if enabled). Note that the handshaking signals across the interface do not change; however, the VLAN type tag 81-00 must be supplied by the client to signify that the frame is VLAN tagged. The client also supplies the two bytes of Tag Control Information, V1 and V2, at the appropriate times in the data stream. More information on the contents of these two bytes can be found in *IEEE*

802.3-2002. For more information about enabling and disabling jumbo frame handling, see “Configuration Registers,” on page 71.

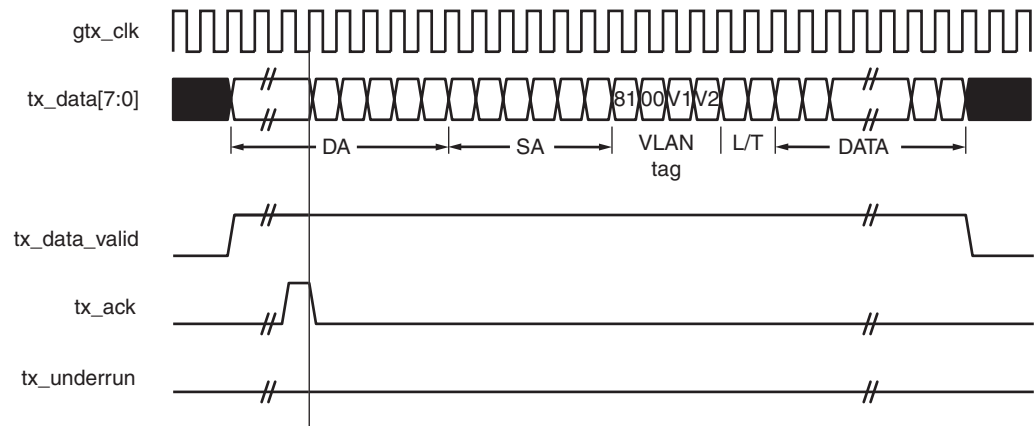


Figure 5-9: Transmission of a VLAN Tagged Frame

## Maximum Permitted Frame Length

The maximum legal length of a frame specified in *IEEE 802.3-2002* is 1518 bytes for non-VLAN tagged frames. VLAN tagged frames may be extended to 1522 bytes. When jumbo frame handling is disabled and the client attempts to transmit a frame that exceeds the maximum legal length, the GEMAC core will insert an error code to corrupt the current frame, and the frame will be truncated to the maximum legal length. When jumbo frame handling is enabled, frames longer than the legal maximum are transmitted error-free. For more information on enabling and disabling Jumbo frame handling, see “Configuration Registers,” on page 71.

## Inter-Frame Gap Adjustment

A configuration bit in the transmitter control register (see “Configuration Registers,” on page 71) allows the user to control the length of the inter-frame gap transmitted by the MAC on the physical interface. If this function is selected, the MAC exerts back pressure on the client interface to delay the transmission of the next frame until the requested number of idle cycles has elapsed. The number of idle cycles is controlled by the value on the tx\_ifg\_delay port seen at the start of frame transmission on the client interface. Figure 5-10 shows the MAC operating in this mode.

Reducing the interframe gap to below the *IEEE 802.3-2002* minimum of 12 idles is supported, but the MAC will transmit an absolute minimum of 4 idles. If the Ethernet Statistics core is used with the MAC, then accuracy cannot be guaranteed if the interframe gap adjustment is set to less than 12 idles. However, the tx\_statistic\_vector and rx\_statistic\_vector values will always remain correct.

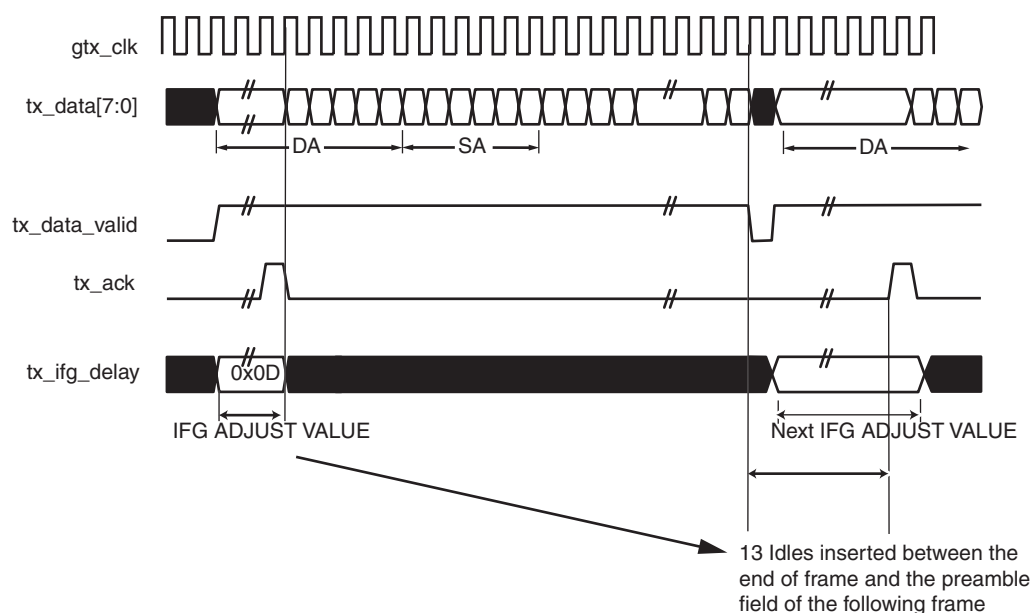


Figure 5-10: Inter-Frame Gap Adjustment

## Transmitter Statistics Vector

The statistics for the transmitted frame are contained within the `tx_statistic_vector`. The vector is driven synchronously by the transmitter clock, `gtx_clk`, following frame transmission. The bit field definition for the vector is defined in [Table 5-3](#).

All bit fields, with the exception of byte valid, are valid only when `tx_statistic_valid` is asserted ([Figure 5-11](#)). Byte valid is significant on every `gtx_clk` cycle.

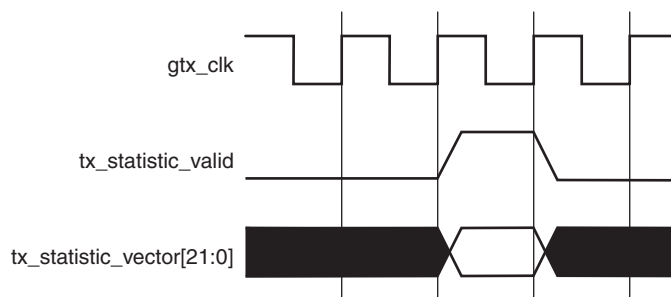


Figure 5-11: Transmitter Statistic Vector Timing

Table 5-3: Bit Definition for the Transmitter Statistics Vector

tx_statistics_vector bit(s)	Name	Description
21	Pause Frame	Asserted if the previous frame was a pause frame that the MAC itself initiated in response to a pause_req assertion.
20	Byte Valid	Asserted if a MAC frame byte (DA to FCS inclusive) is in the process of being transmitted. This is valid on every clock cycle.  Do not use this as an enable signal to indicate that data is present on gmii_txd.
19	VLAN Frame	Asserted if the previous frame contained a VLAN identifier in the length/type field when transmitter VLAN operation is enabled.
18:5	Frame Length	The length of the previous frame in number of bytes. The count will stick at 16,383 for any jumbo frames larger than this value.
4	Control Frame	Asserted if the previous frame had the special MAC Control Type code 88-08 in the length/type field.
3	Underrun Frame	Asserted if the previous frame contained an underrun error.
2	Multicast Frame	Asserted if the previous frame contained a multicast address in the destination address field.
1	Broadcast Frame	Asserted if the previous frame contained a broadcast address in the destination address field.
0	Successful Frame	Asserted if the previous frame was transmitted without error.



## Using Flow Control

This chapter describes the operation of the flow-control logic of the GEMAC core. The flow control block is designed to clause 31 of the *IEEE 802.3-2002* standard. The MAC may be configured to transmit pause requests and to act on their reception; these modes of operation can be independently enabled or disabled. See [“Flow Control Configuration,”](#) on page 74 for more information.

### Overview of Flow Control

#### Flow Control Requirement

Figure 6-1 illustrates the requirements for Flow Control.

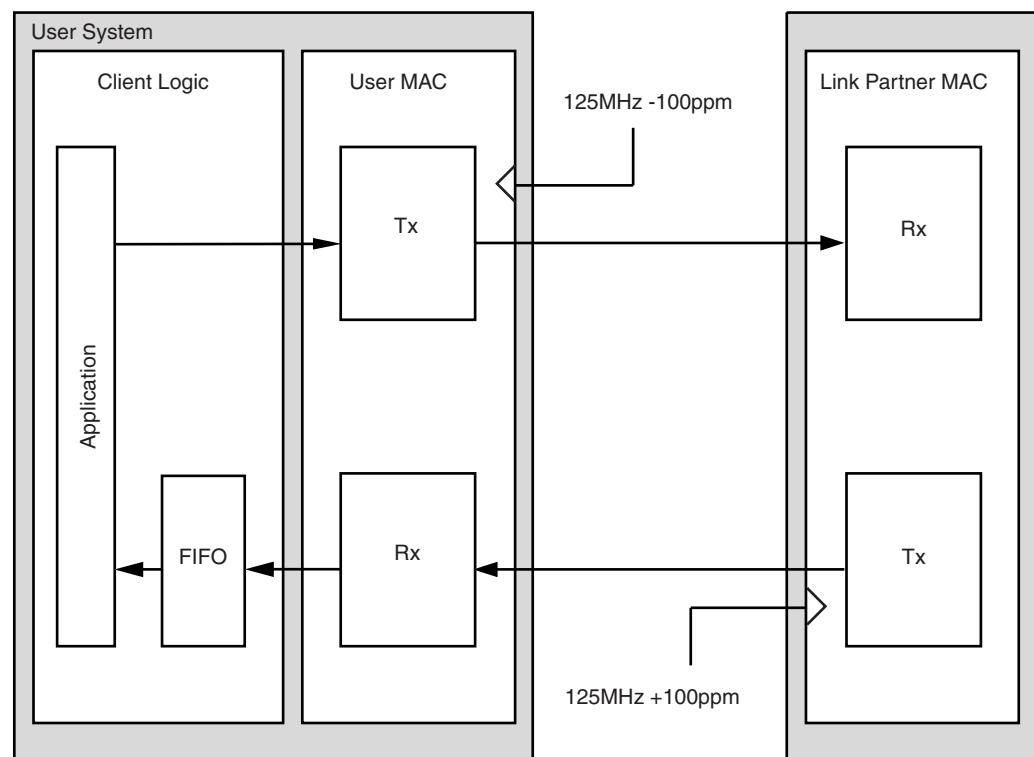


Figure 6-1: Requirement for Flow Control

The user MAC on the left side has a reference clock slightly slower than the nominal 125 MHz. The link partner MAC on the right side has a reference clock slightly faster than

the nominal 125 MHz. As a result, the user MAC receives data at a faster line rate than that at which it can transmit. The MAC on the left is shown performing a loopback implementation which results in the FIFO filling up over time. Without Flow Control, this FIFO will eventually fill and overflow, resulting in the corruption or loss of Ethernet frames. Enabling Flow Control in the MAC provides a mechanism to solve this data rate matching problem.

## Flow Control Basics

A MAC may transmit a pause control frame to request that its link partner cease transmission for a defined period of time. For example, the user MAC on the left side of [Figure 6-1](#) may initiate a pause request when its client FIFO (illustrated) reaches a nearly full state.

A MAC should respond to received pause control frames by ceasing transmission of frames for the period of time defined in the received pause control frame. For example, the link partner MAC in [Figure 6-1](#) may cease transmission after receiving the pause control frame transmitted by the user MAC. In a well designed system, the link partner MAC would cease transmission before the client FIFO experienced an overflow condition. This provides time for the FIFO to be emptied to a safe level before normal operation resumes, thus safeguarding the system against FIFO overflow conditions and frame loss.

## Pause Control Frames

Control frames are a unique type of Ethernet frame, defined in clause 31 of the *IEEE 802.3-2002* standard. Control frames are differentiated from other frame types by a defined value placed in the length/type field (MAC Control Type code). [Figure 6-2](#) illustrates the control frame format.

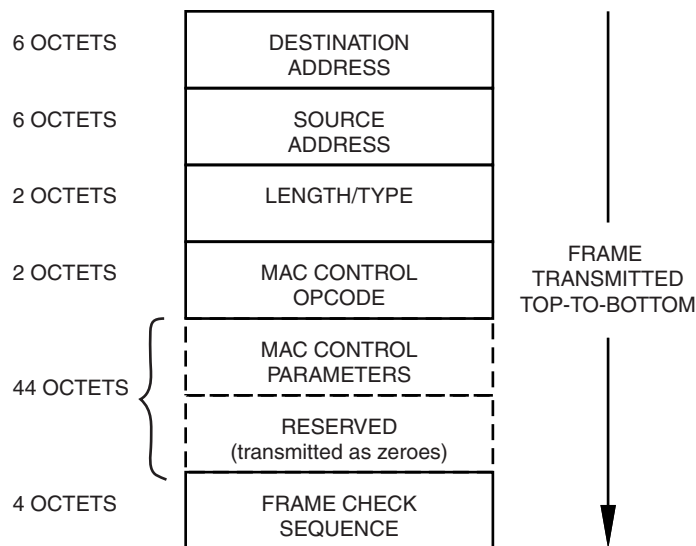


Figure 6-2: MAC Control Frame Format

A pause control frame is a unique type of control frame, identified by a defined value placed in the MAC Control opcode field.

**Note:** MAC Control opcodes other than for pause (flow control) frames have recently been defined for Ethernet Passive Optical Networks.

The MAC control parameter field of the pause control frame contains a 16-bit field containing a binary value directly related to the pause duration. This defines the number of `pause_quantum` (512 bit times of the particular implementation). For 1-Gigabit Ethernet, a single `pause_quantum` corresponds to 512 ns.

## Flow Control Operation of the GEMAC

### Transmitting a PAUSE Control Frame

#### Core-initiated Pause Request

If the GEMAC core is configured to support transmit flow control, the client can initiate a pause control frame by asserting `pause_req` (see “Flow Control Configuration,” on page 74). Figure 6-3 illustrates pause request timing.

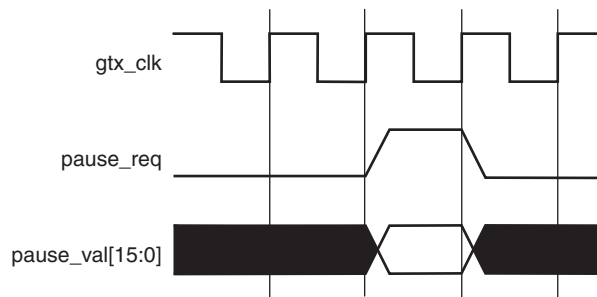


Figure 6-3: Pause Request Timing

This action causes the core to construct and transmit a pause control frame on the link with the following MAC Control frame parameters (Figure 6-2):

- The destination address used is an *IEEE802.3* globally assigned multicast address (to which any flow control-capable MAC will respond).
- The source address used is the configurable pause frame MAC address (see “Receiver Configuration,” on page 72).
- The value sampled from the `pause_val[15:0]` port at the time of the `pause_req` assertion will be encoded into the MAC control parameter field to select the duration of the pause (in units of *pause\_quantum*).

If the transmitter is inactive at the time of the pause request, this pause control frame is transmitted immediately. If the transmitter is currently busy, the current frame being transmitted is allowed to complete, followed by the pause control frame (in preference to any pending client-supplied frame).

A pause control frame initiated by this method is transmitted even if the transmitter has ceased in response to receiving an inbound pause request.

**Note:** Only a single pause control frame request is stored by the transmitter. If `pause_req` is asserted numerous times in a short time period (before the control pause frame transmission has

begun), only a single pause control frame is transmitted. The most recent value sampled will be the `pause_val[15:0]` value used.

## Client Initiated Pause Request

For maximum flexibility, flow control logic can be disabled in the core and alternatively implemented in the client logic connected to the core (see [“Flow Control Configuration,” on page 74](#)). Any type of control frame can be transmitted through the core through the client interface using the same transmission procedure as a standard Ethernet frame (see [“Transmitting Outbound Frames,” on page 43](#)).

## Receiving a Pause Control Frame

### Core Initiated Response to a Pause Request

An error free control frame is a received frame matching the format of [Figure 6-2](#). It must pass all standard receiver frame checks (for example, FCS field checking). In addition, the control frame received must be exactly 64-bytes in length (from destination address through to the FCS field inclusive: this is minimum legal Ethernet MAC frame size and the defined size for control frames).

Any control frame received that does not conform to these checks contains an error and is passed to the receiver client with the `rx_bad_frame` signal asserted.

### Pause Frame Reception Disabled

When pause control reception is disabled, an error free control frame is received through the client interface with `rx_good_frame` asserted (see [“Flow Control Configuration,” on page 74](#)). In this way, the frame is passed to the client logic for interpretation (see [“Client Initiated Response to a Pause Request,” on page 53](#)).

### Pause Frame Reception Enabled

When pause control reception is enabled, and an error-free frame is received by the GEMAC core (see [“Flow Control Configuration,” on page 74](#)), the following frame decoding functions are performed:

- The destination address field is matched against the *IEEE 802.3* globally assigned multicast address or the configurable pause frame MAC address (see [“Configuration Registers,” on page 71](#)).
- The length/type field is matched against the MAC control type code.
- The opcode field contents are matched against the Pause opcode.

If any of the previously described checks are false, the frame is ignored by the Flow Control logic and passed up to the client logic for interpretation by marking it with `rx_good_frame` asserted. It is then the responsibility of the MAC client logic to decode, act on (if required), and drop this control frame.

If all the previously described checks are true, the 16-bit binary value in the MAC Control Parameters field of the control frame is then used to inhibit transmitter operation for the required number of *pause\_quantum*. This inhibit is implemented by delaying the assertion of `tx_ack` at the transmitter client interface until the requested pause duration has expired. The received pause frame is then passed on to the client with `rx_bad_frame` asserted to indicate to the client that the pause frame can be dropped.

**Note:** Any frame in which the length/type field contains the MAC control type should be dropped by the receiver client logic. All control frames are indicated by `rx_statistic_vector` bit 19 (see “Receiver Statistics Vector,” on page 41).

## Client Initiated Response to a Pause Request

For maximum flexibility, flow control logic can be disabled in the core and alternatively implemented in the client logic connected to the core (see “Flow Control Configuration,” on page 74). Any type of error-free control frame is then passed through the core with `rx_good_frame` asserted. The frame is passed to the client for interpretation. It is then the responsibility of the client to drop this control frame and to act on it by ceasing transmission through the core, if applicable.

## Flow Control Implementation Example

This section provides a basic overview of a Flow Control implementation, using Figure 6-1 as a sample. To summarize the example, the user MAC on the left hand side of the figure cannot match the full line rate of the link partner MAC on the right hand side due to clock tolerances. Over time, the FIFO illustrated will fill and overflow. The goal is to implement a flow control method which will (over a long time period) reduce the average line rate of the link partner MAC to that of the user MAC.

### Method

1. Choose a FIFO nearly full to occupancy threshold (7/8 occupancy is used in this description—but the choice of threshold is implementation specific). When the occupancy of the FIFO exceeds this occupancy, initiate a single pause control frame, from the user MAC, with 0xFFFF used as the *pause\_quantum* duration (0xFFFF is placed on `pause_val[15:0]`). This is the maximum pause duration. This causes the link partner MAC to cease transmission, and the FIFO of the user system will start to empty.
2. Choose a second FIFO occupancy threshold (3/4 is used in this description—but the choice of threshold is implementation specific). When the occupancy of the FIFO falls below this occupancy, initiate a second pause control frame from the user MAC, with 0x0000 used as the *pause\_quantum* duration (0x0000 is placed on `pause_val[15:0]`). This indicates a zero pause duration, and upon receiving this pause control frame, the link partner MAC immediately resumes transmission (it does not wait for the original requested pause duration to expire). This pause control frame can therefore be considered a “pause cancel” command.

## Operation

Figure 6-4 illustrates the FIFO occupancy over a period of time.

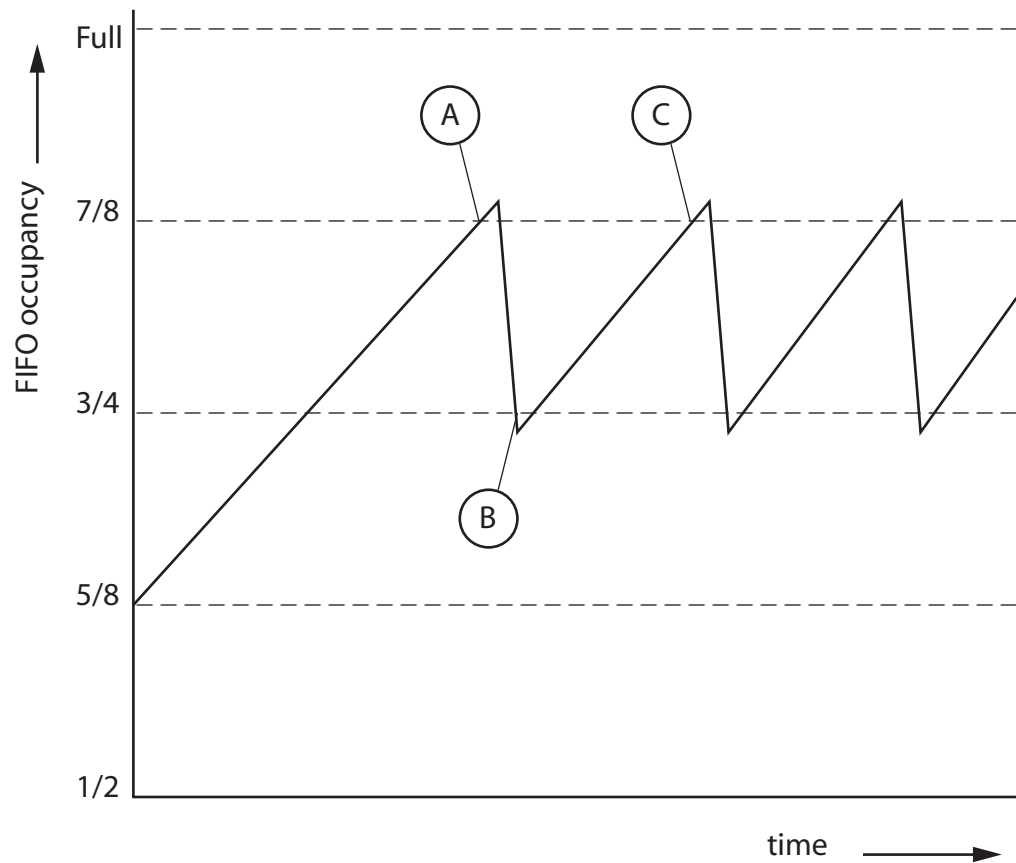


Figure 6-4: Flow Control Implementation Triggered from FIFO Occupancy

The following describes the sequence of flow control operation.

1. The average FIFO occupancy of the user system gradually increases over time due to the clock tolerances. At point A, the occupancy has reached the threshold of  $7/8$  occupancy. This triggers the maximum duration pause control frame request.
2. On receiving the pause control frame, the link partner MAC ceases transmission.
3. After the link partner MAC ceases transmission, the occupancy of the FIFO in the user system rapidly empties. The occupancy falls to the second threshold of  $3/4$  occupancy at point B. This triggers the zero duration pause control frame request (the pause cancel command).
4. On receiving this second pause control frame, the link partner MAC resumes transmission.
5. Normal operation resumes and the FIFO occupancy again gradually increases over time. At point C, this Flow Control cycle repeats.

## Using the Physical Side Interface

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This chapter provides general guidelines for creating designs using the Physical Side Interface of the GEMAC core. The physical side interface implements GMII-style signaling and is typically attached to a physical layer device (PHY), either off-chip or internally integrated. See [“Physical Side Interface” in Chapter 2](#) for more information. For information about using an internal interface in conjunction with the Ethernet 1000BASE-X PCS/PMA or SGMII core, see [Chapter 11, “Interfacing to Other Cores.”](#)

The remainder of this chapter describes how to use the core with an external GMII or RGMII. See also [Chapter 9, “Constraining the Core”](#) for a listing of required constraints.

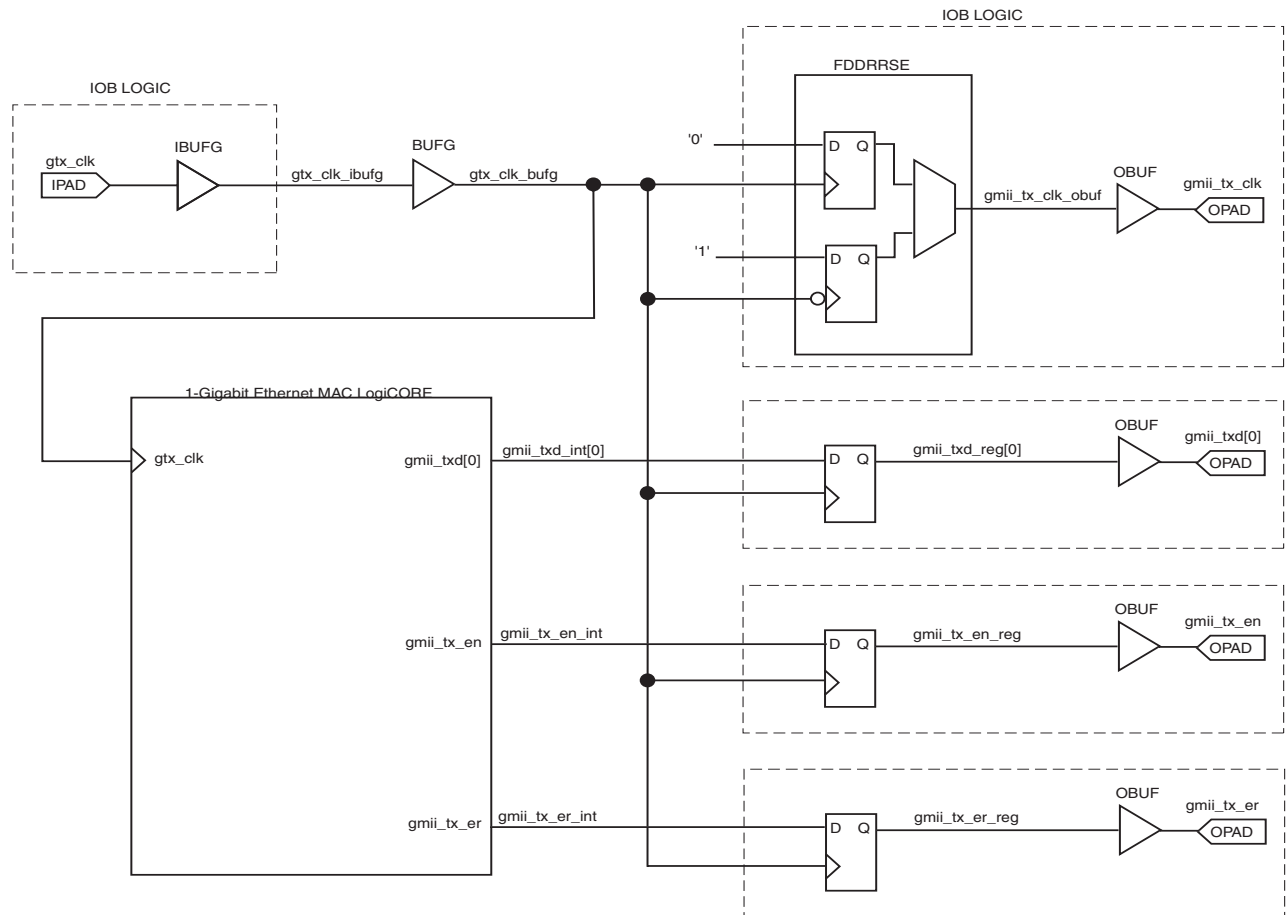
### Implementing External GMII

The HDL example design that is delivered with the core will implement an external GMII when GMII is selected from the CORE Generator GUI (see [Chapter 3, “Generating the Core”](#)). For more information about the example design, see the *1-Gigabit Ethernet MAC Getting Started Guide*.

#### GMII Transmitter Logic

[Figure 7-1](#) illustrates how to use the physical transmitter interface of the core to create an external GMII in a Virtex-II device. The signal names and logic shown in this figure exactly match those delivered with the example design when the GMII is selected. If other families are chosen, equivalent primitives and logic specific to that family is used in the example design.

[Figure 7-1](#) shows that the output transmitter signals are registered in device IOBs before driving them to the device pads. The logic required to forward the transmitter clock is also shown. This logic uses an IOB output Double-Data-Rate (DDR) register so that the clock signal produced incurs exactly the same delay as the data and control signals. This clock signal, `gmii_tx_clk`, is inverted with respect to `gtx_clk` so that the rising edge of `gmii_tx_clk` will occur in the centre of the data valid window, therefore maximizing setup and hold times across the interface.





## GMII Receiver Logic

### Virtex-II Pro and Virtex-II Devices

Figure 7-2 illustrates how to use the physical receiver interface of the core to create an external GMII in a Virtex-II device. The signal names and logic shown on the figure exactly match those delivered with the example design when the GMII is chosen. If other families are chosen, equivalent primitives and logic specific to that family will automatically be used in the example design.

This figure also shows that the input receiver signals are registered in device IOBs before driving them to the device pads. This logic achieves the required setup and hold times across the interface in all these device families.

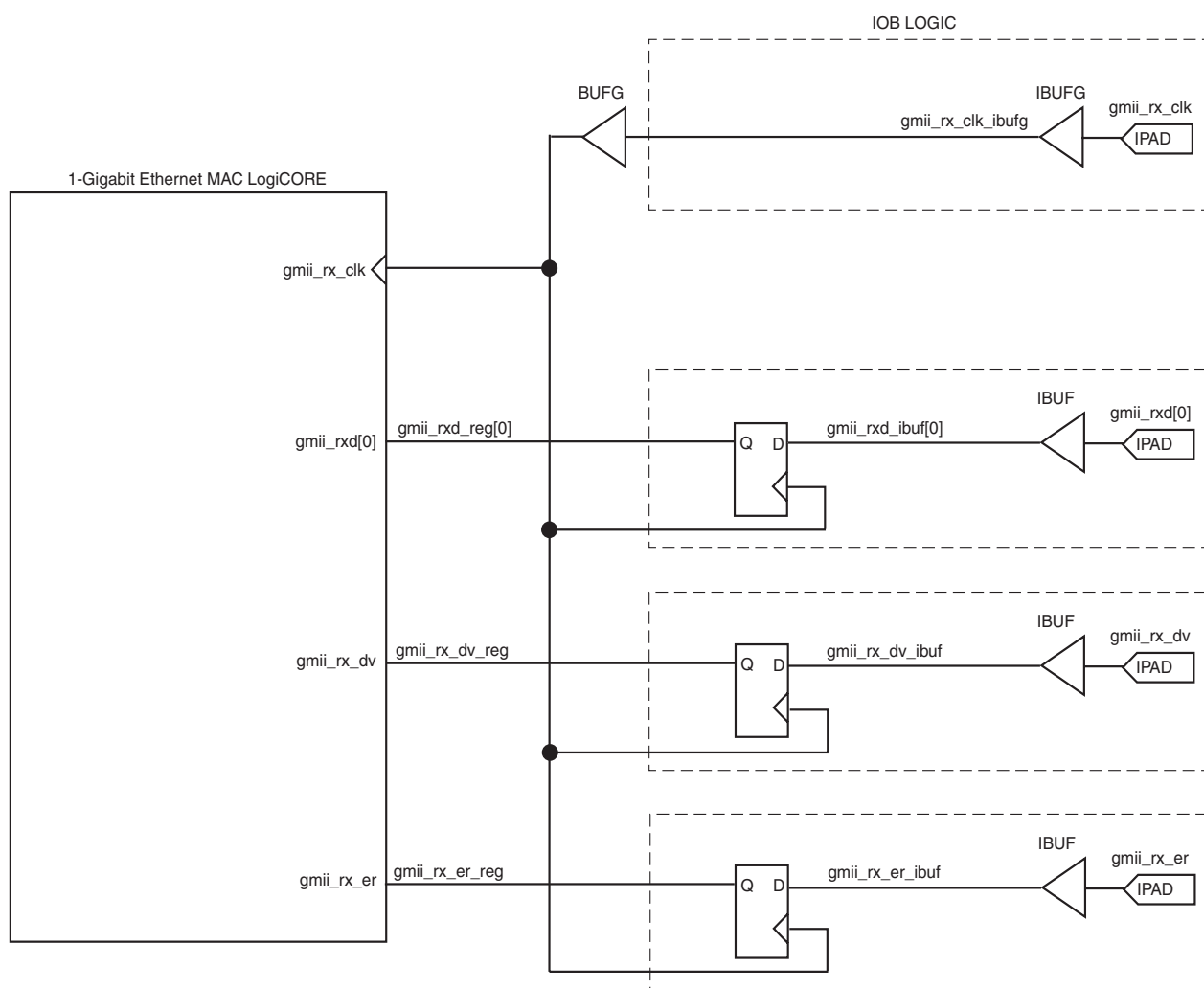


Figure 7-2: External GMII Receiver Logic

## Spartan-3, Spartan-3E and Spartan-3A Devices

The logic required to implement the transmitter GMII remains identical to that described in “GMII Transmitter Logic,” on page 55. However, the logic described in “GMII Receiver Logic” does not meet the input setup and hold requirements for GMII with Spartan-3, Spartan-3E and Spartan-3A devices. A DCM must be used on the `gmii_rx_clk` clock path, as illustrated in Figure 7-3. This is performed by the example design delivered with the core for Spartan-3, Spartan-3E and Spartan-3A devices (all signal names and logic match Figure 7-3). This DCM circuitry may optionally be used in other families.

Phase-shifting may then be applied to the DCM to fine-tune the setup and hold times at the GMII IOB input flip-flops. Fixed phase-shift is applied to the DCM with the example UCF for the example design. See Appendix C, “Calculating DCM Phase-Shifting.”

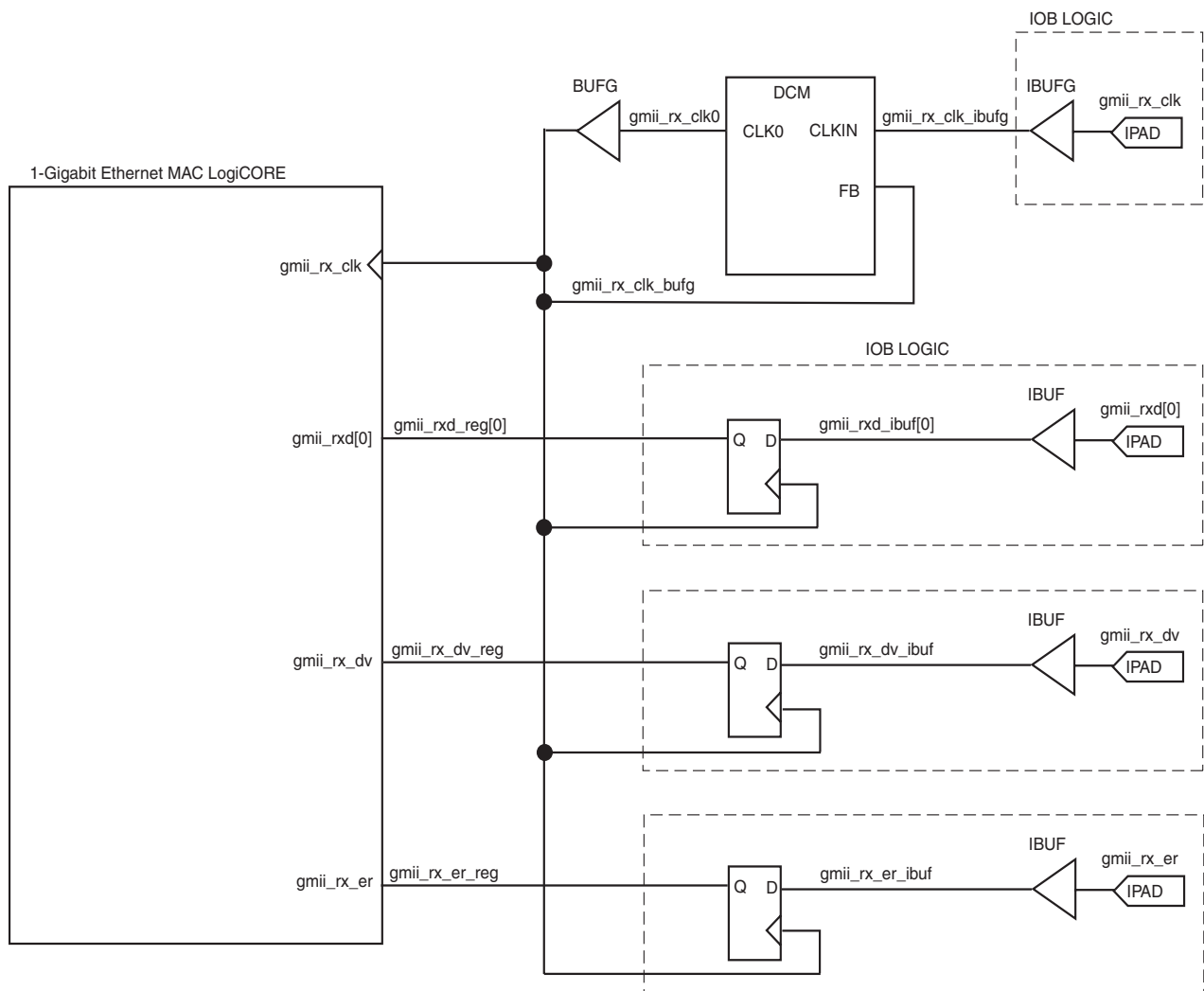


Figure 7-3: External GMII Receiver Logic for Spartan-3 and Spartan-3E Devices

## Virtex-4 Devices

The logic required to implement the transmitter GMII is identical to that described in “GMII Transmitter Logic,” on page 55. However, the logic described in “GMII Receiver Logic” does not meet the input setup and hold requirements for GMII with Virtex-4 devices. An IDELAY component may be used on the clock, data and control paths, as illustrated in Figure 7-4. These can be used to either shift the input clock `gmii_rx_clk` or the data and control signals to meet the setup and hold requirements and to allow for any bus skew across the data and control inputs. The IDELAY components are used in fixed delay mode, where the attribute `IOBDELAY_VALUE` determines the tap delay value. An IDELAYCTRL primitive must be instantiated for this mode of operation. See the *Virtex-4 User Guide* for more information about the IDELAYCTRL and IDELAY components.

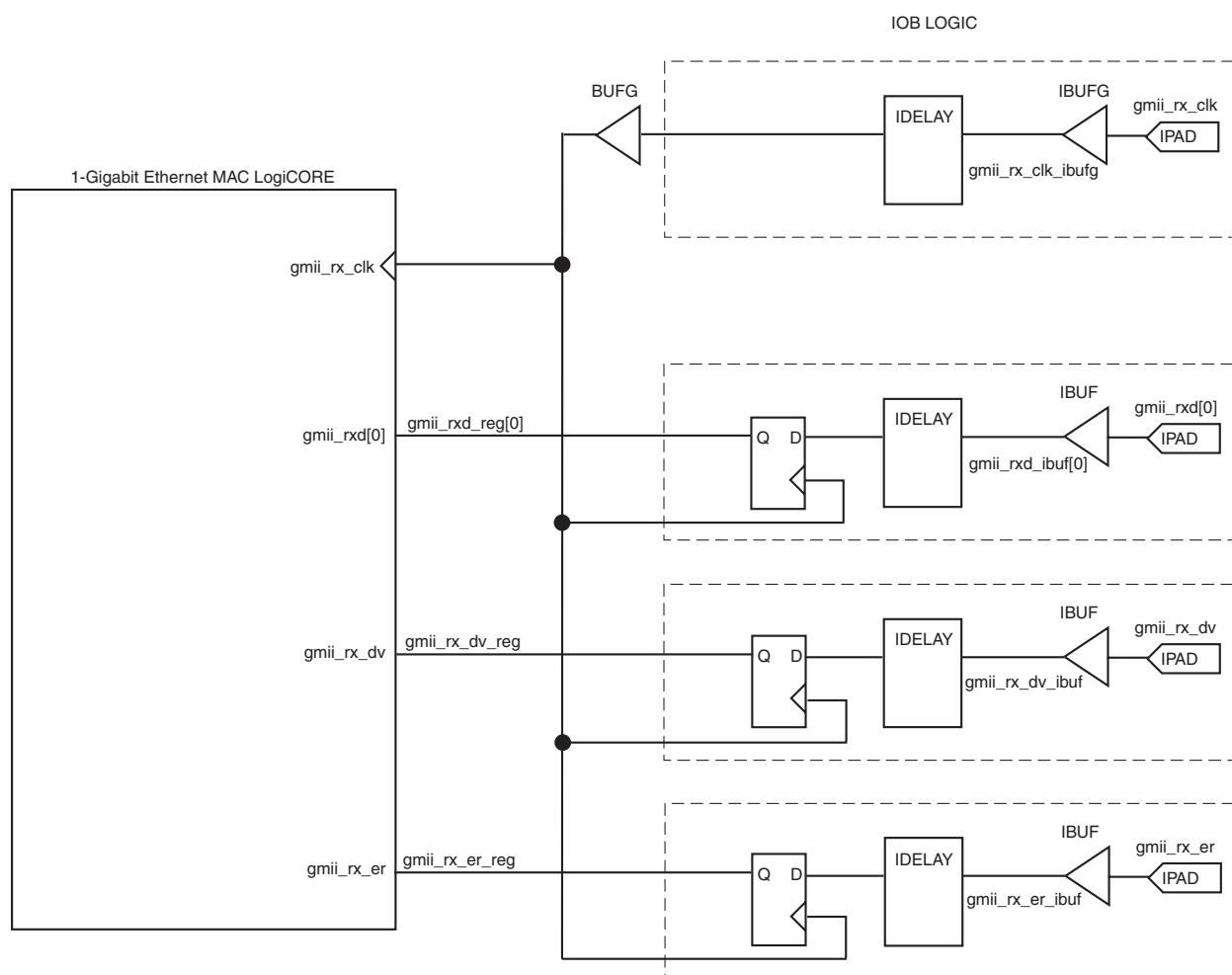


Figure 7-4: External GMII Receiver Logic for Virtex-4 Devices

## Virtex-5 Devices

The logic required to implement the transmitter GMII is identical to that described in “GMII Transmitter Logic,” on page 55. However, the logic described in “GMII Receiver Logic” does not meet the input setup and hold requirements for GMII with Virtex-5 devices. An IODELAY component may be used on the clock, data and control paths, as illustrated in Figure 7-5. These can be used to either shift the input clock gmii\_rx\_clk or the data and control signals to meet the setup and hold requirements and to allow for any bus skew across the data and control inputs. The IODELAY components are used in fixed delay mode, where the attribute IDELAY\_VALUE determines the tap delay value. An IDELAYCTRL primitive must be instantiated for this mode of operation. Refer to the *Virtex-5 User Guide* for more information on the use of IDELAYCTRL and IODELAY components.

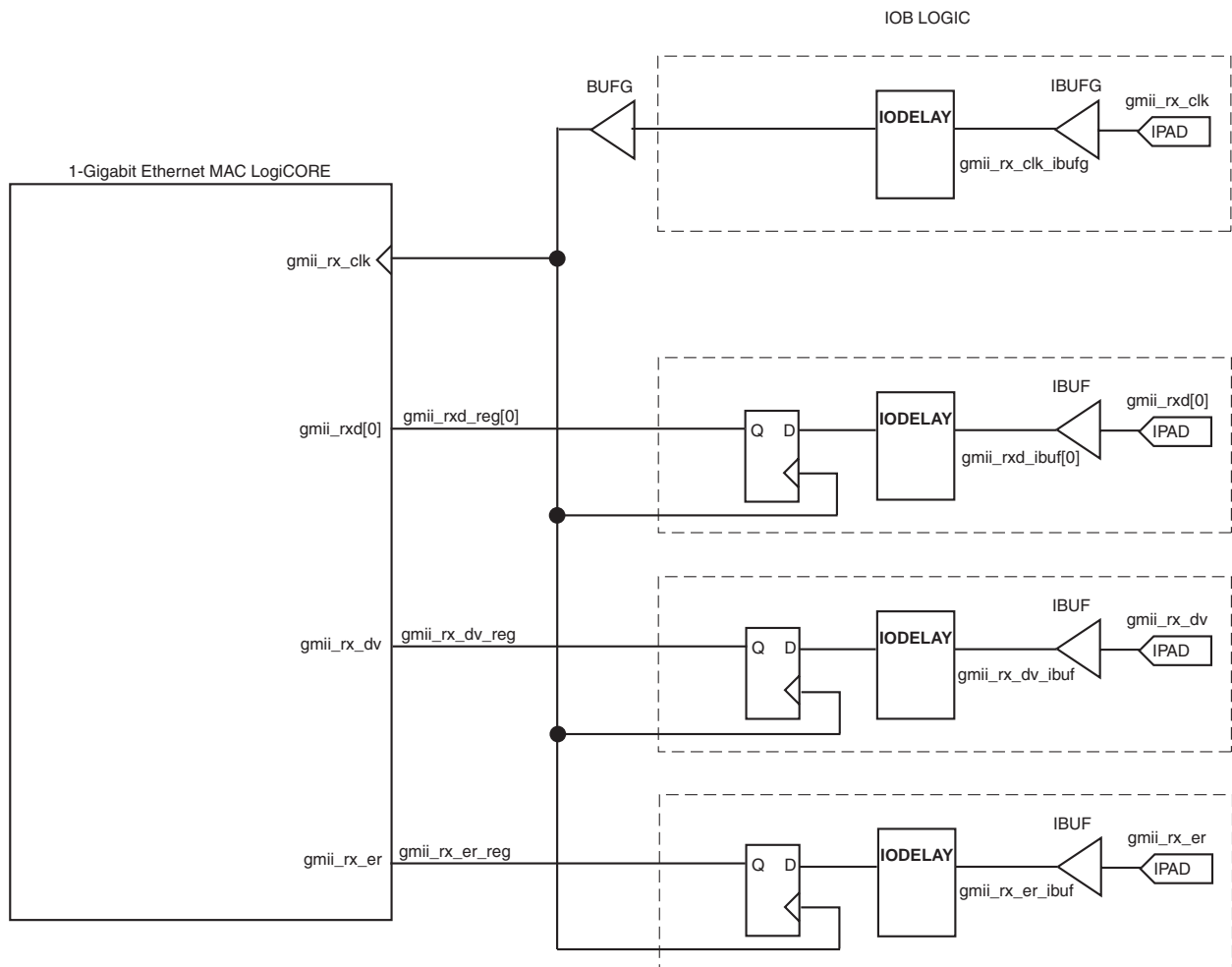


Figure 7-5: External GMII Receiver Logic for Virtex-5 Devices

## Implementing External RGMII

The HDL example design delivered with the core implements an external RGMII when RGMII is selected from the CORE Generator GUI (see Chapter 3, “Generating the Core”). For more information about using the example design, see the *1-Gigabit Ethernet MAC Getting Started Guide*.

## RGMII Transmitter Logic

### Virtex-II Pro, Virtex-II, Spartan-3, and Spartan-3A Devices

Figure 7-6 illustrates how to use the physical transmitter interface of the core to create an external RGMII in a Virtex-II device. The signal names and logic shown in this figure precisely match those delivered with the example design when the RGMII is selected. If other families are used, equivalent primitives and logic specific to that family is used in the example design.

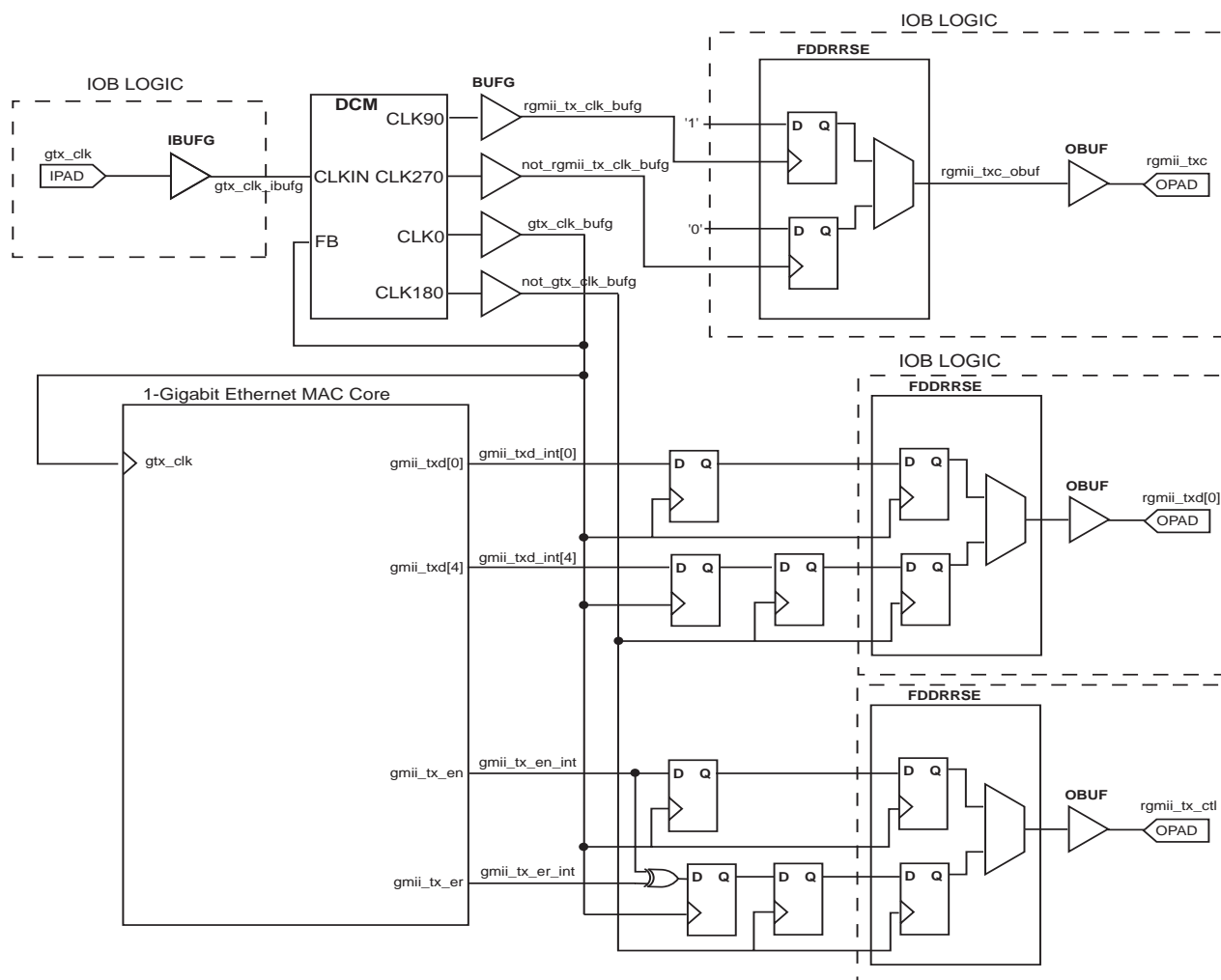


Figure 7-6: External RGMII Transmitter Logic

Figure 7-6 shows that the output transmitter signals are registered on `gtx_clk_bufg`, in the FPGA fabric, including the encoded `rgmii_tx_ctl_int` signal, derived from the logical xor of `gmii_tx_en_int` and `gmii_tx_er_int`. The signals to be transmitted on the RGMII falling clock edge are then registered on the falling edge clock, `not_gtx_clk_bufg`. This ensures that the data is presented to the Double Data Rate registers at the correct time. Finally, the transmitter signals are registered by an IOB output Double-Data-Rate (DDR) register before being driven to output pads.

The logic required to forward the transmitter clock is also shown. This uses an IOB output DDR register so the clock signal produced incurs on exactly the same delay as the data and

control signals. The `rgmii_tx_clk` clock signal is phase-shifted by 90 degrees in the DCM with respect to `gtx_clk_bufg`. This means that the rising edge of `rgmii_txc` occurs in the center of the data valid window—which maximizes setup and hold times across the interface, as specified in the *Reduced Gigabit Media Independent Interface (RGMII) Version 2.0* specification.

## Virtex-4 Devices

Figure 7-7 shows using the physical transmitter interface of the core to create an external RGMII in a Virtex-4 device. The signal names and logic shown exactly match those delivered with the example design when the RGMII is selected

Figure 7-7 also shows that the output transmitter signals are registered in the IOBs in ODDR components. These components convert the input signals into one double-data-rate signal. These signals are then output through OBUFs before being driven to output pads.

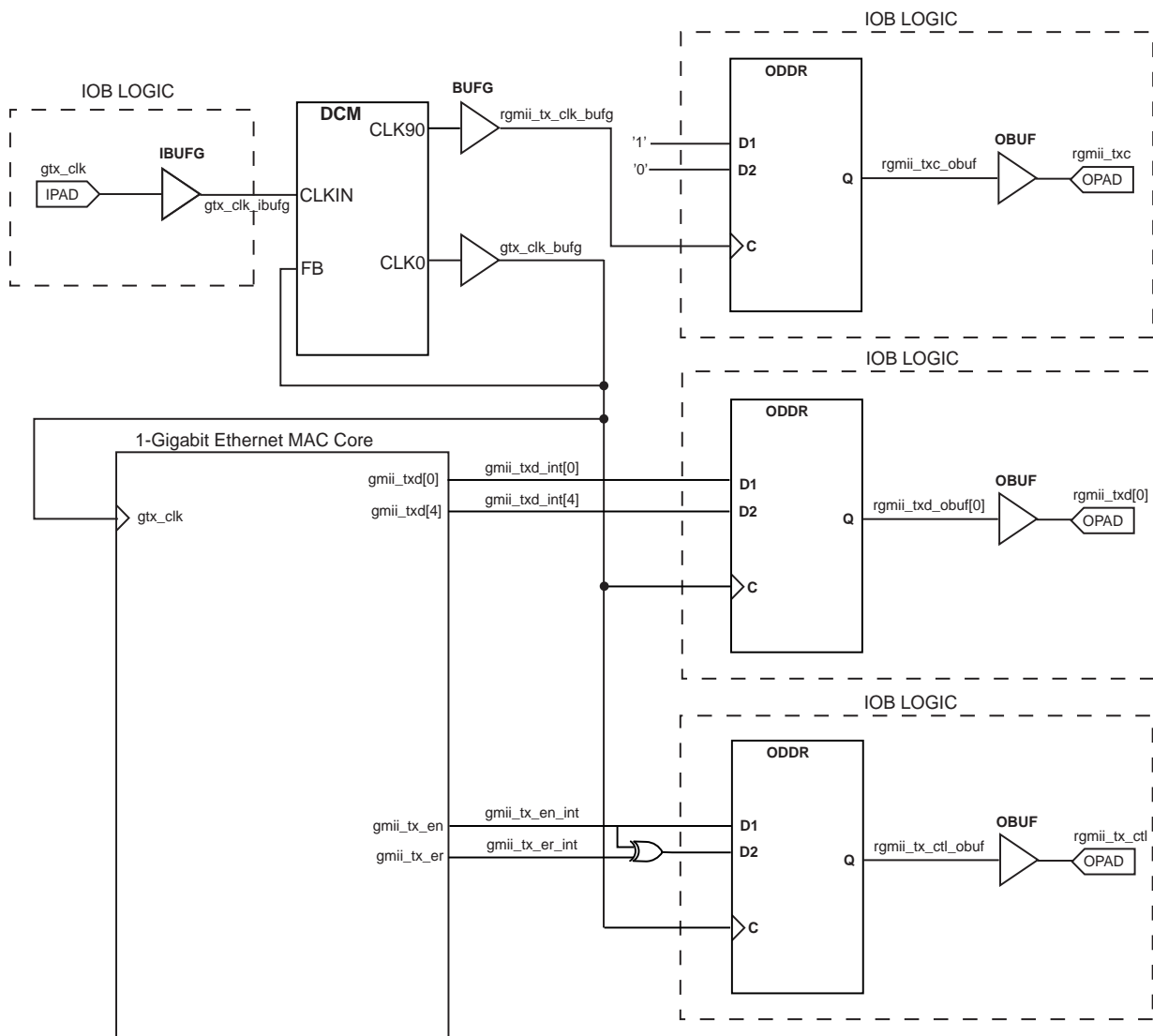


Figure 7-7: External RGMII Transmitter Logic in Virtex-4 Devices

The logic required to forward the transmitter clock is also shown: this uses an ODDR register so that the clock signal produced incurs exactly on the same delay as the data and control signals. The `rgmii_tx_clk` clock signal is phase-shifted by 90 degrees in the DCM with respect to `gtx_clk_bufg`. This means that the rising edge of `rgmii_txc` occurs in the center of the data valid window—which maximizes setup and hold times across the interface, as specified in the RGMII v2.0 specification.

## Virtex-5 Devices

The same logic that is used in [Figure 7-7](#) can also be used without modification for Virtex-5 devices. However, an alternative solution has been adopted for the example design delivered with the core. [Figure 7-8](#) shows using the physical transmitter interface of the core to create an external RGMII in a Virtex-5 device. The signal names and logic shown exactly match those delivered with the example design when the RGMII is selected.

[Figure 7-8](#) also shows that the output transmitter signals are registered in the IOBs in ODDR components. These components convert the input signals into one double-data-rate signal. The ODDR outputs are passed through IODELAYS—and these can be used to adjust the relationship between the individual signals. These signals are then output through OBUFs before being driven to output pads.

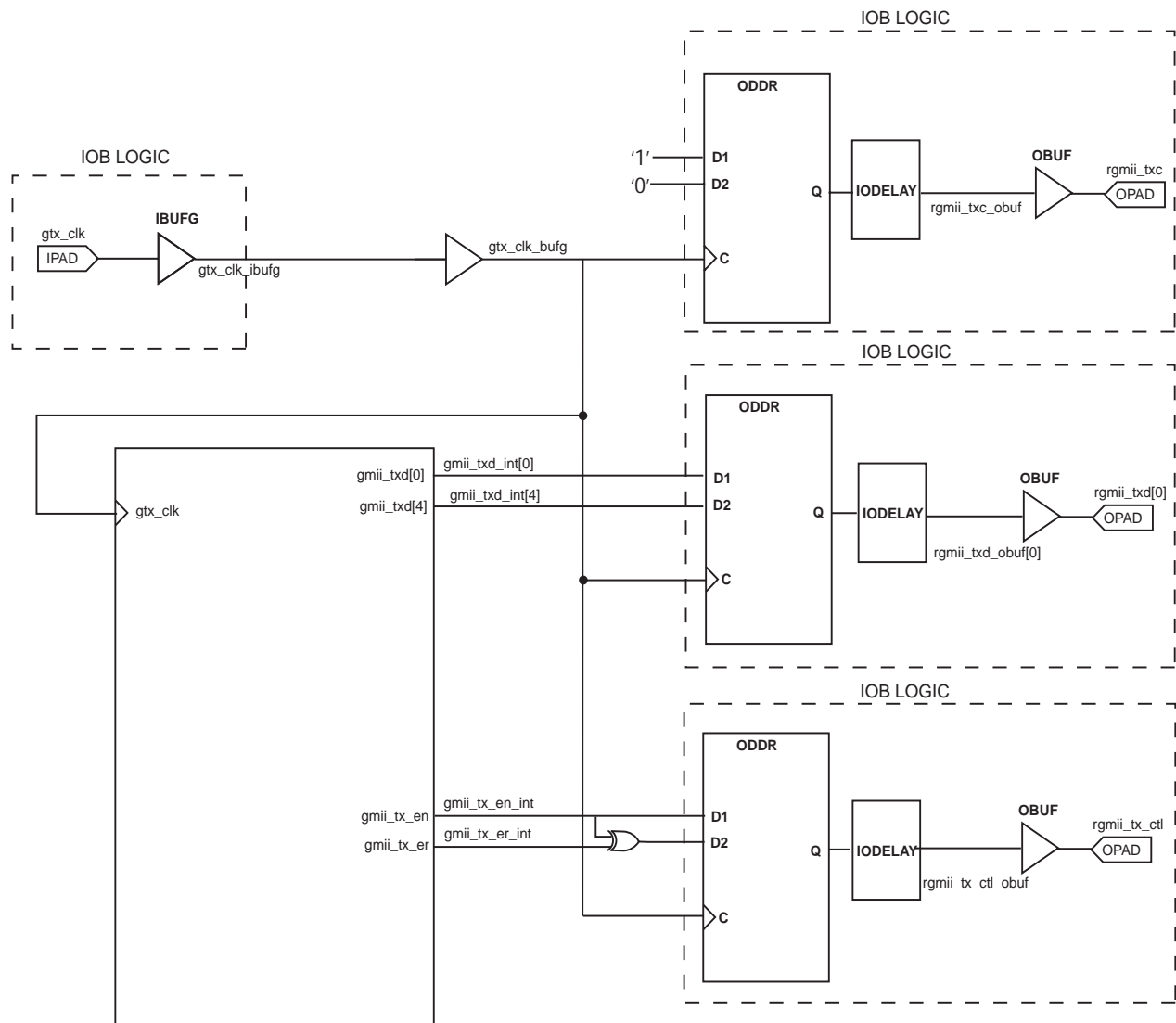


Figure 7-8: External RGMII Transmitter Logic in Virtex-5 Devices

The logic required to forward the transmitter clock is also shown. It has matching logic to the data and control signals to provide a known relationship between the signals. An IODELAY component is used to phase-shift the `rgmii_txc` clock signal by 90 degrees with respect to `gtx_clk_bufg`. This allows the rising edge of `rgmii_txc` to occur in the center of the data valid window—which maximizes setup and hold times across the interface, as specified in the RGMII v2.0 specification. The IODELAY component is used in fixed delay mode, where the attribute `ODELAY_VALUE` determines the tap delay value. An IDELAYCTRL primitive must be instantiated for this mode of operation. See the *Virtex-5 User Guide* for more information on the use of IDELAYCTRL and IODELAY components.



## RGMII Receiver Logic

### Virtex-II Pro, Virtex-II, Spartan-3, and Spartan-3A Devices

Figure 7-9 shows using the physical receiver interface of the core to create an external RGMII in a Virtex-II device. The signal names and logic exactly match those delivered with the example design when the RGMII is selected. If other families are used, equivalent primitives and logic specific to that family is used in the example design.

Figure 7-9 also shows that the input receiver signals are registered in device IOBs on rising edges of both `gmii_rx_clk_bufg` and `not_gmii_rx_clk_bufg`. The signals are then registered inside the FPGA fabric, before a final register stage to synchronize signals to the rising edge clock. To achieve the required setup and hold times across the interface, the DCM uses a phase-shift to adjust the clock relative to the data. See [Appendix C](#), “Calculating DCM Phase-Shifting.”

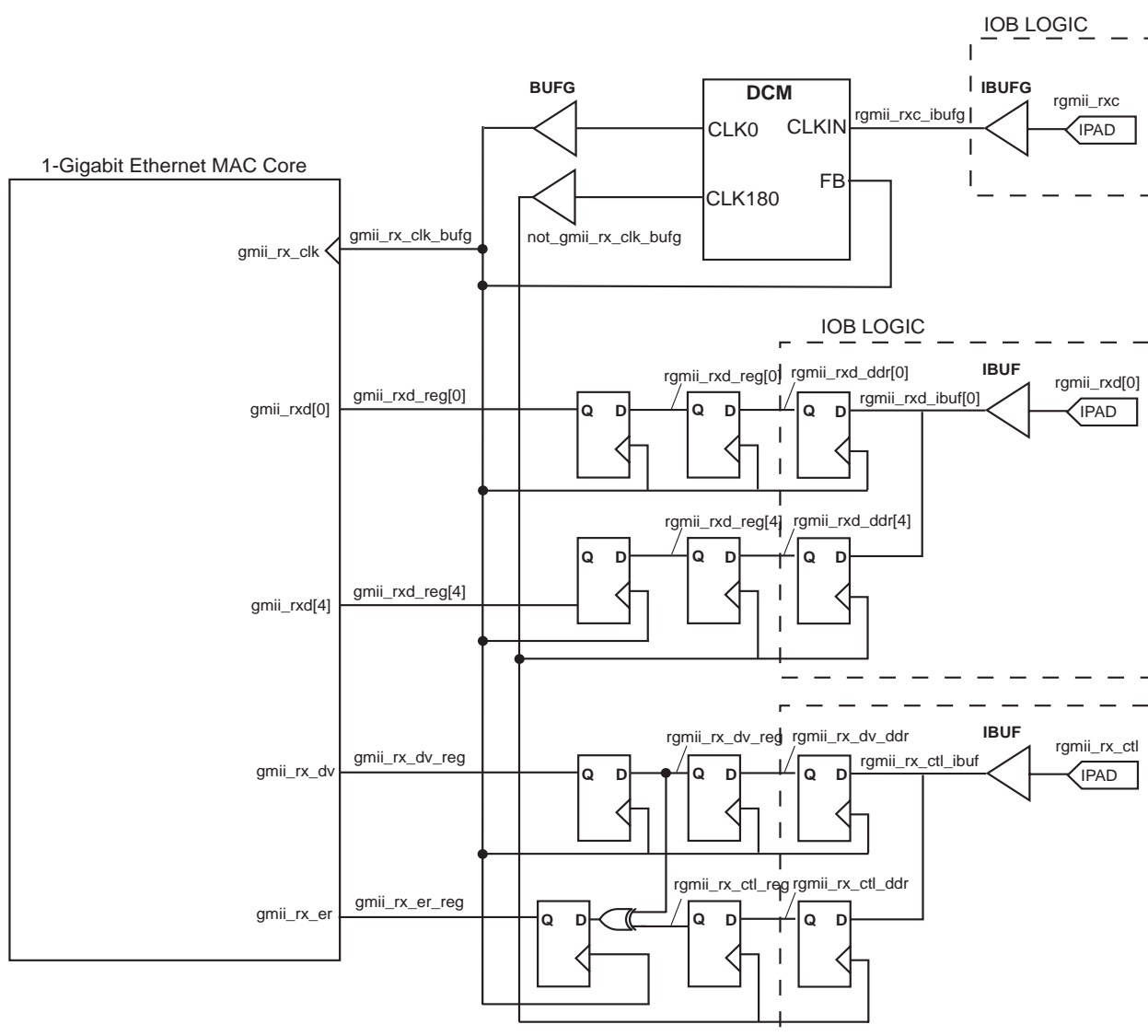


Figure 7-9: External RGMII Receiver Logic

## Virtex-4 Devices

Figure 7-10 shows using the physical receiver interface of the core to create an external RGMII in a Virtex-4 device. The signal names and logic exactly match those delivered with the example design when RGMII is selected.

Figure 7-10 also shows that the input receiver signals are registered in the IOBs in IDDR components. These components convert the input double data rate signals into GMII specification signals. The `gmii_rx_er_int` signal is derived in the FPGA fabric from the outputs of the control IDDR component.

The IDELAY components can be used to phase-shift the input RGMII clock, data, and control signals to meet the setup and hold margins and counter any bus skew. The IDELAY components are used in fixed delay mode, where the attribute `IOBDELAY_VALUE` determines the tap delay value. An IDELAYCTRL primitive must be instantiated for this mode of operation. See the *Virtex-4 User Guide* for more information about the IDELAYCTRL and IDELAY components.

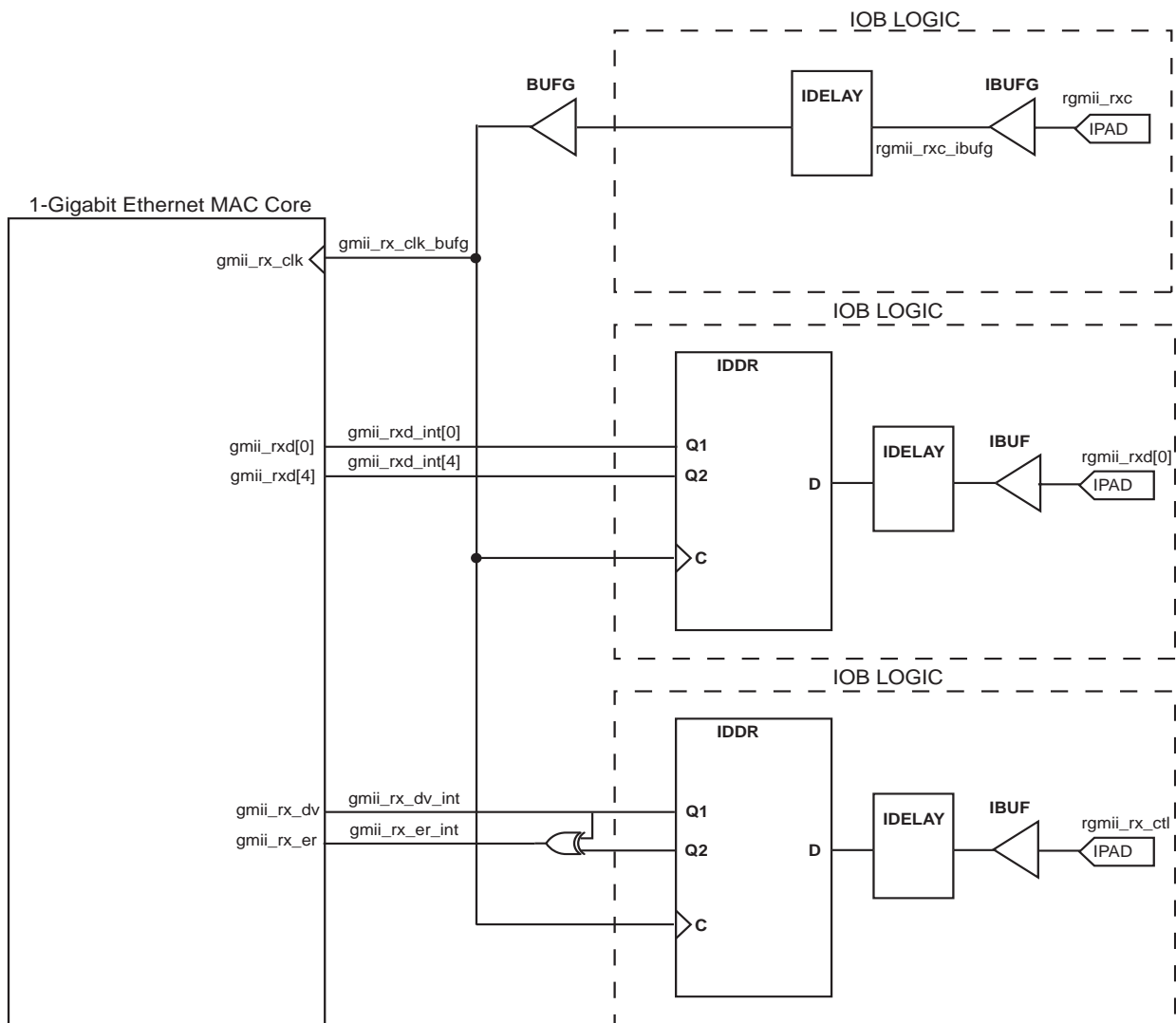


Figure 7-10: External RGMII Receiver Logic for Virtex-4 Devices

## Virtex-5 Devices

Figure 7-11 shows using the physical receiver interface of the core to create an external RGMII in a Virtex-5 device. The signal names and logic exactly match those delivered with the example design when RGMII is selected.

Figure 7-11 also shows that the input receiver signals are registered in the IOBs in IDDR components. These components convert the input double data rate signals into GMII specification signals. The `gmii_rx_er_int` signal is derived in the FPGA fabric from the outputs of the control IDDR component.

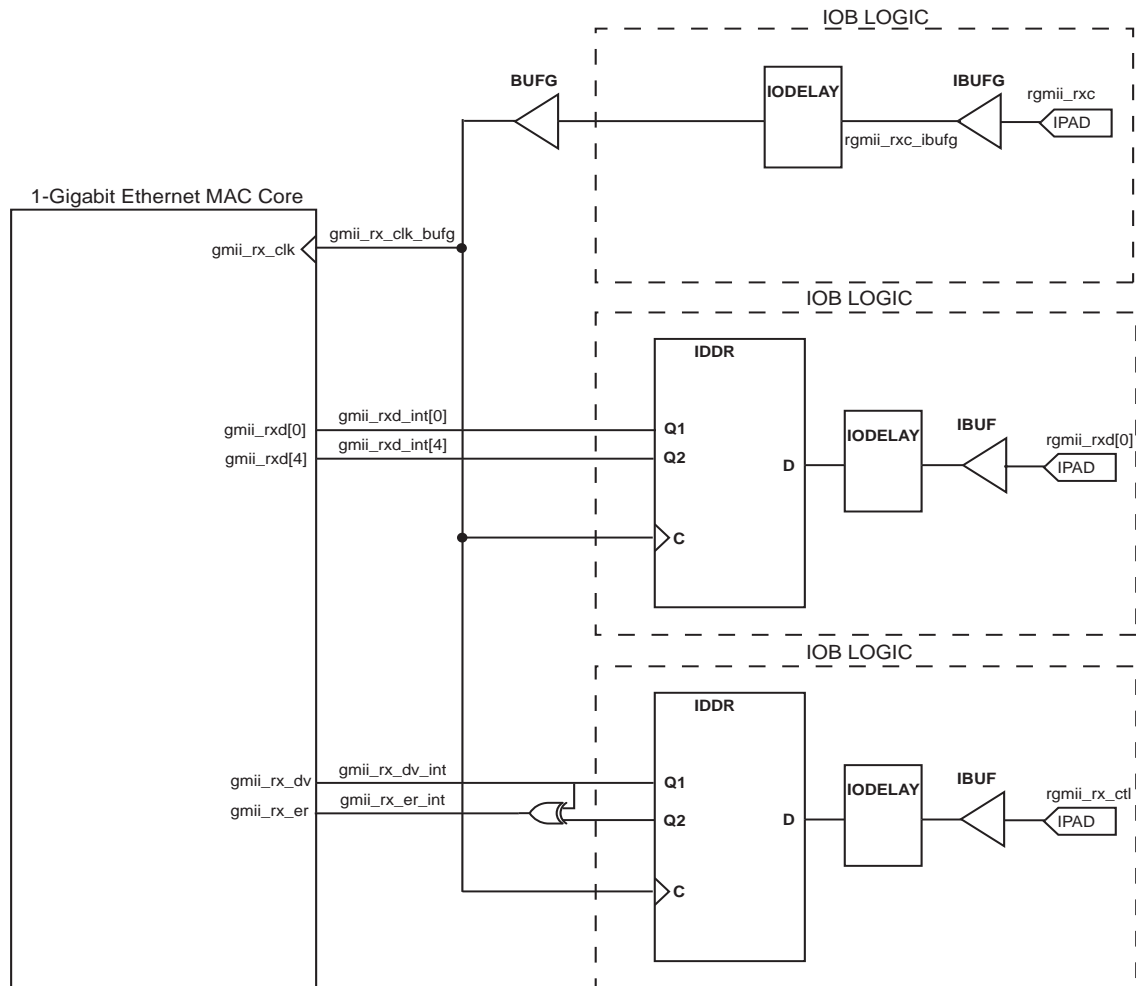


Figure 7-11: External RGMII Receiver Logic for Virtex-5 Devices

The IODELAY components are used to phase-shift the input RGMII clock, data and control signals to meet the setup and hold margins. The IODELAY components are used in fixed delay mode, where the attribute `IDELAY_VALUE` determines the tap delay value. An IDELAYCTRL primitive must be instantiated for this mode of operation. See the *Virtex-5 User Guide* for more information on the use of IDELAYCTRL and IODELAY components.

## RGMII Inband Status Decoding Logic

The inband status decoding logic is common to all device families. Figure 7-12 illustrates the decoding of RGMII inband status information. This information is received through the RGMII interface between frames in a Virtex-II device. The signal names and logic shown exactly match those delivered with the example design when the RGMII is selected. If other families are used, equivalent primitives and logic specific to that family is used in the example design.

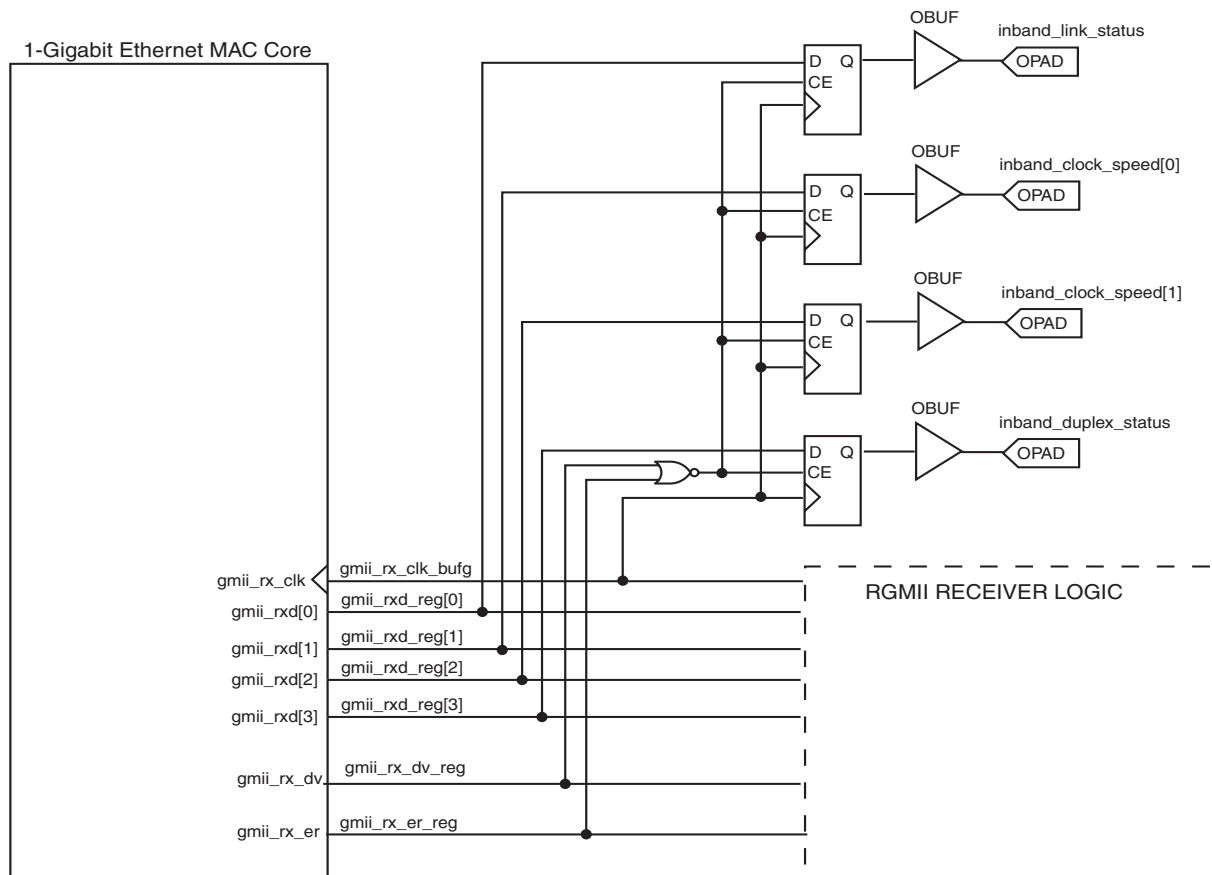


Figure 7-12: RGMII Inband Status Decoding Logic

## Using the MDIO interface

The MDIO interface is accessed through the optional management interface and is typically connected to the MDIO port of a physical-layer device to access its configuration and status registers (see “MDIO Interface,” on page 79). The MDIO format is defined in IEEE 802.3, clause 22.

## Connecting the MDIO to an Internally Integrated PHY

The MDIO ports of the GEMAC core can be connected to the MDIO ports of an internally integrated physical-layer device, such as the MDIO port of the Xilinx Ethernet 1000BASE-X PCS/PMA or SGMII core. See Chapter 11, “Interfacing to Other Cores” for more information.

## Connecting the MDIO to an External PHY

The MDIO ports of the GEMAC core can be connected to the MDIO of an external physical-layer device. In this situation, `mdio_in`, `mdio_out`, and `mdio_tri` must be connected to a tri-state buffer to create a bidirectional wire, `mdio`. This tri-state buffer can be either external to the FPGA, or internally integrated by using an IOB IOBUF component with an appropriate SelectIO™ standard for the external PHY (illustrated in Figure 7-13).

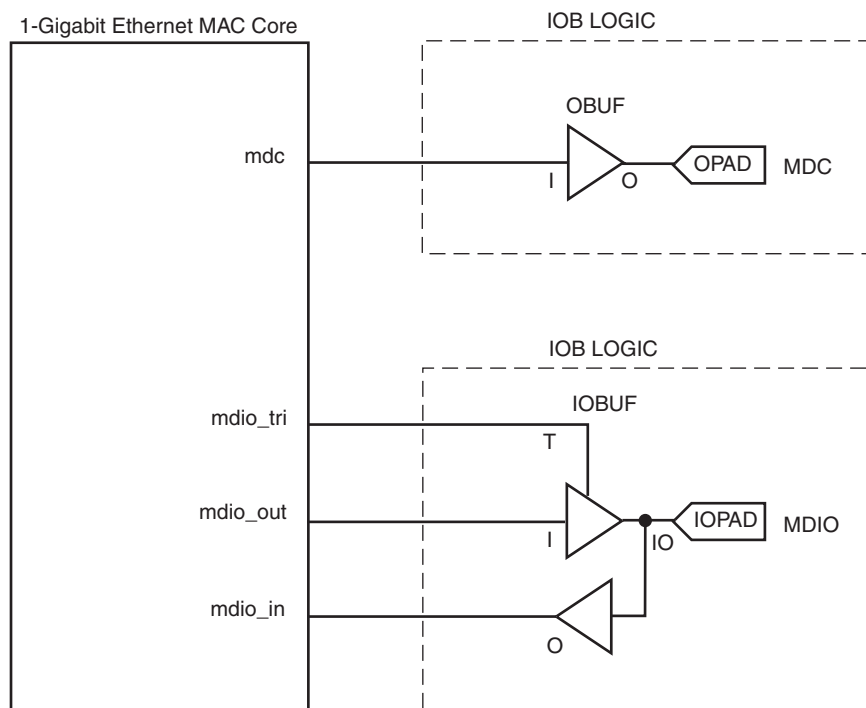


Figure 7-13: Creating an External MDIO Interface



## Configuration and Status

This chapter provides general guidelines for configuring and monitoring the GEMAC core, including a detailed description of the client-side management interface and registers present in the core. It also describes the alternative to the optional management interface which is the Configuration Vector.

### Using the Optional Management Interface

The Management Interface is a processor-independent interface with standard address, data, and control signals. It may be used as is, or a wrapper (not supplied) may be used to interface to common bus architectures. For port definition, see [“Management Interface \(Optional\),” on page 26](#).

This interface is used for:

- Configuring of the GEMAC core via the configuration registers.
- Access through the MDIO interface to the management registers located in the PHY connected to the GEMAC core.

The Management Interface can be accessed in different ways, depending on the type of transaction. A truth table showing which access method is required for each transaction type is shown in [Table 8-1](#). These access methods are described in the following sections.

**Table 8-1: Management Interface Transaction Types**

Transaction	host_miim_sel	host_addr[9]
Configuration	0	1
MIIM access	1	X

### Host Clock Frequency

The Management Interface clock, `host_clk`, is used to derive the MDIO clock, `mdc`, and for this reason is subject to the following frequency restriction:

$$\geq 10 \text{ MHz}$$

Configuring the GEMAC core to derive the `mdc` signal from this clock is detailed in [“MDIO Interface,” on page 79](#).

### Configuration Registers

After a power-up or system reset, the client may reconfigure the core parameters using their defaults. Configuration changes can be written at any time. Both the receiver and

transmitter logic responds only to configuration changes during inter-frame gaps. The exceptions to this are the configurable resets that take effect immediately.

Configuration of the GEMAC core is performed through a register bank that is accessed through the management interface. The configuration registers available in the core are listed in [Table 8-2](#). As shown, the address has some implicit *don't care* bits; any access to an address in the ranges shown performs a 32-bit read or write from the same configuration word.

**Table 8-2: Configuration Registers**

Address	Description
0x200-0x23F	Receiver Configuration (Word 0)
0x240-0x27F	Receiver Configuration (Word 1)
0x280-0x2BF	Transmitter Configuration
0x2C0-0x2FF	Flow Control Configuration
0x300-0x33F	Reserved
0x340-0x37F	Management Configuration
0x380-0x383	Unicast Address (Word 0) (if Address Filter is present)
0x384-0x387	Unicast Address (Word 1) (if Address Filter is present)
0x388-0x38B	Address Table Configuration (Word 0) (if Address Filter is present)
0x38C-0x38F	Address Table Configuration (Word 1) (if Address Filter is present)
0x390-0x393	Address Filter Mode (if Address Filter is present)

## Receiver Configuration

The register contents for the two receiver configuration words are shown in [Table 8-3](#) and [Table 8-4](#).

**Table 8-3: Receiver Configuration Word 0**

Bit	Default Value	Description
31-0	All 0s	<p>Pause frame MAC Source Address[31:0]. This address is used by the MAC to match against the destination address of any incoming flow control frames. It is also used by the flow control block as the source address (SA) for any outbound flow control frames. See <a href="#">Chapter 6, "Using Flow Control."</a></p> <p>The address is ordered so the first byte transmitted/received is the lowest positioned byte in the register; for example, a MAC address of AA-BB-CC-DD-EE-FF would be stored in Address[47:0] as 0xFFEEDDCCBBAA.</p>



Table 8-4: Receiver Configuration Word 1

Bit	Default Value	Description
15-0	All 0s	Pause frame MAC Source Address[47:32]
24-16	n/a	Reserved
25	0	<b>Length/Type Error Check Disable</b> When this bit is set to '1,' the core will not perform the length/type field error checks as described in <a href="#">“Length/Type Field Error Checks,” on page 40</a> . When this bit is set to '0,' the length/type field checks will be performed; this is normal operation.
26	n/a	Reserved
27	0	<b>VLAN Enable</b> When this bit is set to '1,' VLAN tagged frames will be accepted by the receiver.
28	1	<b>Receiver Enable.</b> If set to '1,' the receiver block will be operational. If set to '0,' the block will ignore activity on the physical interface RX port.
29	0	<b>In-band FCS Enable</b> When this bit is '1,' the MAC receiver will pass the FCS field up to the client. When at '0,' the client will not be passed the FCS. In both cases, the FCS will be verified on the frame.
30	0	<b>Jumbo Frame Enable</b> When this bit is set to '1,' the MAC receiver will accept frames over the specified <i>IEEE 802.3-2002</i> maximum legal length. When this bit is '0,' the MAC will only accept frames up to the specified maximum.
31	0	<b>Reset</b> When this bit is set to '1,' the receiver will be reset. The bit will then automatically revert to '0.' Note that this reset also sets all of the receiver configuration registers to their default values.

## Transmitter Configuration

The register contents for the Transmitter Configuration Word are described in [Table 8-5](#).

Table 8-5: Transmitter Configuration Word

Bit	Default Value	Description
24-0	n/a	Reserved
25	0	<b>Interframe Gap Adjust Enable</b> If '1,' the transmitter will read the value on the port tx_ifg_delay at the start of frame transmission and adjust the interframe gap following the frame accordingly.
26	n/a	Reserved
27	0	<b>VLAN Enable</b> When this bit is set to '1,' the transmitter will allow the transmission of VLAN tagged frames.

Table 8-5: Transmitter Configuration Word (Continued)

Bit	Default Value	Description
28	1	<b>Transmit Enable</b> When this bit is '1,' the transmitter is operational. When it is '0,' the transmitter is disabled.
29	0	<b>In-band FCS Enable</b> When this bit is '1,' the MAC transmitter will expect the FCS field to be passed in by the client. When this bit is '0,' the MAC transmitter will append padding as required, compute the FCS and append it to the frame.
30	0	<b>Jumbo Frame Enable</b> When this bit is set to '1,' the MAC transmitter will send frames that are greater than the specified IEEE 802.3-2002 maximum legal length. When this bit is '0,' the MAC will only send frames up to the specified maximum.
31	0	<b>Reset</b> When set to '1,' the transmitter will be reset. The bit will then automatically revert to '0.' Note that this reset will also set all of the transmitter configuration registers to their default values.

## Flow Control Configuration

Table 8-6 lists the register contents for the Flow Control Configuration Word.

Table 8-6: Flow Control Configuration Word

Bit	Default Value	Description
28-0	n/a	Reserved
29	1	<b>Flow Control Enable (RX)</b> When this bit is '1,' received flow control frames will inhibit the transmitter operation. When this bit is '0,' received flow control frames will always be passed up to the client.
30	1	<b>Flow Control Enable (TX)</b> When this bit is '1,' asserting the PAUSE_REQ signal will send a flow control frame out from the transmitter. When this bit is '0,' asserting the PAUSE_REQ signal has no effect.
31	n/a	Reserved

## MDIO Configuration

The register contents for the Management Configuration Word are described in [Table 8-7](#).

**Table 8-7: Management Configuration Word**

Bits	Default Value	Description
4-0	All 0s	<b>Clock Divide[4:0]</b> This value enters a logical equation which enables the mdc frequency to be set as a divided down ratio of the host_clk frequency.
5	0	<b>MDIO Enable</b> When this bit is '1,' the MDIO interface can be used to access attached PHY devices. When this bit is '0,' the MDIO interface is disabled and the MDIO signals remain inactive.
31-6	n/a	Reserved

## Address Filter Configuration

[Table 8-8](#) through [Table 8-12](#) describe the registers used to access the Address Filter configuration when the GEMAC core implemented with an Address Filter. The register contents for the two unicast address registers are found in [Table 8-8](#) and [Table 8-9](#).

**Table 8-8: Unicast Address Word 0**

Bits	Default Value	Description
31-0	All 0s	Address filter unicast address[31:0]. The address is ordered so the first byte received is the lowest positioned byte in the register; for example, a MAC address of AA-BB-CC-DD-EE-FF would be stored in Address[47:0] as 0xFFEEDDCCBBAA.

**Table 8-9: Unicast Address Word 1**

Bits	Default Value	Description
15-0	All 0s	Address filter unicast address[47:32].
31-16	N/A	Reserved

The Address Filter can be programmed to respond to four separate additional addresses stored in an address table in the Address Filter. [Table 8-10](#) and [Table 8-11](#) describe how the contents of the address table are set.

**Table 8-10: Address Table Configuration Word 0**

Bits	Default Value	Description
31–0	All 0s	MAC Address[31:0]. The address is ordered so the first byte received is the lowest positioned byte in the register; for example, a MAC address of AA-BB-CC-DD-EE-FF would be stored in Address[47:0] as 0xFFEEDDCCBBA.

**Table 8-11: Address Table Configuration Word 1**

Bits	Default Value	Description
15–0	All 0s	MAC Address[47:32].
17–16	All 0s	The location in the address table that the MAC address is to be written to or read from. There are up to 4 entries in the table (Location 0 to 3).
22–18	N/A	Reserved
23	0	<b>Read not write</b> This bit is set to '1' to read from the address table. If it is set to '1,' the contents of the table entry that is being accessed by the bits 17-16 will be output on the hostrddata bus in consecutive cycles (least significant word first). If it is set to '0,' the data on bits 15-0 is written into the table at the address specified by bits 17-16.
31–24	N/A	Reserved

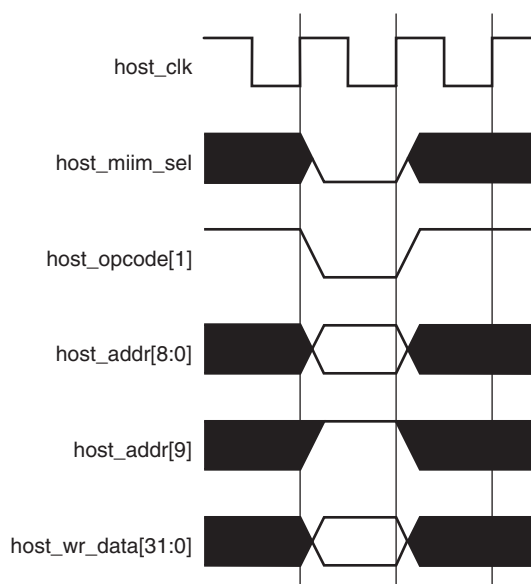
The contents of the Address Filter mode register are described in [Table 8-12](#).

**Table 8-12: Address Filter Mode**

Bits	Default Value	Description
30–0	N/A	Reserved
31	0	<b>Promiscuous Mode</b> If this bit is set to '1,' the Address Filter operates in promiscuous mode. All frames are passed to the receiver client, regardless of the destination address.

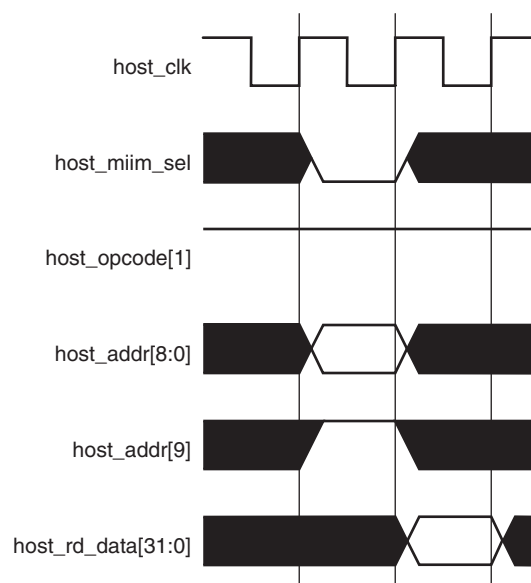
## Writing and Reading to and from the Configuration Registers

Writing to the configuration registers through the management interface is shown in [Figure 8-1](#). When accessing the configuration registers (when `host_addr[9] = '1'` and `host_miim_sel = '0'`), the upper bit of `host_opcode` functions as an active low write enable signal. The lower `host_opcode` bit is a *don't care* bit.



**Figure 8-1: Configuration Register Write Timing**

Reading from the configuration register words is similar, but the upper `host_opcode` bit should be '1,' as shown in Figure 8-2. In this case, the contents of the register appear on `host_rd_data` the `host_clk` edge after the register address is asserted onto `host_addr`.



**Figure 8-2: Configuration Register Read Timing**

## Accessing the Address Table

To write to a specific entry in the address table, you must first write the least significant 32-bits of the address into the Address Table Configuration (Word 0) register. You then write the most significant 16 bits together with the location in the table (bits 17–16) to the

Address Table Configuration (Word1) register with bit 23 (read not write) set to '0.' This is shown in Figure 8-3. Although it is shown in the figure, there is no requirement for the two writes to be on adjacent cycles.

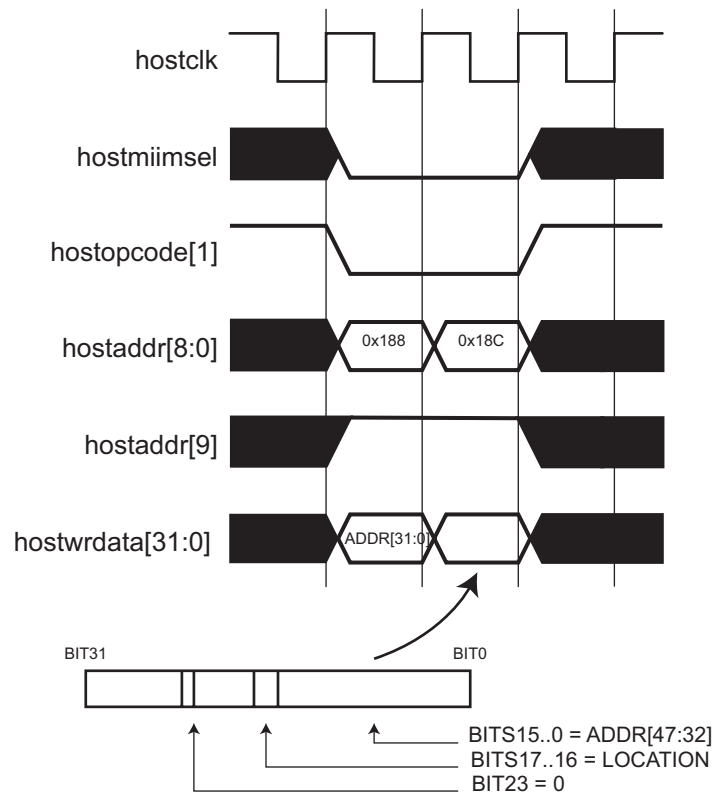


Figure 8-3: Address Table Write Timing

As shown in Figure 8-4, you must write to the Address Table Configuration register (Word 1) with the location set to the desired table entry and bit 23 set to '1' to read from the address table. On the next cycle the least significant word appears on the hostwrrdata bus. One cycle afterwards, the most significant 16-bits are output on the lower 16 bits of the bus.

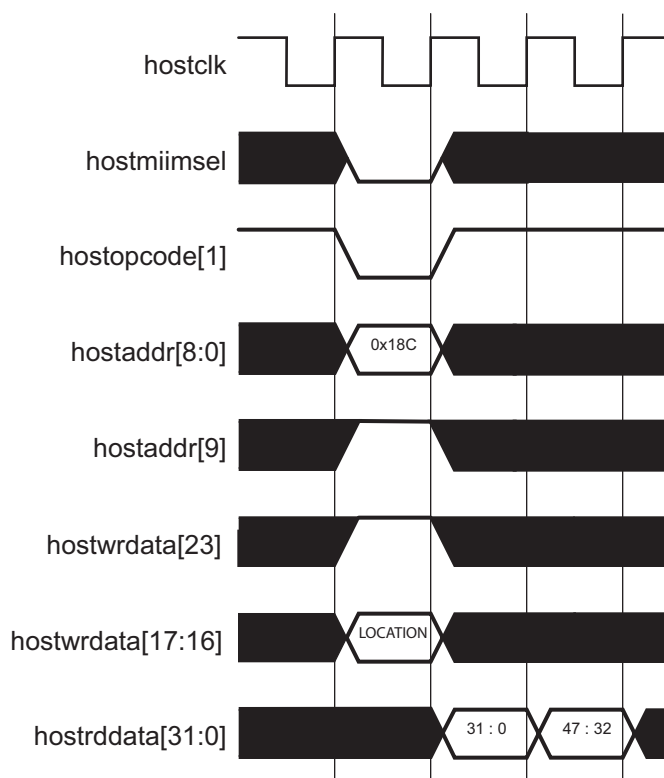


Figure 8-4: Address Table Read Timing

## MDIO Interface

### Introduction to MDIO

The MDIO interface for 1 Gbps operation (and slower speeds) is defined in *IEEE 802.3* clause 22. This is a two wire interface consisting of a clock, *mdc*, and a shared serial data line, *mdio*.

An MDIO bus in a system consists of a single Station Management (STA) master management entity and a number of MDIO Managed Device (MMD) slave entities. [Figure 8-5](#) illustrates a typical system. All transactions are initiated by the STA entity. The GEMAC core implements a STA and can be connected to MMDs (PHY devices) to access their management registers.

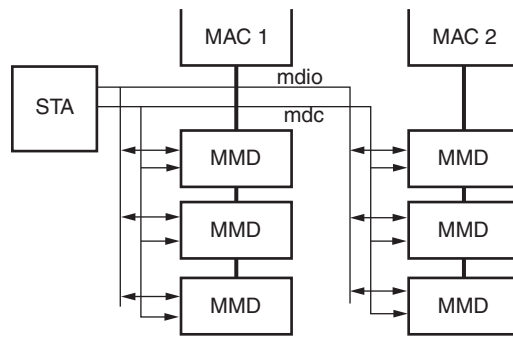


Figure 8-5: Typical MDIO-managed System

There are two different transaction types of MDIO for write and read. They are described in this section.

### Abbreviations Used

The following abbreviations apply for the remainder of this chapter.

- **PRE** Preamble
- **ST** Start of frame
- **OP** Operation code
- **PHYAD** PHY address
- **REGAD** Register address
- **TA** Turnaround

### Write Transaction

Figure 8-6 shows a write transaction across the MDIO, as defined by OP='01.' The addressed MMD (PHYAD) device takes the 16-bit word in the data field and writes it to the register at REGAD.

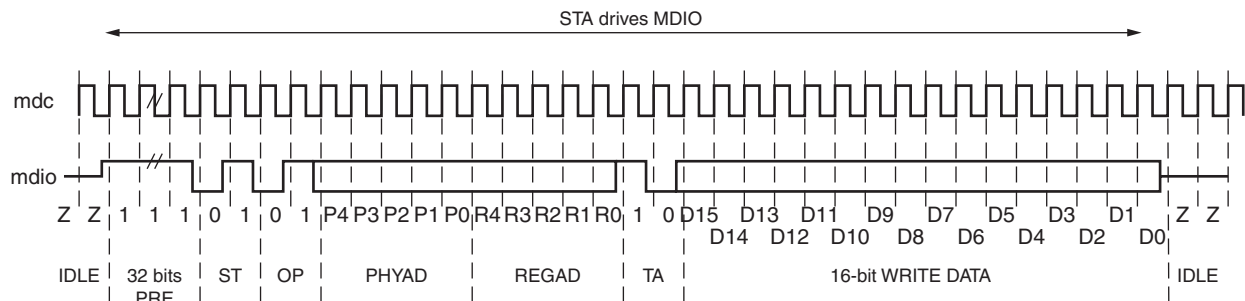


Figure 8-6: MDIO Write Transaction

### Read Transaction

Figure 8-7 shows a Read transaction; this is defined by OP="10". The addressed MMD (PHYAD) device returns the 16-bit word from the register at REGAD.



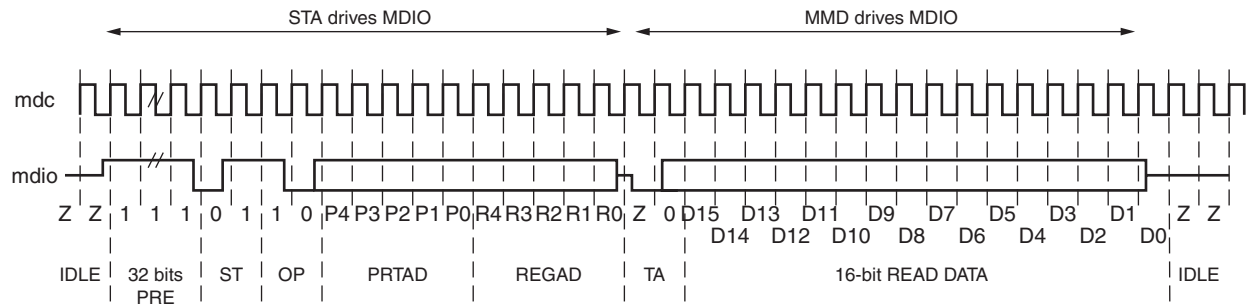


Figure 8-7: MDIO Read Transaction

For details of the register map of MMD (PHY layer devices) and a detailed description of the operation of the MDIO Interface itself, see *IEEE 802.3-2002*.

## Accessing MDIO With GEMAC

More information about MDIO with GEMAC can be found in the following sections of this guide:

- For the GEMAC port definition of the MDIO, see [“MDIO Interface” in Chapter 2](#)
- [“Connecting the MDIO to an Internally Integrated PHY,” on page 68](#)
- [“Connecting the MDIO to an External PHY,” on page 69](#)

The management interface is also used to access the MDIO interface of the GEMAC core. The MDIO interface supplies a clock to the connected PHY, `mdc`. This clock is derived from the `host_clk` signal using the value in the `Clock Divide[4:0]` configuration register. The frequency of `mdc` is given by the following equation:

$$f_{MDC} = \frac{f_{HOST\_CLK}}{(1 + \text{Clock Divide}[4:0]) \times 2}$$

The frequency of `mdc` given by this equation should not exceed 2.5 MHz in order to comply with the *IEEE 802.3-2002* specification for this interface. To prevent `mdc` from being out of specification, the `Clock Divide[4:0]` value powers up at 00000. While this value is in the register, it is impossible to enable the MDIO interface.

For details of the register map of PHY layer devices and a detailed description of the operation of the MDIO interface itself, see *IEEE 802.3-2002*.

Figure 8-8 shows access to the MDIO interface through the Management Interface.

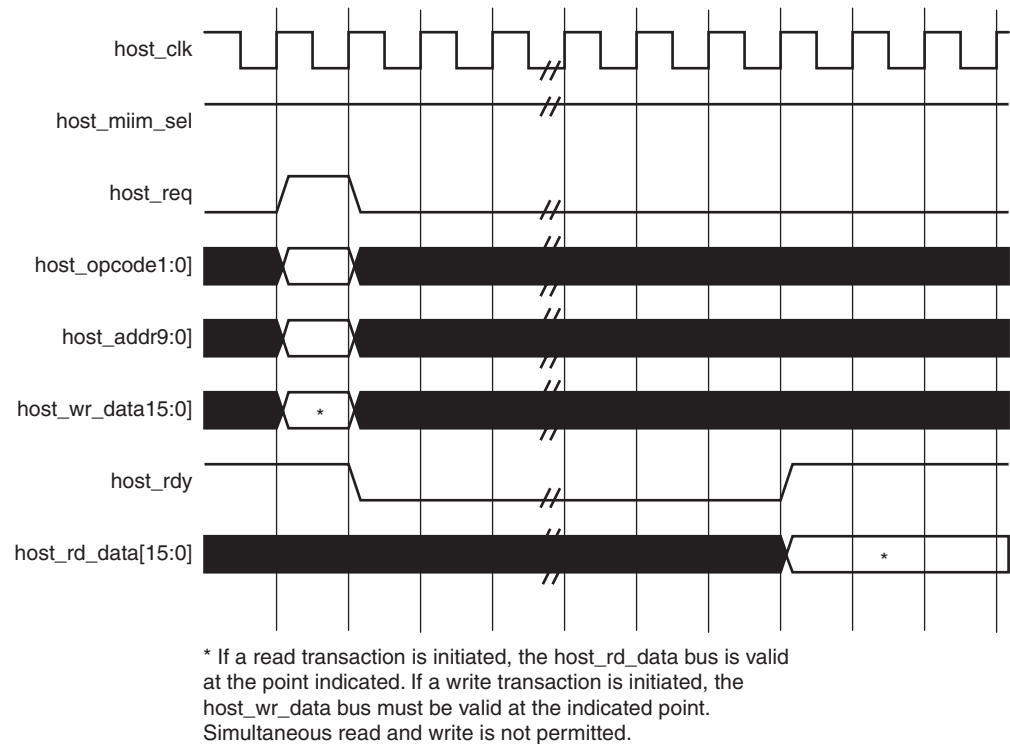


Figure 8-8: MDIO Access through Management Interface

For MDIO transactions, the following applies:

- host\_miim\_sel is '1'
- host\_opcode[1:0] maps to the OP (opcode) field of the MDIO frame
- host\_addr maps to the two address fields of the MDIO frame; PHYAD is host\_addr[9:5], and REGAD is host\_addr[4:0]
- host\_wr\_data[15:0] maps into the data field of the MDIO frame when performing a write operation
- The data field of the MDIO frame maps into host\_rd\_data[15:0] when performing a read operation

The GEMAC core signals the host that it is ready for an MDIO transaction by asserting host\_miim\_rdy. A read or write transaction on the MDIO is initiated by a pulse on the host\_req signal. This pulse is ignored if the MDIO interface already has a transaction in progress. The GEMAC core then deasserts the host\_miim\_rdy signal while the transaction across the MDIO is in progress. When the transaction across the MDIO interface has been completed, the host\_miim\_rdy signal is asserted by the GEMAC core; if the transaction is a read, the data is available on the host\_rd\_data[15:0] bus at this time.

## Access without the Management Interface

If the optional management interface is omitted from the core, all of the relevant configuration settings described in [Table 8-3](#) through [Table 8-6](#) are brought out of the core as signals. These signals are bundled into the `configuration_vector[64:0]` signal as described in [Table 8-13](#).

These signals may permanently set by connecting to logic 0 or 1, or may be changed by the user application at any time; however, with the exception of the reset and the flow control configuration signals, any changes do not take effect until the current frame has completed transmission or reception.

The Clock heading in [Table 8-13](#) denotes which clock domain the configuration signal is registered into before use by the core. It is not necessary to drive the signal from this clock domain.

**Table 8-13: Configuration Vector Bit Definition**

Bit(s)	Configuration Register cross reference	Clock	Description
47:0	"Receiver Configuration Word 0" bits 31-0 and "Receiver Configuration Word 1" bits 15-0	gmii_rx_clk	<p><b>Pause frame MAC Source Address[47:0]</b> This address is used by the GEMAC core to match against the destination address of any incoming flow control frames, and as the source address for any outbound flow control frames.</p> <p>The address is ordered such that the first byte transmitted or received is the least significant byte in the register; for example, a MAC address of AA-BB-CC-DD-EE-FF will be stored in bite [47:0] as 0xFFEEDDCCBBAA.</p>
48	n/a	n/a	This input is unused.
49	"Receiver Configuration Word 1" bit 27	gmii_rx_clk	<b>Receiver VLAN Enable</b> When this bit is set to '1,' VLAN tagged frames are accepted by the receiver.
50	"Receiver Configuration Word 1" bit 28	gmii_rx_clk	<b>Receiver Enable</b> If set to '1,' the receiver block is operational. If set to '0,' the block ignores activity on the physical interface RX port.
51	"Receiver Configuration Word 1" bit 29	gmii_rx_clk	<b>Receiver In-band FCS Enable</b> When this bit is '1,' the MAC receiver will pass the FCS field up to the client. When it is '0,' the MAC receiver will not pass the FCS field. In both cases, the FCS field will be verified on the frame.
52	"Receiver Configuration Word 1" bit 30	gmii_rx_clk	<b>Receiver Jumbo Frame Enable</b> When this bit is '0,' the receiver will not pass frames longer than the maximum legal frame size specified in <i>IEEE 802.3-2002</i> . At '1,' the receiver will not have an upper limit on frame size.

Table 8-13: Configuration Vector Bit Definition (*Continued*)

Bit(s)	Configuration Register cross reference	Clock	Description
53	<a href="#">“Receiver Configuration Word 1” bit 31</a>	n/a	<b>Receiver Reset.</b> When this bit is ‘1,’ the receiver is held in reset. This signal is an input to the reset circuit for the receiver block.
54	<a href="#">“Transmitter Configuration Word” bit 25</a>	gtx_clk	<b>Transmitter Interframe Gap Adjust Enable</b> If ‘1,’ then the transmitter will read the value of the tx_ifg_delay port and set the interframe gap accordingly. If ‘0,’ the transmitter will always insert at least the legal minimum interframe gap.
55	n/a	n/a	Not used.
56	<a href="#">“Transmitter Configuration Word” bit 27</a>	gtx_clk	<b>Transmitter VLAN Enable</b> When this bit is set to ‘1,’ the transmitter allows the transmission of VLAN tagged frames.
57	<a href="#">“Transmitter Configuration Word” bit 28</a>	gtx_clk	<b>Transmitter Enable</b> When this bit is ‘1,’ the transmitter will be operational. When it is ‘0,’ the transmitter is disabled.
58	<a href="#">“Transmitter Configuration Word” bit 29</a>	gtx_clk	<b>Transmitter In-Band FCS Enable</b> When this bit is ‘1,’ the MAC transmitter will expect the FCS field to be pass in by the client. When it is ‘0,’ the MAC transmitter will append padding as required, compute the FCS and append it to the frame.
59	<a href="#">“Transmitter Configuration Word” bit 30</a>	gtx_clk	<b>Transmitter Jumbo Frame Enable</b> When this bit is ‘1,’ the MAC transmitter will allow frames larger than the maximum legal frame length specified in <i>IEEE 802.3-2002</i> to be sent. When set to ‘0,’ the MAC transmitter will only allow frames up to the legal maximum to be sent.
60	<a href="#">“Transmitter Configuration Word” bit 31</a>	n/a	<b>Transmitter Reset.</b> When this bit is ‘1,’ the MAC transmitter is held in reset. This signal is an input to the reset circuit for the transmitter block.
61	<a href="#">“Flow Control Configuration Word” bit 29</a>	gtx_clk	<b>Transmit Flow Control Enable.</b> When this bit is ‘1,’ asserting the pause_req signal causes the GEMAC core to send a flow control frame out from the transmitter. When this bit is ‘0,’ asserting the pause_req signal will have no effect.

Table 8-13: Configuration Vector Bit Definition (*Continued*)

Bit(s)	Configuration Register cross reference	Clock	Description
62	<a href="#">“Flow Control Configuration Word”</a> bit 30	gtx_clk	<b>Receive Flow Control Enable.</b> When this bit is ‘1,’ received flow control frames will inhibit the transmitter operation. When at ‘0,’ received flow frames are passed up to the client.
63	<a href="#">“Receiver Configuration Word 1”</a> bit 25	gtx_clk	<b>Length/Type Error Check Disable.</b> When this bit is set to ‘1,’ the core will not perform the length/type field error checks as described in <a href="#">“Length/Type Field Error Checks,”</a> on page 40. When this bit is set to ‘0,’ the length/type field checks will be performed: this is normal operation.
64	<a href="#">“Address Filter Mode”</a> bit 31	gmii_rx_clk	<b>Address Filter Enable.</b> When this bit is ‘0,’ the Address Filter is enabled. If it is set to ‘1,’ the Address Filter will operate in promiscuous mode.



# Constraining the Core

---

This chapter defines the GEMAC core constraint requirements. An example UCF that implements the constraints defined in this chapter is provided with the HDL example design for the core.

See the *1-Gigabit Ethernet MAC Getting Started Guide* for more information about the CORE Generator output files and detailed information about the HDL example design.

## Required Constraints

### Device, Package, and Speedgrade Selection

The GEMAC can be implemented in Virtex-II, Virtex-II Pro, Virtex-4, Virtex-5, Spartan-3, Spartan-3E, and Spartan-3A devices with the following attributes:

- Large enough to accommodate the core
- Contains a sufficient number of IOBs
- Operates at the following speed grades:
  - ◆ –4 for Virtex-II, Spartan-3, Spartan-3E, and Spartan-3A
  - ◆ –5 for Virtex-II Pro
  - ◆ –10 for Virtex-4
  - ◆ –1 for Virtex-5

### I/O Location Constraints

No specific I/O location constraints are required.

### Placement Constraints

No specific placement constraints are required.

### Timing Constraints

The core can have up to three separate clock domains: `gtx_clk` for the transmitter logic, `gmii_rx_clk` for the receiver logic, and `host_clk` for the optional management logic. These clock nets and the signals within the core that cross these clock domains must be constrained appropriately in a UCF.

The constraints defined in this section are implemented in the UCF for the example design delivered with the core. Sections from this UCF are copied into the following descriptions

to provide examples. These examples should be studied in conjunction with the HDL source code for the example design.

## PERIOD Constraints for Clock Nets

### gtx\_clk

The clock provided to gtx\_clk must be constrained for a clock frequency of 125 MHz.

The following UCF syntax shows the necessary constraints being applied to the gtx\_clk\_bufg signal (which is routed to the gtx\_clk port of the core):

```
# Set the Transmitter clock period constraints: please do not relax
NET "gtx_clk_bufg"          TNM_NET = "clk_tx";
TIMEGRP "tx_clock"         = "clk_tx";
TIMESPEC "TS_tx_clk"       = PERIOD "tx_clock" 8000 ps HIGH 50 %;
```

### gmii\_rx\_clk

The clock provided to gmii\_rx\_clk must be constrained for a clock frequency of 125 MHz.

The following UCF syntax shows the necessary constraints being applied to the gmii\_rx\_clk\_bufg signal (which is routed to the gmii\_rx\_clk port of the core):

```
# Set the Receiver clock period constraints: please do not relax
NET "gmii_rx_clk_bufg"     TNM_NET = "clk_rx";
TIMEGRP "rx_clock"         = "clk_rx";
TIMESPEC "TS_rx_clk"       = PERIOD "rx_clock" 8000 ps HIGH 50 %;
```

### host\_clk

The clock provided to host\_clk must be constrained to the desired frequency within the allowable range (see [“Host Clock Frequency,” on page 71](#)).

The following UCF syntax shows a 100 MHz period constraint being applied to the host\_clk signal (which is routed to the host\_clk port of the core):

```
# Set the Management Clock period constraints: relax as required
NET "host_clk" TNM_NET = "host_clk";
TIMEGRP "host" = "host_clk" EXCEPT "mdio_logic";
TIMESPEC "TS_host_clk" = PERIOD "host" 10000 ps HIGH 50 %;
```

## MDIO Logic

The MDIO logic (see [“MDIO Interface,” on page 28](#)) is synchronous to host\_clk, but data only changes at the mdc output rate (as configured in the [“MDIO Configuration,” on page 75](#)). Nominally mdc will be set to a frequency of 2.5 MHz. Every flip-flop in the MDIO logic is clocked with host\_clk, but is sent a clock enable pulse at the mdc frequency. To prevent this logic being over constrained by the host\_clk period, the relevant flip-flops for the MDIO logic can be grouped together and removed from the host\_clk period constraint. This is shown in the previous UCF syntax for host\_clk. The UCF syntax which follows targets the MDIO logic flip-flops and groups them together. Reduced clock period constraints are then applied.

```
#####
# MDIO Constraints: please do not edit                                     #
#####

# Place the MDIO logic in it's own timing groups
INST "*gmac_core/BU2/U0/MANIFGEN?MANAGEN/PHY/ENABLE_REG"  TNM = "mdc_rising";
INST "*gmac_core/BU2/U0/MANIFGEN?MANAGEN/PHY/READY_INT"   TNM = "mdc_rising";
INST "*gmac_core/BU2/U0/MANIFGEN?MANAGEN/PHY/STATE_COUNT*" TNM = FFS "mdc_rising";
INST "*gmac_core/BU2/U0/MANIFGEN?MANAGEN/PHY/MDIO_TRISTATE" TNM = "mdc_falling";
```



```
INST " *gmac_core/BU2/U0/MANIFGEN?MANAGEN/PHY/MDIO_OUT"      TNM = "mdc_falling";
TIMEGRP "mdio_logic" = "mdc_rising" "mdc_falling";

TIMESPEC "TS_mdio1" = PERIOD "mdio_logic" 400 ns;
TIMESPEC "TS_mdio2" = FROM "mdc_rising" TO "mdc_falling" 200 ns;
```

## Timespecs for Critical Logic within the Core

Signals must cross clock domains at certain points in the core, as described in the following sections.

### Flow Control

Pause requests are received and decoding in the `gmii_rx_clk` domain and must be transferred into the `gtx_clk` domain to pause the transmitter. Therefore, whenever `gmii_rx_clk` and `gtx_clk` are derived from different clock sources, the following constraints must always be applied:

```
# Flow Control logic reclocking
INST " *gmac_core/BU2/U0/FLOW/RX_PAUSE/GOOD_FRAME_TO_TX"      TNM="flow_rx_to_tx";
INST " *gmac_core/BU2/U0/FLOW/RX_PAUSE/PAUSE_REQ_TO_TX"      TNM="flow_rx_to_tx";
INST " *gmac_core/BU2/U0/FLOW/RX_PAUSE/AUSE_VALUE_TO_TX*"      TNM="flow_rx_to_tx";
TIMESPEC "TS_flow_rx_to_tx" = FROM "flow_rx_to_tx" TO "tx_clock" 8000 ps;
```

### Configuration

When the optional Management Interface is used with the core, configuration information is written synchronously to `host_clk`. Receiver configuration data must be transferred onto the `gmii_rx_clk` clock domain for use with the receiver and into the `gtx_clk` clock domain for use with the transmitter. The following UCF syntax targets this logic and a timing ignore attribute (TIG) is applied. It does not matter when configuration changes take place—the current configurations are sampled between frames by both the receiver and transmitter.

```
# Configuration Register reclocking
INST " *gmac_core/BU2/U0/MANIFGEN?MANAGEN/CONF/RX0_OUT*"      TNM="config_to_rx";
INST " *gmac_core/BU2/U0/MANIFGEN?MANAGEN/CONF/RX1_OUT*"      TNM="config_to_rx";
INST " *gmac_core/BU2/U0/MANIFGEN?MANAGEN/CONF/FC_OUT_29"      TNM="config_to_rx";
TIMESPEC "TS_config_to_rx" = FROM "config_to_rx" TO "rx_clock" TIG;

INST " *gmac_core/BU2/U0/MANIFGEN?MANAGEN/CONF/TX_OUT*"      TNM="config_to_tx";
INST " *gmac_core/BU2/U0/MANIFGEN?MANAGEN/CONF/FC_OUT_30"      TNM="config_to_tx";
TIMESPEC "TS_config_to_tx" = FROM "config_to_tx" TO "tx_clock" TIG;
```

## Timespecs for Reset Logic within the Core

Internally, the core is divided into clock/reset domains that group elements with common clock and reset signals. The reset circuitry for one of these domains is illustrated in [Figure 10-5](#). This circuit provides controllable skews on the reset nets within the design. More information on the operation and rationale behind this circuit can be found in Ken Chapman's Xilinx TechXclusive, "Get Smart About Reset" at:

[www.xilinx.com/support/techxclusives/global-techX19.htm](http://www.xilinx.com/support/techxclusives/global-techX19.htm)

The following UCF syntax identifies the relevant reset logic and groups them together. Timing constraints are then applied to constrain the skews on the reset nets:

```
INST " *gmac_core/BU2/U0/SYNC_RX_RESET_I/RESET_OUT" TNM = "reset_dist_grp";
INST " *gmac_core/BU2/U0/SYNC_TX_RESET_I/RESET_OUT" TNM = "reset_dist_grp";
INST " *gmac_core/BU2/U0/G_SYNC_MGMT_RESET?SYNC_MGMT_RESET_HOST_I/RESET_OUT" TNM =
"reset_dist_grp";
TIMESPEC "TS_reset_dist" = FROM "reset_dist_grp" 6100 ps;
```

**Note:** The third line in the above UCF syntax is only required when the optional Management Interface is used.

## Constraints when Implementing an External GMII

The constraints defined in this section are implemented in the UCF for the example design delivered with the core. Sections from this UCF are copied into the descriptions below to provide examples. These examples should be studied in conjunction with the HDL source code for the example design and with the description [“Implementing External GMII,” on page 55.](#)

### GMII IOB Constraints

The following constraints target the flip-flops that are inferred in the top-level HDL file for the example design; constraints are set to ensure that these are placed in IOBs.

```
# GMII Transmitter Constraints: place flip-flops in IOB
INST "*gmii_interface/gmii_txd_reg*"      IOB = true;
INST "*gmii_interface/gmii_tx_en_reg"     IOB = true;
INST "*gmii_interface/gmii_tx_er_reg"     IOB = true;

# GMII Receiver Constraints: place flip-flops in IOB
INST "*gmii_interface/gmii_rxd_reg*"      IOB = true;
INST "*gmii_interface/gmii_rx_dv_reg"     IOB = true;
INST "*gmii_interface/gmii_rx_er_reg"     IOB = true;
```

The GMII is a 3.3 volt signal-level interface. The 3.3 volt LVTTTL SelectIO standard is the default for Virtex-II devices. The following constraints may be added without harm. The 3.3 volt LVTTTL SelectIO standard is not the default for Virtex-5, Virtex-4, Virtex-II Pro, Spartan-3, Spartan-3E and Spartan-3A devices. Use the following constraints with the device IO Banking rules.

```
INST "gmii_txd<?>"      IOSTANDARD = LVTTTL;
INST "gmii_tx_en"        IOSTANDARD = LVTTTL;
INST "gmii_tx_er"        IOSTANDARD = LVTTTL;

INST "gmii_rxd<?>"      IOSTANDARD = LVTTTL;
INST "gmii_rx_dv"        IOSTANDARD = LVTTTL;
INST "gmii_rx_er"        IOSTANDARD = LVTTTL;

INST "gmii_tx_clk"        IOSTANDARD = LVTTTL;
INST "gmii_rx_clk"        IOSTANDARD = LVTTTL;
```

In addition, the example design provides pad locking on the GMII for several families. This is provided as a guideline only; there are no specific I/O location constraints for this core.

## GMII Input Setup/Hold Timing

Figure 9-1 and Table 9-1 illustrate the setup and hold time window for the input GMII signals. This is the worst-case data valid window presented to the FPGA device pins.

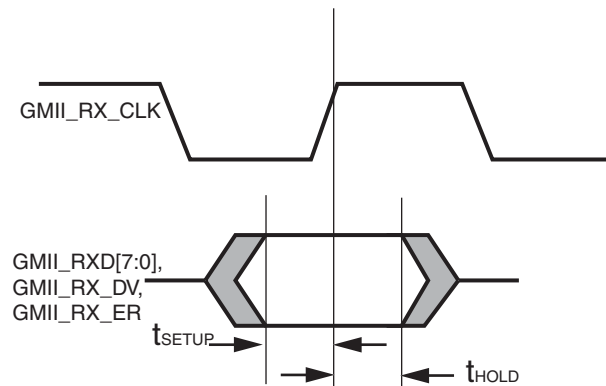


Figure 9-1: Input GMII Timing

Observe that there is a 2 ns data valid window which is presented across the GMII input bus. This must be correctly sampled by the FPGA devices.

Table 9-1: Input GMII Timing

Symbol	Min	Max	Units
$t_{\text{SETUP}}$	2.00	-	ns
$t_{\text{HOLD}}$	0.00	-	ns

### Virtex-II, and Virtex-II Pro Devices

These families have input delay elements (which are always of a fixed delay) that are automatically inserted by the Xilinx tools and are set to provide a zero-hold time. These input delays will automatically meet input setup and hold timing on the GMII without any specific constraints.

### Spartan-3, Spartan-3E, and Spartan-3A Devices

The GMII design uses a DCM on the receiver clock domain for Spartan-3, Spartan-3E, and Spartan-3A devices. Phase-shifting is then applied to the DCM to align the resultant clock so that it will correctly sample the 2 ns GMII data valid window at the input flip-flops.

The fixed phase-shift is applied to the DCM using the following UCF syntax:

```
INST *gmii_interface/gmii_rxc_dcm CLKOUT_PHASE_SHIFT = FIXED;
INST *gmii_interface/gmii_rxc_dcm PHASE_SHIFT = 0;
```

The value of PHASE\_SHIFT is preconfigured in the example designs to meet the setup and hold constraints for the example GMII pinout in the particular device. The setup/hold timing which is achieved after place-and-route is reported in the datasheet section of the TRCE report (created by the implement script).

For customers fixing their own pinout, the setup and hold figures reported in the TRCE report can be used to initially setup the approximate DCM phase-shift. [Appendix C, "Calculating DCM Phase-Shifting"](#) describes a more accurate method for fixing the phase-shift by using hardware measurement of a unique PCB design.

## Virtex-4 Devices

The GMII design uses IDELAY components on the receiver clock, data and control signals for Virtex-4 devices. A fixed tap delay can be applied to either delay the data and control signals or delay the clock so that the data/control are correctly sampled by the `gmii_rx_clk` clock at the IOB flip-flop, meeting GMII setup and hold timing.

The choice of delaying data/control or clock is dependent upon a number of factors, not least being the required shift. There are trade-offs to be made with either choice. Delaying the clock is clock-period specific as we move the clock to line up each edge with data from the following edge. Delaying the data/control introduces more jitter which degrades the overall setup/hold window. The interface timing report in the two cases is also quite different. See [“Understanding Timing Reports for GMII Setup/Hold Timing.”](#)

The following constraint shows an example of setting the delay value for one of these IDELAY components. Data/Control bits can be adjusted individually, if desired, to compensate for any PCB routing skew.

```
INST *gmii_interface/delay_gmii_rx_dv IOBDELAY_VALUE = 53;
```

The value of `IOBDELAY_VALUE` is preconfigured in the example designs to meet the setup and hold constraints for the example GMII pinout in the particular device. The setup/hold timing which is achieved after place-and-route is reported in the datasheet section of the TRCE report (created by the implement script).

When IDELAY or IODELAY primitives are instantiated with a fixed delay attribute, an IDELAYCTRL component must be also instantiated to continuously calibrate the individual input delay elements. The IDELAYCTRL module requires a reference clock, which is assumed to be an input to the example design delivered by CORE Generator. The most efficient way to use the IDELAYCTRL module is to lock the placement of the instance to the clock region of the device where the IDELAY/IODELAY components are placed. An example LOC constraint for the IDELAYCTRL module is shown commented out in the UCF. See the *Virtex-4 User Guide* and code comments for more information.

## Virtex-5 Devices

The GMII design uses IODELAY components on the receiver clock, data and control signals for Virtex-5 devices. A fixed tap delay can be applied to either delay the data and control signals or delay the clock so that the data/control are correctly sampled by the `gmii_rx_clk` clock at the IOB flip-flop, meeting GMII setup and hold timing.

The choice of delaying data/control or clock is dependant upon a number of factors, not least being the required shift. There are trade-offs to be made with either choice. Delaying the clock is clock-period specific as we move the clock to line up each edge with data from the following edge. Delaying the data/control introduces more jitter which degrades the overall setup/hold window. The interface timing report in the two cases is also quite different. See [“Understanding Timing Reports for GMII Setup/Hold Timing.”](#)

The following constraint shows an example of setting the delay value for one of these IODELAY components. Data/Control bits can be adjusted individually, if desired, to compensate for any PCB routing skew.

```
INST *gmii_interface/delay_gmii_rx_dv IDELAY_VALUE = 33;
```

The value of `IDELAY_VALUE` is preconfigured in the example designs to meet the setup and hold constraints for the example GMII pinout in the particular device. The setup/hold timing which is achieved after place-and-route is reported in the datasheet section of the TRCE report (created by the implement script).

When IDELAY or IODELAY primitives are instantiated with a fixed delay attribute, an IDELAYCTRL component must be also instantiated to continuously calibrate the individual input delay elements. The IDELAYCTRL module requires a reference clock, which is assumed to be an input to the example design delivered by CORE Generator. The most efficient way to use the IDELAYCTRL module is to lock the placement of the instance to the clock region of the device where the IDELAY/IODELAY components are placed. An example LOC constraint for the IDELAYCTRL module is shown commented out in the UCF. See the *Virtex-5 User Guide* and code comments for more information.

## Understanding Timing Reports for GMII Setup/Hold Timing

### Non-Virtex-4 or Virtex-5 devices

Setup and Hold results for the GMII input bus can be found in the data sheet section of the Timing Report. The results are self-explanatory and it is easy to see how they relate to [Figure 9-1](#). Here follows an example for the GMII report from a Virtex-2 device. The implementation requires 1.531 ns of setup—this is less than the 2 ns required so there is some slack. The implementation requires -0.125 ns of hold— this is less than the 0 ns required so there is some slack.

Data Sheet report:

-----

All values displayed in nanoseconds (ns)

Setup/Hold to clock gmii\_rx\_clk

Source	Setup to clk (edge)	Hold to clk (edge)	Internal Clock(s)	Clock Phase
gmii_rx_dv	1.531 (R)	-0.141 (R)	gmii_rx_clk_bufg	0.000
gmii_rx_er	1.531 (R)	-0.141 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<0>	1.531 (R)	-0.141 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<1>	1.525 (R)	-0.135 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<2>	1.531 (R)	-0.141 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<3>	1.525 (R)	-0.135 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<4>	1.515 (R)	-0.125 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<5>	1.515 (R)	-0.125 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<6>	1.520 (R)	-0.130 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<7>	1.520 (R)	-0.130 (R)	gmii_rx_clk_bufg	0.000

### Virtex-4 or Virtex-5 devices with Delayed Data/Control

Setup and Hold results for the GMII input bus can be found in the data sheet section of the Timing Report. The results are self-explanatory and it is easy to see how they relate to [Figure 9-1](#). Here follows an example for the GMII report from a Virtex-5 device. The implementation requires 1.962 ns of setup—this is less than the 2 ns required, so there is slack. The implementation requires -0.008 ns of hold—this is less than the 0 ns required, so there is slack.

Data Sheet report:

-----

All values displayed in nanoseconds (ns)

Setup/Hold to clock gmii\_rx\_clk

-----

Source	Setup to clk (edge)	Hold to clk (edge)	Internal Clock(s)	Clock Phase
gmii_rx_dv	1.955 (R)	-0.017 (R)	gmii_rx_clk_bufg	0.000
gmii_rx_er	1.962 (R)	-0.031 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<0>	1.949 (R)	-0.013 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<1>	1.944 (R)	-0.009 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<2>	1.947 (R)	-0.012 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<3>	1.942 (R)	-0.008 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<4>	1.950 (R)	-0.015 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<5>	1.962 (R)	-0.026 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<6>	1.957 (R)	-0.022 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<7>	1.952 (R)	-0.020 (R)	gmii_rx_clk_bufg	0.000

## Virtex-4 or Virtex-5 Devices with Delayed Clock

Setup and Hold results for the GMII input bus can be found in the data sheet section of the Timing Report. However, depending on how the setup/hold requirements have been met, it may not be immediately obvious how the results relate to Figure 9-1. The following is an example for the GMII report from a Virtex-4 device where the clock has been delayed to meet the setup/hold requirements.

Data Sheet report:

-----  
All values displayed in nanoseconds (ns)

Setup/Hold to clock gmii\_rx\_clk

Source	Setup to clk (edge)	Hold to clk (edge)	Internal Clock(s)	Clock Phase
gmii_rx_dv	-6.198 (R)	7.526 (R)	gmii_rx_clk_bufg	0.000
gmii_rx_er	-6.225 (R)	7.554 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<0>	-6.149 (R)	7.484 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<1>	-6.152 (R)	7.486 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<2>	-6.206 (R)	7.532 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<3>	-6.207 (R)	7.533 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<4>	-6.134 (R)	7.476 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<5>	-6.134 (R)	7.476 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<6>	-6.170 (R)	7.506 (R)	gmii_rx_clk_bufg	0.000
gmii_rxd<7>	-6.170 (R)	7.506 (R)	gmii_rx_clk_bufg	0.000

The implementation requires -6.134 ns of setup. Figure 9-2 illustrates that this represents a figure of 1.866 ns relative to the following rising edge of the clock (since the IDELAY has acted to delay the clock by an entire period when measured from the input flip-flop). This is less than the 2 ns required, so there is slack.

The implementation requires 7.554 ns of hold. Figure 9-2 illustrates that this represents a figure of -0.446 ns relative to the following rising edge of the clock (since the IDELAY has acted to delay the clock by an entire period when measured from the input flip-flop). This is less than the 0 ns required so there is slack.

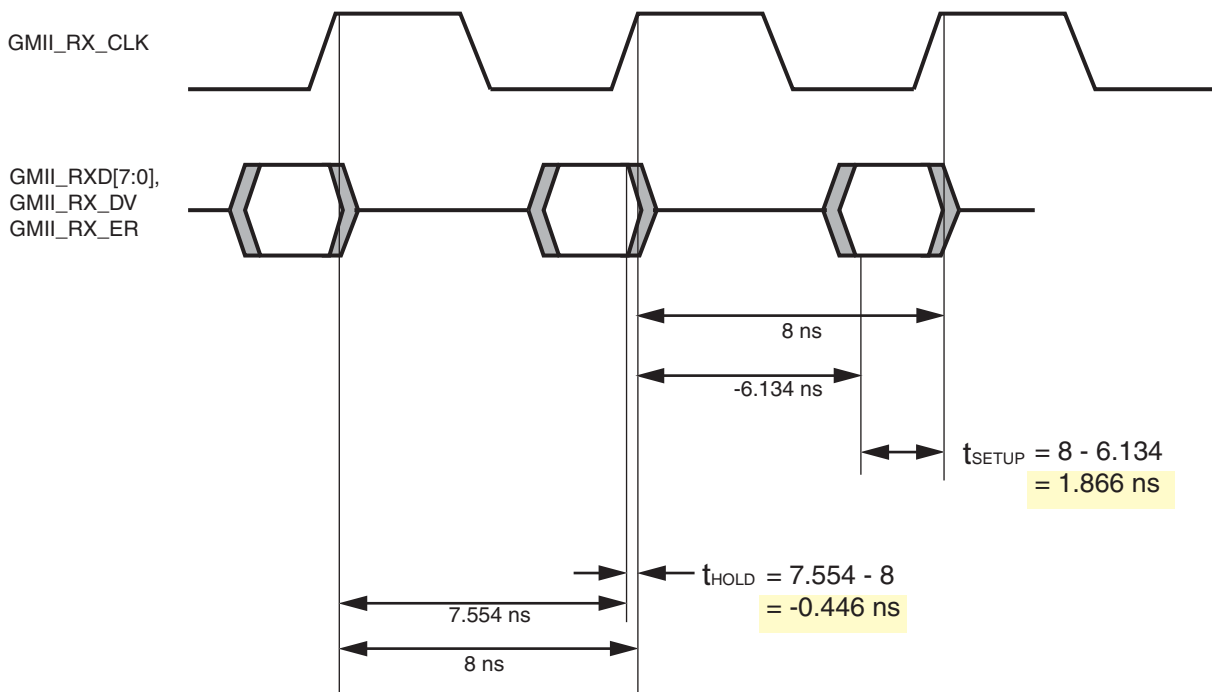


Figure 9-2: Timing Report Setup/Hold Illustration

## Constraints when Implementing an External RGMII

The constraints defined in this section are implemented in the UCF for the example design delivered with the core. Sections from this UCF are copied into the descriptions below to provide examples. These examples should be studied in conjunction with the HDL source code for the example design and with the description “Implementing External GMII,” on page 55.

### RGMII IOB Constraints

The following constraints target the flip-flops that are inferred in the top level HDL file for the example design. Constraints are set to ensure that these are placed in IOBs. The DDR register constraints are not present for a Virtex-4 or Virtex-5 device where DDR components are instantiated rather than inferred.

```
# RGMII Receiver Constraints: place DDR registers in IOB
INST "*rgmii_interface/rgmii_rxd_ddr*" IOB = true;
INST "*rgmii_interface/rgmii_rx_dv_ddr" IOB = true;
INST "*rgmii_interface/rgmii_rx_ctl_ddr" IOB = true;

# Inband Status Registers: place registers in IOB
INST "*rgmii_interface/link_status" IOB = true;
INST "*rgmii_interface/clock_speed*" IOB = true;
INST "*rgmii_interface/duplex_status" IOB = true;
```

The RGMII v2.0 is a 1.5 volt signal-level interface. The 1.5 volt HSTL Class I SelectIO standard is used for RGMII interface pins. Use the following constraints with the device IO Banking rules. The IO slew rate is set to fast to ensure that the interface can meet setup and hold times.

```

INST "rgmii_txd<?>"      IOSTANDARD = HSTL_I;
INST "rgmii_tx_ctl"      IOSTANDARD = HSTL_I;
INST "rgmii_rxd<?>"      IOSTANDARD = HSTL_I;
INST "rgmii_rx_ctl"      IOSTANDARD = HSTL_I;

INST "rgmii_txc"          IOSTANDARD = HSTL_I;
INST "rgmii_rxc"          IOSTANDARD = HSTL_I;

INST "rgmii_txd<?>"      SLEW = FAST;
INST "rgmii_tx_ctl"      SLEW = FAST;
INST "rgmii_txc"          SLEW = FAST;

```

In addition, the example design provides pad locking on the RGMII for several families. This is provided as a guideline only; there are no specific I/O location constraints for this core.

## RGMII Input Setup/Hold Timing

Figure 9-3 and Table 9-2 illustrate the setup and hold time window for the input RGMII signals. This is the worst-case data valid window presented to the FPGA device pins.

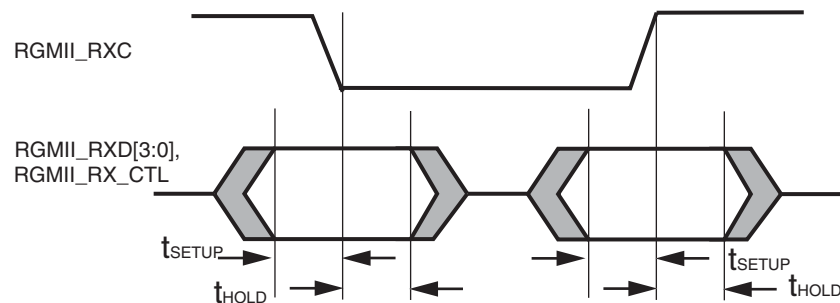


Figure 9-3: Input RGMII Timing

Note the 2 ns data valid window which is presented across the RGMII input bus. This must be correctly sampled on both clock edges by the FPGA devices.

Table 9-2: Input RGMII Timing

Symbol	Min	Typical	Units
$t_{\text{SETUP}}$	1.0	2.0	ns
$t_{\text{HOLD}}$	1.0	2.0	ns

For RGMII, the lower data bits, `rgmii_rxd[3:0]`, should be sampled internally on the rising edge of `rgmii_rxc`, and the upper data bits, `rgmii_rxd[7:4]`, should be sampled internally on the falling edge of `rgmii_rxc`.

### Virtex-II, Virtex-II Pro, Spartan-3, Spartan-3E, and Spartan-3A Devices

The RGMII design uses a DCM on the receiver clock domain for all devices except Virtex-4 and Virtex-5. Phase-shifting is then applied to the DCM to align the resultant clock so that it will correctly sample the 2 ns RGMII data valid window at the input flip-flops.



The fixed phase-shift is applied to the DCM using the following UCF syntax.

```
INST *rgmii_interface/rgmii_rxc_dcm CLKOUT_PHASE_SHIFT = FIXED;
INST *rgmii_interface/rgmii_rxc_dcm PHASE_SHIFT = 40;
```

The value of `PHASE_SHIFT` is preconfigured in the example designs to meet the setup and hold constraints for the example RGMII pinout in the particular device. The setup/hold timing which is achieved after place-and-route is reported in the datasheet section of the TRCE report (created by the implement script).

For customers fixing their own pinout, the setup and hold figures reported in the TRCE report can be used to initially setup the approximate DCM phase-shift. [Appendix C, “Calculating DCM Phase-Shifting”](#) describes a more accurate method for fixing the phase-shift by using hardware measurement of a unique PCB design.

### Virtex-4 Devices

The RGMII design uses IDELAY components on the receiver clock, data and control signals for Virtex-4 devices. A fixed tap delay can be applied to either delay the data and control signals or to delay the clock so that the data/control are correctly sampled by the `rgmii_rxc` clock at the IOB IDDR registers—thus meeting RGMII setup and hold timing.

The choice of delaying data/control or clock is dependent upon a number of factors, not least being the required shift. There are trade-offs to be made with either choice. Delaying the clock is clock-period specific as we move the clock to line up each edge with data from the following edge. Delaying the data/control introduces more jitter which degrades the overall setup/hold window. The interface timing report in the two cases is also quite different. See [“Understanding Timing Reports for RGMII Setup/Hold timing.”](#)

The following constraint shows an example of setting the delay value for one of these IDELAY components. Data/Control bits can be adjusted individually to compensate for any PCB routing skew if desired.

```
INST *rgmii_interface/delay_rgmii_rx_ctl IOBDELAY_TYPE = "FIXED";
INST *rgmii_interface/delay_rgmii_rx_ctl IOBDELAY_VALUE = 40;
```

The value of `IOBDELAY_VALUE` is preconfigured in the example designs to meet the setup and hold constraints for the example RGMII pinout in the particular device. The setup/hold timing which is achieved after place-and-route is reported in the datasheet section of the TRCE report (created by the implement script).

When IDELAY or IODELAY primitives are instantiated with a fixed delay attribute, an IDELAYCTRL component must be also instantiated to continuously calibrate the individual input delay elements. The IDELAYCTRL module requires a reference clock, which is assumed to be an input to the example design delivered by CORE Generator. The most efficient way to use the IDELAYCTRL module is to lock the placement of the instance to the clock region of the device where the IDELAY/IODELAY components are placed. An example LOC constraint for the IDELAYCTRL module is shown commented out in the UCF. See the *Virtex-4 User Guide* and code comments for more information.

### Virtex-5 Devices

The RGMII design uses IODELAY components on both the receiver and transmitter clock domains for Virtex-5 devices. A fixed tap delay is applied to the `rgmii_txc` output clock to move the rising edge of this clock to the centre of the output data window. For the receiver clock, data and control signals, a fixed tap delay can be applied to either delay the data and control signals or delay the clock so that the data/control are correctly sampled by the `rgmii_rxc` clock at the IOB IDDR registers, meeting RGMII setup and hold timing.

The choice of delaying data/control or clock is dependant upon a number of factors, not least being the required shift. There are trade-offs to be made with either choice. Delaying the clock is clock-period specific as we move the clock to line up each edge with data from the following edge. Delaying the data/control introduces more jitter which degrades the overall setup/hold window. The interface timing report in the two cases is also quite different. See [“Understanding Timing Reports for RGMII Setup/Hold timing.”](#)

The following constraint shows an example of setting the delay value for two of these IODELAY components. Data/Control bits can be adjusted individually to compensate for any PCB routing skew, if desired.

```
INST *rgmii_interface/delay_rgmii_tx_clk IDELAY_TYPE = "FIXED";
INST *rgmii_interface/delay_rgmii_tx_clk ODELAY_VALUE = 25;
INST *rgmii_interface/delay_rgmii_tx_clk DELAY_SRC = "O";

INST *rgmii_interface/delay_rgmii_rx_ctl IDELAY_TYPE = "FIXED";
INST *rgmii_interface/delay_rgmii_rx_ctl IDELAY_VALUE = 20;
INST *rgmii_interface/delay_rgmii_rx_ctl DELAY_SRC = "I";
```

The value of IDELAY\_VALUE is preconfigured in the example designs to meet the setup and hold constraints for the example RGMII pinout in the particular device. The setup/hold timing which is achieved after place-and-route is reported in the datasheet section of the TRCE report (created by the implement script).

When IDELAY or IODELAY primitives are instantiated with a fixed delay attribute, an IDELAYCTRL component must be also instantiated to continuously calibrate the individual input delay elements. The IDELAYCTRL module requires a reference clock, which is assumed to be an input to the example design delivered by CORE Generator. The most efficient way to use the IDELAYCTRL module is to lock the placement of the instance to the clock region of the device where the IDELAY/IODELAY components are placed. An example LOC constraint for the IDELAYCTRL module is shown commented-out in the UCF. See the *Virtex-5 User Guide* and code comments for more information.

## RGMII DDR Constraints

The following constraints are present for RGMII designs in all devices with the exception of Virtex-4 and Virtex-5 devices. Due to the use of IDDR and ODDR primitives in the Virtex-4 design, these extra clocking constraints are not required. The RGMII design uses extra clock signals for falling-edge data. These clocks require further period constraints and are constrained relative to their rising edge clock counterparts.

```
NET "not_gtx_clk_bufg"      TNM_NET = "not_clk_tx";
TIMESPEC "TS_not_tx_clk"   = PERIOD "not_clk_tx" "TS_tx_clk" PHASE + 4
nS;
NET "*rgmii_interface/not_rgmii_rx_clk_bufg" TNM_NET =
"not_rgmii_rx_clk_bufg";
TIMESPEC "TS_rnot_gmii_rx_clk_bufg" = PERIOD "not_rgmii_rx_clk_bufg"
"TS_rgmii_rx_clk_bufg" PHASE + 4 nS;
```

The RGMII design requires further clock crossing constraints to ensure timing is met when crossing from rising to falling clock edges and vice versa. A stringent time constraint ensures that timing is met with the worst-case timing allowed in the RGMII specification.

```
INST "*rgmii_interface/rgmii_rxd_reg_4"      TNM="rgmii_falling";
INST "*rgmii_interface/rgmii_rxd_reg_5"      TNM="rgmii_falling";
INST "*rgmii_interface/rgmii_rxd_reg_6"      TNM="rgmii_falling";
INST "*rgmii_interface/rgmii_rxd_reg_7"      TNM="rgmii_falling";
INST "*rgmii_interface/rgmii_rx_ctl_reg"      TNM="rgmii_falling";
INST "*rgmii_interface/gmii_rxd_reg_4"      TNM="rgmii_rising";
```

```

INST "rgmii_interface/gmii_rxd_reg_5"      TNM="rgmii_rising";
INST "rgmii_interface/gmii_rxd_reg_6"      TNM="rgmii_rising";
INST "rgmii_interface/gmii_rxd_reg_7"      TNM="rgmii_rising";
INST "rgmii_interface/gmii_rx_er_reg"      TNM="rgmii_rising";
TIMESPEC "TS_rgmii_falling_to_rising" = FROM "rgmii_falling" TO
"rgmii_rising" 3200 ps;

```

## Understanding Timing Reports for RGMII Setup/Hold timing

### Non-Virtex-4 or Virtex-5 Devices

Setup and Hold results for the RGMII input bus can be found in the data sheet section of the Timing Report. The results are self-explanatory and it is easy to see how they relate to [Figure 9-3](#).

Following is an example for the RGMII report from a Virtex-2 device. Each Input lists two sets of values—one corresponding to the –ve edge of the clock and one to the +ve edge. The first set listed corresponds to –ve edge which occurs at time 4 ns. The implementation requires 0.648 ns of setup to the –ve edge and 0.661 ns to the +ve edge. This is less than the 1 ns required, so there is slack. The implementation requires 0.300 ns of hold to the –ve edge and 0.316 ns to the +ve edge. This is less than the 1 ns required, so there is slack.

Data Sheet report:

-----

All values displayed in nanoseconds (ns)

Setup/Hold to clock rgmii\_rxc

Source	Setup to clk (edge)	Hold to clk (edge)	Internal Clock(s)	Clock Phase
rgmii_rx_ctl	-3.352 (R) 0.661 (R)	4.300 (R) 0.284 (R)	not_rgmii_rx_clk_bufg rgmii_rx_clk_bufg	4.938 0.938
rgmii_rxd<0>	-3.384 (R) 0.629 (R)	4.332 (R) 0.316 (R)	not_rgmii_rx_clk_bufg rgmii_rx_clk_bufg	4.938 0.938
rgmii_rxd<1>	-3.348 (R) 0.665 (R)	4.296 (R) 0.280 (R)	not_rgmii_rx_clk_bufg rgmii_rx_clk_bufg	4.938 0.938
rgmii_rxd<2>	-3.360 (R) 0.653 (R)	4.308 (R) 0.292 (R)	not_rgmii_rx_clk_bufg rgmii_rx_clk_bufg	4.938 0.938
rgmii_rxd<3>	-3.428 (R) 0.585 (R)	4.382 (R) 0.366 (R)	not_rgmii_rx_clk_bufg rgmii_rx_clk_bufg	4.938 0.938

### Virtex-4 or Virtex-5 devices with delayed Data/Control

Setup and Hold results for the RGMII input bus can be found in the data sheet section of the Timing Report. The results are self-explanatory and it is easy to see how they relate to [Figure 9-3](#).

Following is an example for the RGMII report from a Virtex-5 device. Each Input lists two sets of values—one corresponding to the –ve edge of the clock and one to the +ve edge. The first set listed corresponds to +ve edge which occurs at time 0 ns. The implementation requires 0.818 ns of setup to the +ve edge and 0.794 ns to the –ve edge. This is less than the 1 ns required, so there is slack. The implementation requires 0.946 ns of hold to the –ve edge and 0.972 ns to the +ve edge. This is less than the 1 ns required, so there is slack.

Data Sheet report:

-----

All values displayed in nanoseconds (ns)

Setup/Hold to clock rgmii\_rxc

Source	Setup to clk (edge)	Hold to clk (edge)	Internal Clock(s)	Clock Phase
rgmii_rx_ctl	0.810 (R)	0.933 (R)	rgmii_rx_clk_bufg	0.000
	-3.214 (F)	4.959 (F)	rgmii_rx_clk_bufg	4.000
rgmii_rxd<0>	0.811 (R)	0.940 (R)	rgmii_rx_clk_bufg	0.000
	-3.213 (F)	4.966 (F)	rgmii_rx_clk_bufg	4.000
rgmii_rxd<1>	0.801 (R)	0.946 (R)	rgmii_rx_clk_bufg	0.000
	-3.223 (F)	4.972 (F)	rgmii_rx_clk_bufg	4.000
rgmii_rxd<2>	0.818 (R)	0.929 (R)	rgmii_rx_clk_bufg	0.000
	-3.206 (F)	4.955 (F)	rgmii_rx_clk_bufg	4.000
rgmii_rxd<3>	0.809 (R)	0.936 (R)	rgmii_rx_clk_bufg	0.000
	-3.215 (F)	4.962 (F)	rgmii_rx_clk_bufg	4.000

## Virtex-4 or Virtex-5 Devices with Delayed Clock

Setup and Hold results for the RGMII input bus can be found in the data sheet section of the Timing Report. However, depending on how the setup/hold requirements have been met, it may not be immediately obvious how the results relate to Figure 9-3. Following is an example for the RGMII report from a Virtex-4 device where the clock has been delayed to meet the setup/hold requirements.

Data Sheet report:

-----

All values displayed in nanoseconds (ns)

Setup/Hold to clock rgmii\_rxc

Source	Setup to clk (edge)	Hold to clk (edge)	Internal Clock(s)	Clock Phase
rgmii_rx_ctl	-7.178 (R)	8.880 (R)	rgmii_rx_clk_bufg	0.000
	-11.178 (F)	12.880 (F)	rgmii_rx_clk_bufg	4.000
rgmii_rxd<0>	-7.192 (R)	8.893 (R)	rgmii_rx_clk_bufg	0.000
	-11.192 (F)	12.893 (F)	rgmii_rx_clk_bufg	4.000
rgmii_rxd<1>	-7.182 (R)	8.884 (R)	rgmii_rx_clk_bufg	0.000
	-11.182 (F)	12.884 (F)	rgmii_rx_clk_bufg	4.000
rgmii_rxd<2>	-7.180 (R)	8.882 (R)	rgmii_rx_clk_bufg	0.000
	-11.180 (F)	12.882 (F)	rgmii_rx_clk_bufg	4.000
rgmii_rxd<3>	-7.179 (R)	8.881 (R)	rgmii_rx_clk_bufg	0.000
	-11.179 (F)	12.881 (F)	rgmii_rx_clk_bufg	4.000

Each Input lists two sets of values—one corresponding to the +ve edge of the clock and one to the -ve edge. The first set listed corresponds to +ve edge, which occurs at time 8 ns as we have delayed the clock to use the following +ve edge.

The implementation requires -7.179 ns of setup to the +ve edge. Figure 9-4 illustrates that this represents 0.821 ns relative to the following rising edge of the clock (since the IDELAY has acted to delay the clock by an entire period when measured from the input flip-flop).

This is less than the 1 ns required, so there is slack. Equally for the -ve edge, we have -11.179 ns of setup—this edge is at time 12 ns and therefore this equates to a setup of 0.821 ns.

The implementation requires 8.893 ns of hold to the +ve edge. Figure 9-4 illustrates that this represents 0.893 ns relative to the following rising edge of the clock (since the IDELAY has acted to delay the clock by an entire period when measured from the input flip-flop). This is less than the 1 ns required, so there is slack. Equally for the -ve edge, we have 12.893 ns of hold —this edge is at time 12 ns and therefore equates to a hold time of 0.893 ns.

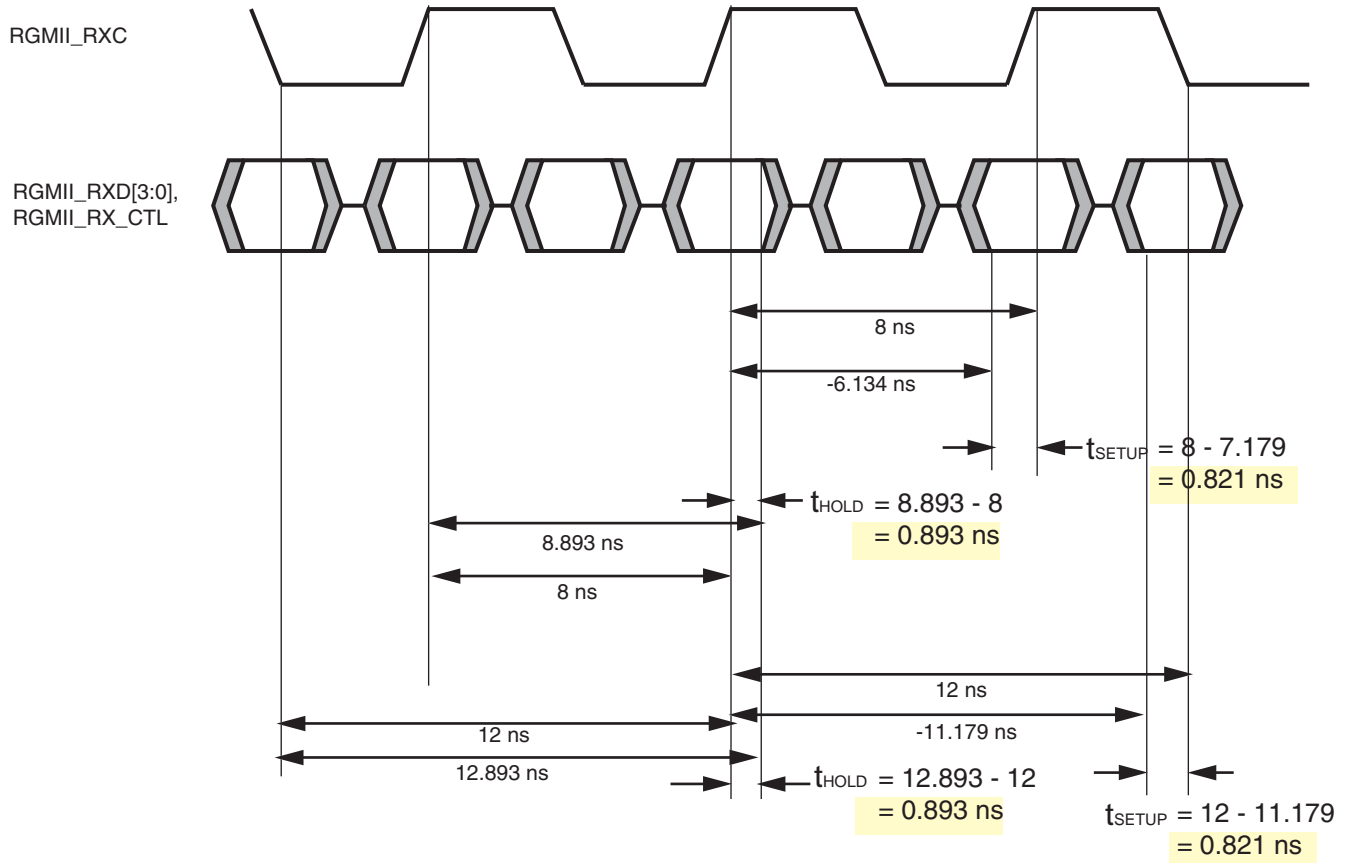


Figure 9-4: Timing Report Setup/Hold Illustration



## Clocking and Resetting

This chapter describes clock management considerations that are associated with implementing the GEMAC core. It describes the clock management logic for all implementations of the core and how clock management logic can be shared across multiple instantiations of the core. The reset circuitry within the core is also described.

### Clocking the Core

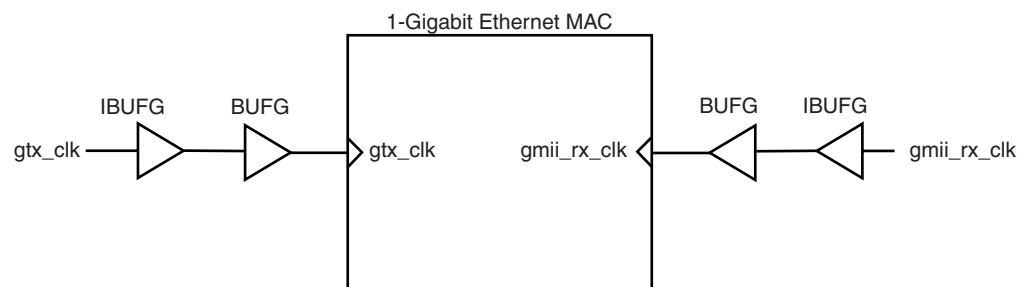
#### With Internal GMII

When the GMII-style interface of the core is used as an internal interface (for example, with an internally connected PHY core), it is likely that `gtx_clk` and `gmii_rx_clk` will be derived from the same clock source. See [Chapter 11, “Interfacing to Other Cores”](#) for an example.

#### With External GMII

[Figure 10-1](#) illustrates the clock management used with an external GMII interface. All clocks illustrated have a frequency of 125 MHz. The clock `gtx_clk` must be provided to the GEMAC core. This is a high-quality clock that satisfies the *IEEE 802.3-2002* requirements. It is expected that this clock will be derived from an external oscillator and connected into the device through an IBUFG, as illustrated in [Figure 10-1](#).

When an external GMII, `gmii_rx_clk` will usually be derived from a different clock source to `gtx_clk`. In this case, `gmii_rx_clk` is received through an IBUFG. This clock is then usually routed onto a global clock network by connecting it to a BUFG. The resulting global clock is used by all MAC receiver logic. The exception to this is the Spartan-3, Spartan-3E, and Spartan-3A devices, which require a DCM on the `gmii_rx_clk` path to meet GMII setup and hold requirements. See [“Spartan-3, Spartan-3E and Spartan-3A Devices.”](#)



**Figure 10-1: Clock Management Logic with External GMII**

## With RGMII

### Standard Clocking Scheme

Figure 10-2 illustrates the clock management used with an external RGMII interface. All clocks illustrated have a frequency of 125 MHz. The `gtx_clk` clock must be provided to the GEMAC core. This is a high-quality clock that satisfies *IEEE 802.3-2002* requirements. It is expected that this clock will be derived from an external oscillator and connected into the device through an IBUFG as illustrated in Figure 10-2. This clock is used as the input clock to a DCM from where phase-shifted clock signals are generated for use in the RGMII transmitter logic. The zero phase-shifted clock is used as the input `gtx_clk` to the GEMAC core.

The receiver clock, `rgmii_rxc` is usually derived from a different clock source to `gtx_clk`. In this case, `rgmii_rxc` will be received through an IBUFG. This clock is routed into a DCM where it is used to generate phase-shifted clock signals for use in the RGMII receiver logic. A fixed phase-shift value is applied to the DCM to meet RGMII setup and hold requirements. See Appendix C, “Calculating DCM Phase-Shifting.” The zero clock is used as the input `gmii_rx_clk` to the GEMAC core.

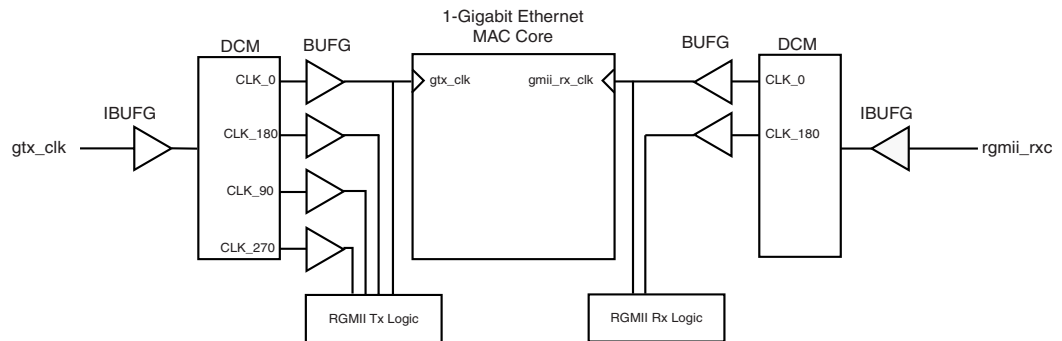


Figure 10-2: Clock Management with External RGMII

### Reducing Global Clock Buffers in Virtex-II Pro Devices

The clock schemes shown previously for RGMII are fairly expensive in global clock buffer resources, due to their use in providing both the in-phase and out-of-phase clocks for the DDR registers. Seven global clock buffers are used in the worst-case scenario (including a global clock buffer for `host_clk`).

To halve the number of clock buffers used in each DDR register implementation, consult Xilinx Application Note [XAPP685, High-Speed Clock Architecture for DDR Designs Using Local Inversion](#).

## Multiple Cores

### With External GMII

Figure 10-3 illustrates how to share clock resources across multiple instantiations of the core when using the optional GMII. `gtx_clk` may be shared between multiple cores as illustrated, resulting in a common transmitter clock domain across the device.



A common receiver clock domain is usually not possible as each core will receive an independent receiver clock from the PHY attached to the other end of the GMII. As illustrated in Figure 10-3, this results in a separate receiver clock domain for each core.

**Note:** Although not illustrated, if the optional Management Interface is used, `host_clk` can also be shared between cores.

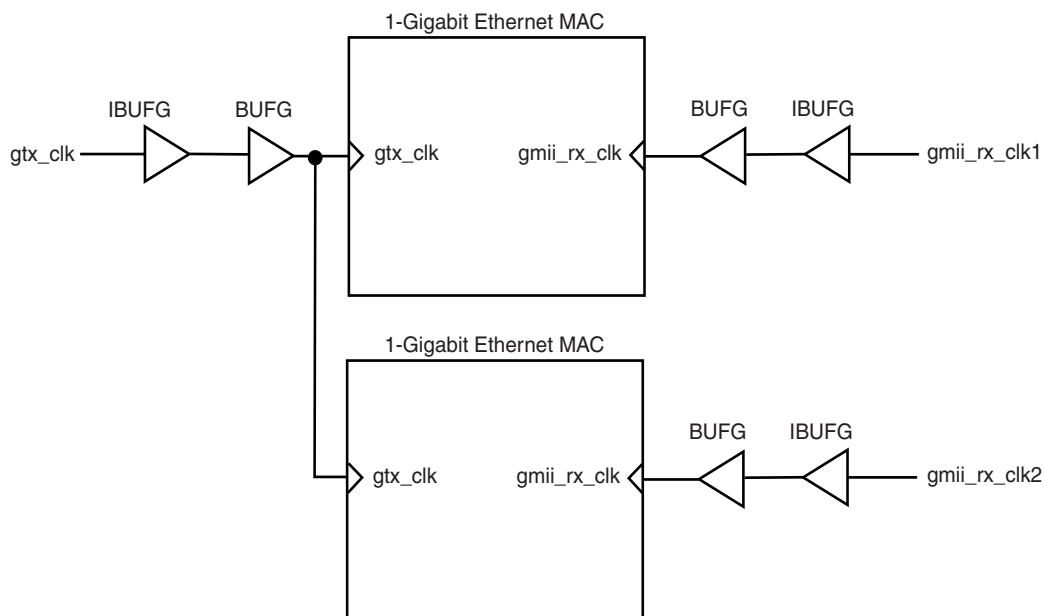


Figure 10-3: Clock Management Logic with External GMII (Multiple Cores)

## With RGMII

Figure 10-4 illustrates sharing clock resources across multiple instantiations of the core using the optional RGMII. `gtx_clk` may be shared between multiple cores as illustrated, resulting in a common transmitter clock domain across the device.

As a general rule, a common receiver clock domain is not possible. Each core receives an independent receiver clock from the PHY attached to the other end of the RGMII—as illustrated in Figure 10-4. This results in a separate receiver clock domain for each core.

**Note:** Although not illustrated, if the optional Management Interface is used, `host_clk` can also be shared between cores.

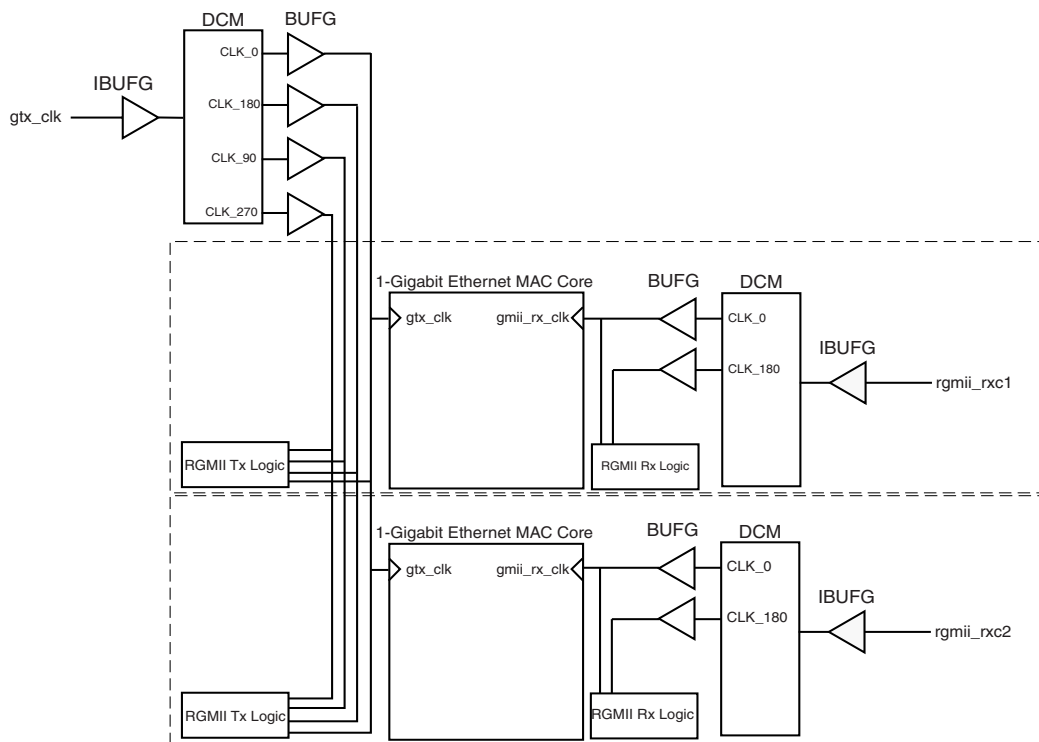


Figure 10-4: Clock Management Logic with External RGMII (Multiple Cores)

## Reset Conditions

Internally, the core is divided up into clock/reset domains that group together elements with common clock and reset signals. The reset circuitry for one of these domains is illustrated in Figure 10-5. This circuit provides controllable skews on the reset nets within the design.

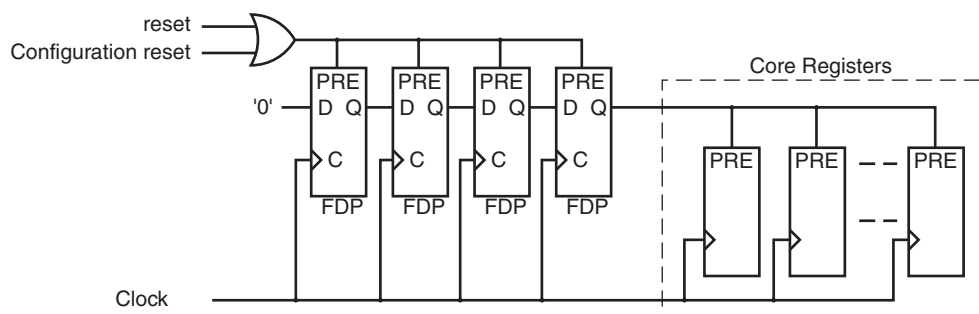


Figure 10-5: Reset Circuit for a Single Clock/reset Domain

More information on the operation and rationale behind this circuit can be found in Ken Chapman's Xilinx TechXclusive, "Get Smart About Reset" which can be found at:

[www.xilinx.com/support/techxclusives/global-techX19.htm](http://www.xilinx.com/support/techxclusives/global-techX19.htm)

## Interfacing to Other Cores

---

### Ethernet 1000Base-X PCS/PMA or SGMII Core

The GEMAC core can be integrated in a single device with the Ethernet 1000BASE-X PCS/PMA or SGMII core to extend core functionality to provide the following:

- 1000BASE-X Physical Coding Sublayer (PCS) logic designed to the *IEEE 802.3* specification with either:
  - 1000BASE-X Physical Medium Attachment (PMA) using an integrated Virtex-II Pro, Virtex-4 or Virtex-5 RocketIO™ transceiver.
  - 1000BASE-X parallel Ten-Bit-Interface (TBI) for connection to external SERDES.
- Alternatively, the Ethernet 1000BASE-X PCS/PMA or SGMII core can function as a GMII to Serial-GMII (SGMII) bridge, meaning that this can be used to provide the GEMAC core with an SGMII for serial connection to appropriate PHYs.

A description of the latest available IP Update containing the Ethernet 1000BASE-X PCS/PMA or SGMII core and instructions for obtaining the IP Update can be found on the Ethernet 1000BASE-X PCS/PMA or SGMII product page at:  
[www.xilinx.com/systemio/1gbsx\\_phy/index.htm](http://www.xilinx.com/systemio/1gbsx_phy/index.htm).

A full description of the Ethernet 1000BASE-X PCS/PMA or SGMII core is outside the scope of this document.

### Integration to Provide 1000BASE-X PCS with TBI

Figure 11-1 illustrates the connections and clock management logic required to interface the GEMAC core to the Ethernet 1000BASE-X PCS/PMA or SGMII core (when used in 1000BASE-X mode with the parallel TBI). It depicts the following:

- Direct internal connections are made between the GMII interfaces of the two cores.
- If the GEMAC has been built with the optional management logic, the MDIO port can be connected up to that of the Ethernet 1000BASE-X PCS/PMA or SGMII core to access its embedded configuration and status registers. See “[Using the Optional Management Interface](#).”
- Due to the embedded Receiver Elastic Buffer in the Ethernet 1000BASE-X PCS/PMA or SGMII core, the entire GMII is synchronous to a single clock domain. For this reason, `gtx_clk` is used as the 125 MHz reference clock for both cores and the transmitter and receiver logic of the GEMAC core now operate in the same clock domain. This allows clock crossing constraints between the `gtx_clk` and `gmii_rx_clk` clock domains to be removed from the GEMAC user constraints file (UCF). See “[Timespecs for Critical Logic within the Core](#).”

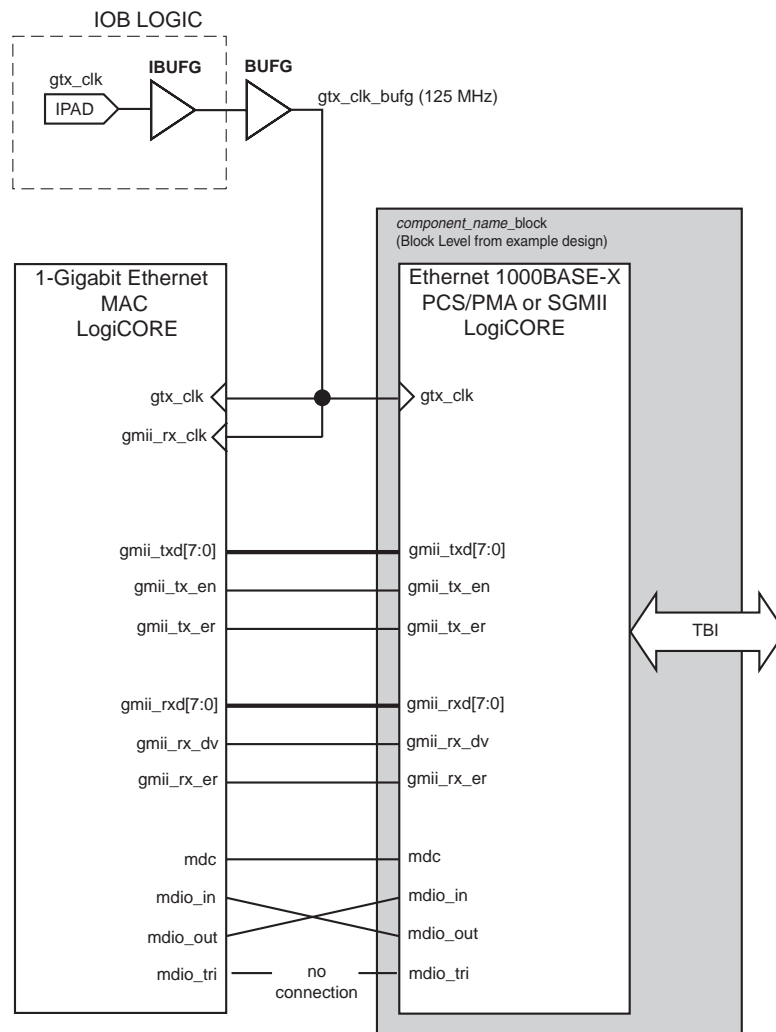


Figure 11-1: 1-Gigabit Ethernet MAC Extended to Include 1000BASE-X PCS with TBI

## Integration to Provide 1000BASE-X PCS and PMA using RocketIO

### Virtex-II Pro Devices

Figure 11-2 illustrates the connections and clock management logic required to interface the GEMAC core to the Ethernet 1000BASE-X PCS/PMA or SGMII core (when used in 1000BASE-X mode with PMA using the Virtex-II Pro RocketIO transceiver).

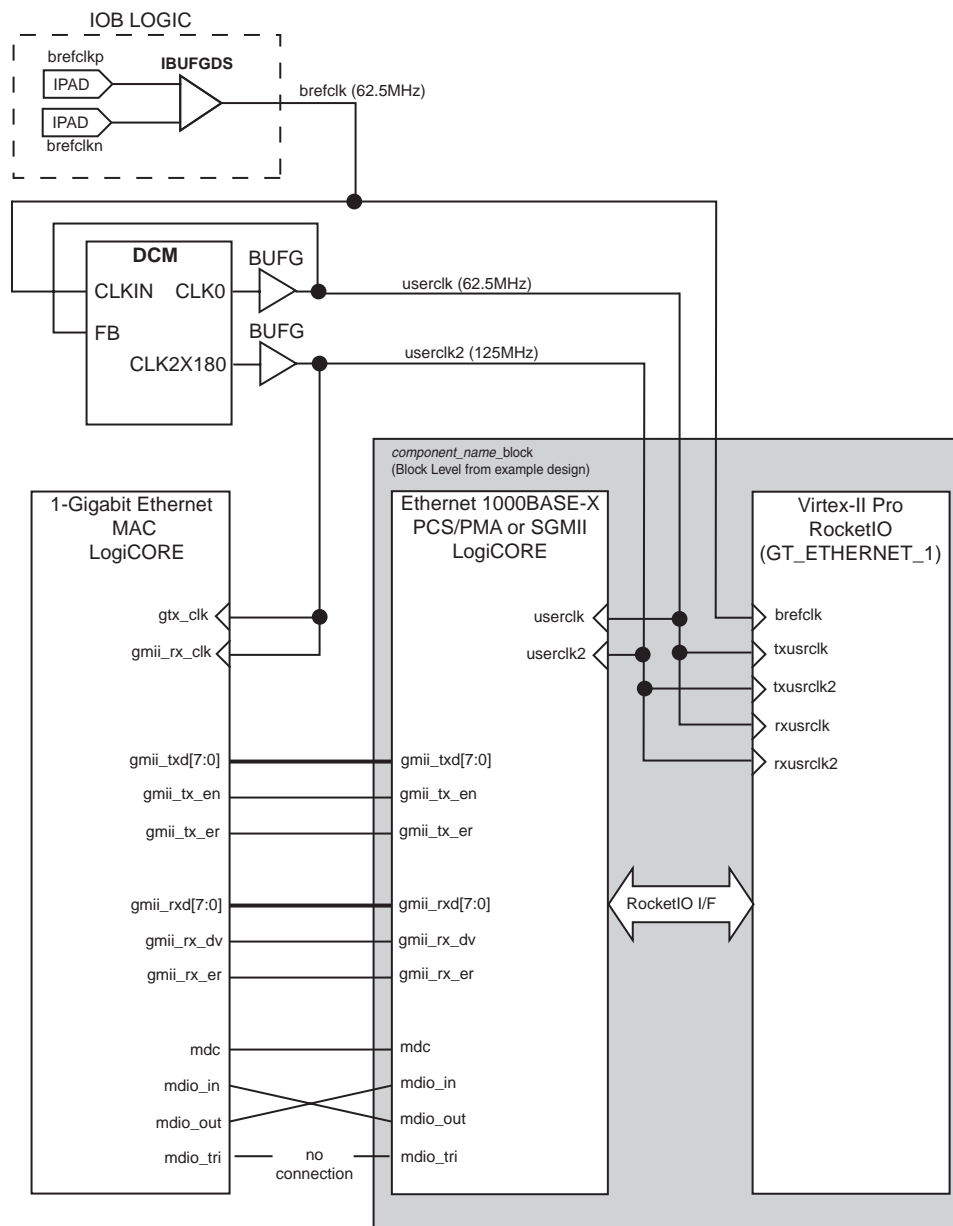


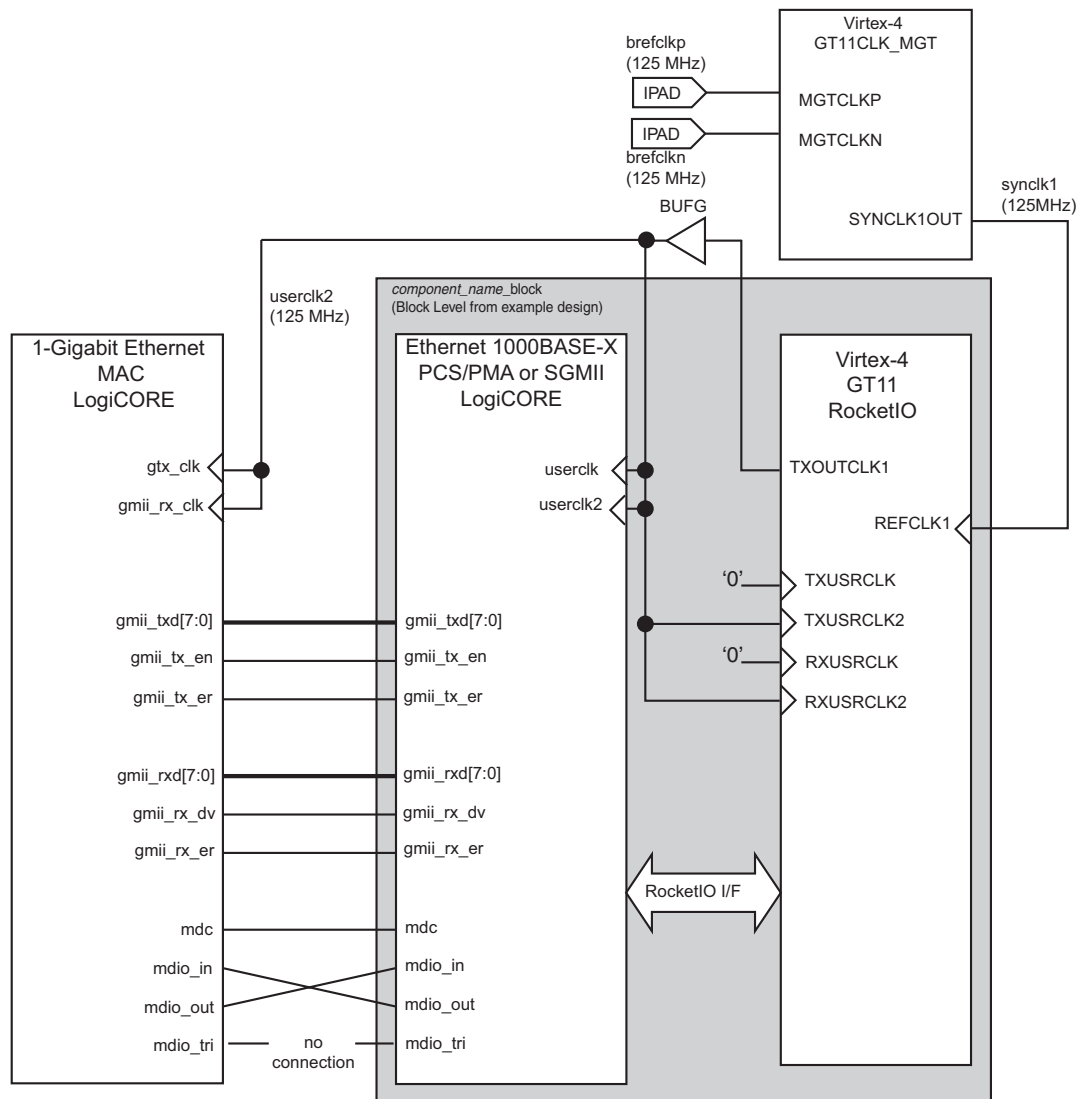
Figure 11-2: 1-Gigabit Ethernet MAC Extended to Include 1000BASE-X PCS and PMA using a Virtex-II Pro RocketIO Transceiver

Figure 11-2 illustrates the following:

- Direct internal connections are made between the GMII interfaces between the two cores.
- If the GEMAC is built with the optional management logic, the MDIO port can be connected to that of the Ethernet 1000BASE-X PCS/PMA or SGMII core to access its embedded configuration and status registers. See [“Using the Optional Management Interface.”](#)
- Due to the embedded Receiver Elastic Buffer in the Ethernet 1000BASE-X PCS/PMA or SGMII core, the entire GMII is synchronous to a single clock domain. For this reason, `userclk2` is used as the 125 MHz reference clock for both cores and the transmitter and receiver logic of the GEMAC core now operate in the same clock domain. This allows clock crossing constraints between the `gtx_clk` and `gmii_rx_clk` clock domains to be removed from the GEMAC UCF. See [“Timespecs for Critical Logic within the Core.”](#)

## Virtex-4 Devices

Figure 11-3 illustrates the connections and clock management logic required to interface the GEMAC core to the Ethernet 1000BASE-X PCS/PMA or SGMII core (when used in 1000BASE-X mode with PMA using the Virtex-4 RocketIO transceiver).



**Figure 11-3: 1-Gigabit Ethernet MAC Extended to Include 1000BASE-X PCS and PMA using a Virtex-4 RocketIO Transceiver**

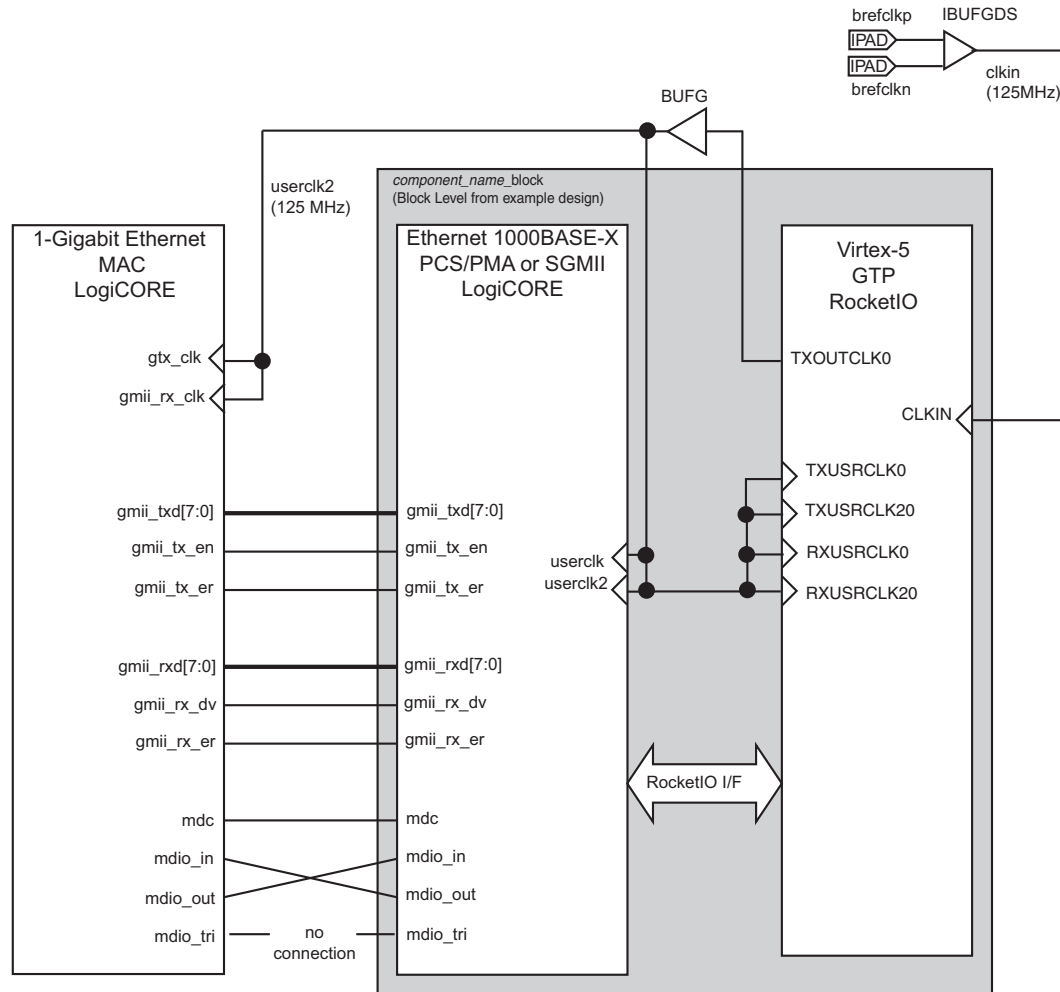
Figure 11-3 illustrates the following:

- Direct internal connections are made between the GMII interfaces between the two cores.
- If the GEMAC has been generated with the optional Management Interface, the MDIO port can be connected up to that of the Ethernet 1000BASE-X PCS/PMA or SGMII core to access its embedded configuration and status registers. See [“Using the Optional Management Interface.”](#)
- Due to the embedded Receiver Elastic Buffer in the Ethernet 1000BASE-X PCS/PMA or SGMII core, the entire GMII is synchronous to a single clock domain. For this reason, `userclk2` is used as the 125 MHz reference clock for both cores and the transmitter and receiver logic of the GEMAC core now operate in the same clock domain. This allows clock crossing constraints between the `gtx_clk` and

gmii\_rx\_clk clock domains to be removed from the GEMAC UCF. See [“Timespecs for Critical Logic within the Core.”](#)

## Virtex-5 Devices

Figure 11-4 illustrates the connections and clock management logic required to interface the GEMAC core to the Ethernet 1000BASE-X PCS/PMA or SGMII core (when used in 1000BASE-X mode with PMA using the Virtex-5 RocketIO transceiver).



**Figure 11-4: 1-Gigabit Ethernet MAC Extended to Include 1000BASE-X PCS and PMA using a Virtex-5 RocketIO Transceiver**

Figure 11-4 illustrates the following:

- Direct internal connections are made between the GMII interfaces between the two cores.
- If the GEMAC has been generated with the optional Management Interface, the MDIO port can be connected up to that of the Ethernet 1000BASE-X PCS/PMA or SGMII core to access its embedded configuration and status registers. See [“Using the Optional Management Interface.”](#)
- Due to the embedded Receiver Elastic Buffer in the Ethernet 1000BASE-X PCS/PMA or SGMII core, the entire GMII is synchronous to a single clock domain. For this



reason, `userclk2` is used as the 125 MHz reference clock for both cores and the transmitter and receiver logic of the GEMAC core now operate in the same clock domain. This allows clock crossing constraints between the `gtx_clk` and `gmii_rx_clk` clock domains to be removed from the GEMAC UCF. See “[Timespecs for Critical Logic within the Core.](#)”

## Integration to Provide SGMII Functionality

The connections between the two cores to provide SGMII functionality are identical to the connections required for 1000BASE-X PCS and PMA using the RocketIO transceiver. The only difference is that the Ethernet 1000BASE-X PCS/PMA or SGMII core is generated with the SGMII option. See “[Integration to Provide 1000BASE-X PCS and PMA using RocketIO](#)” for a description of SGMII integration.

## Ethernet Statistics Core

The Ethernet Statistics core can be integrated in a single device with the 1-Gigabit Ethernet MAC core. Using the Ethernet Statistics core allows you to:

- Count statistics based on the `rx_statistics_vector` and `tx_statistics_vector` outputs from the MAC.
- Select 32-bit or 64-bit counters.
- Specify which statistics are counted and have precise control of the conditions under which the counters are incremented.
- Read statistics optionally through the host interface of the GEMAC or independently of the MAC.

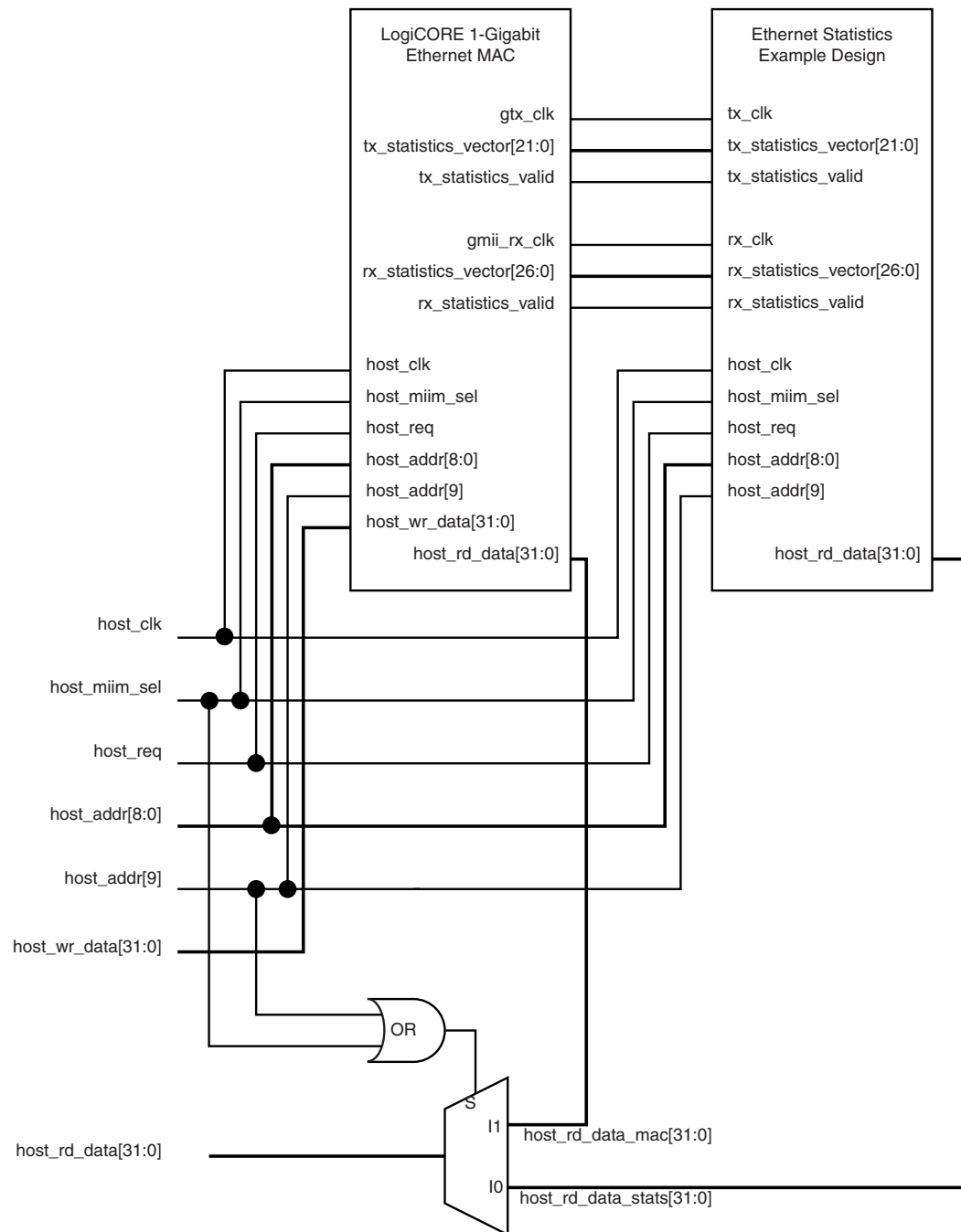
A description of the latest available IP Update containing the Ethernet Statistics core and instructions on obtaining the IP update can be found on the Ethernet Statistics product page at:

[www.xilinx.com/xlnx/xebiz/designResources/ip\\_product\\_details.jsp?key=ETHERNET\\_STATS](http://www.xilinx.com/xlnx/xebiz/designResources/ip_product_details.jsp?key=ETHERNET_STATS)

## Connecting the Ethernet Statistics core to Provide Statistics Gathering

The Ethernet Statistics core interfaces directly to the 1-Gigabit Ethernet MAC core. The Ethernet Statistics core takes in the `tx_statistics_vector` and `rx_statistics_vector` as inputs. Statistics values gathered can then be read out through the Statistics core Management Interface, that can be shared with the MAC Management Interface.

[Figure 11-5](#) illustrates connecting the Ethernet Statistics core to the MAC.



**Figure 11-5: Interfacing the Ethernet Statistics to the 1-Gigabit Ethernet MAC**

The management interfaces of the two cores can be shared by avoiding bus conflict, as follows:

- Selecting a different address range for the statistics to that of the MAC configuration registers. This is achieved by setting `host_addr[9]` to logic 0 when reading from the statistics and logic 1 when writing and reading to the MAC configuration registers.
- Using the `host_miim_sel` signal to differentiate between a statistical counter read and a MAC initiated MDIO transaction. This is achieved by setting `host_miim_sel`

to logic 0 for a statistical counter read and logic 1 for a MAC initiated MDIO transaction.

Table 11-1 describes the type of host transactions that occur if the host interface is shared (as illustrated in Figure 11-5).

**Table 11-1: Management Interface Transaction Types**

Transaction	host_miim_sel	host_addr[9]
Configuration	0	1
MIIM access	1	X
Statistics Read	0	0



# Implementing Your Design

---

This chapter describes how to simulate and implement your design containing the GEMAC core.

## Pre-implementation Simulation

A unit delay structural model of the GEMAC core netlist is provided as a CORE Generator output file. This can be used for simulation of the block in the design phase of the project.

### Using the Simulation Model

For information about setting up your simulator to use the pre-implemented model, see the Xilinx *Synthesis and Verification Design Guide*, included in your Xilinx software installation.

The unit delay structural model of the GEMAC core can be found in the CORE Generator project directory. Details of the CORE Generator outputs can be found in the *1-Gigabit Ethernet MAC Getting Started Guide*.

#### VHDL

- `<component_name>.vhd`

#### Verilog

- `<component_name>.v`

## Synthesis

### XST—VHDL

A component and instantiation template for the core named `<component_name>.vho` is provided in the CORE Generator project directory. Use this to help instance the GEMAC core into your VHDL source.

After your entire design is complete, create:

- An XST project file `top_level_module_name.prj` listing all the user source code files.
- An XST script file `top_level_module_name.scr` containing your required synthesis options.

To synthesize the design, run:

```
$ xst -ifn top_level_module_name.scr
```

See the *XST User Guide* for more information on creating project and synthesis script files, and running the xst program.

## XST—Verilog

A module declaration for the GEMAC core is provided in the CORE Generator project directory:

```
<component_name>/example_design/<component_name>.mod.v
```

Use this module to help instance the GEMAC core into your Verilog source.

After your entire design is complete, create:

- An XST project file *top\_level\_module\_name.prj* listing all the user source code files. Make sure you include

```
%XILINX%/verilog/src/ise/unisim_comp.v
```

and

```
<component_name>/example_design/component_name.mod.v
```

as the first two files in the project list.

- An XST script file *top\_level\_module\_name.scr* containing your required synthesis options.

To synthesize the design, run:

```
$ xst -ifn top_level_module_name.scr
```

See the *XST User Guide* for more information on creating project and synthesis script files, and running the xst program.

## Implementation

### Generating the Xilinx Netlist

To generate the Xilinx netlist, the ngdbuild tools are used to translate and merge the individual design netlists into a single design database, the NGD file. Also merged at this stage is the UCF for the design.

An example of the ngdbuild command is:

```
$ ngdbuild -sd path_to_core_netlist -sd path_to_user_synth_results \
    -uc top_level_module_name.ucf top_level_module_name
```

### Mapping the Design

To map the logic gates of the user design netlist into the CLBs and IOBs of the FPGA, run the map command. The map command writes out a physical design to an NCD file. An example of the map command is:

```
$ map -o top_level_module_name_map.ncd top_level_module_name.ngd \
    top_level_module_name.pcf
```

## Placing-and-Routing the Design

Execute the **par** command to place-and-route your design logic components (mapped physical logic cells) contained within an NCD file in accordance with the layout and timing requirements specified within the PCF file. The **par** command outputs the placed and routed physical design to an NCD file.

An example of the **par** command is:

```
$ par top_level_module_name_map.ncd top_level_module_name.ncd \  
    top_level_module_name.pcf
```

## Static Timing Analysis

Execute the **trce** command to evaluate timing closure on a design and create a Timing Report file (TWR) derived from static timing analysis of the Physical Design file (NCD). The analysis is typically based on constraints included in the optional PCF file.

An example of the **trce** command is:

```
$ trce -o top_level_module_name.twr top_level_module_name.ncd \  
    top_level_module_name.pcf
```

## Generating a Bitstream

To create the configuration bitstream (BIT) file based on the contents of a physical implementation file (NCD), the **bitgen** command must be executed. The BIT file defines the behavior of the programmed FPGA. An example of the **bitgen** command is:

```
$ bitgen -w top_level_module_name.ncd
```

## Post-Implementation Simulation

The purpose of post-implementation simulation is to verify that the design as implemented in the FPGA works as expected.

### Generating a Simulation Model

Run the **netgen** command to generate a chip-level simulation netlist for your design.

**VHDL**

```
$ netgen -sim -ofmt vhd1 -ngm top_level_module_name_map.ngm \  
    -tm netlist top_level_module_name.ncd \  
    top_level_module_name_postimp.vhd
```

**Verilog**

```
$ netgen -sim -ofmt verilog -ngm top_level_module_name_map.ngm \  
    -tm netlist top_level_module_name.ncd \  
    top_level_module_name_postimp.v
```

## Using the Model

For information about setting up your simulator to use the pre-implemented model, see the Xilinx *Synthesis and Verification Design Guide* included in your Xilinx software installation.

## Other Implementation Information

For more information about using the Xilinx implementation tool flow, including command line switches and options, see the Xilinx ISE software manuals.

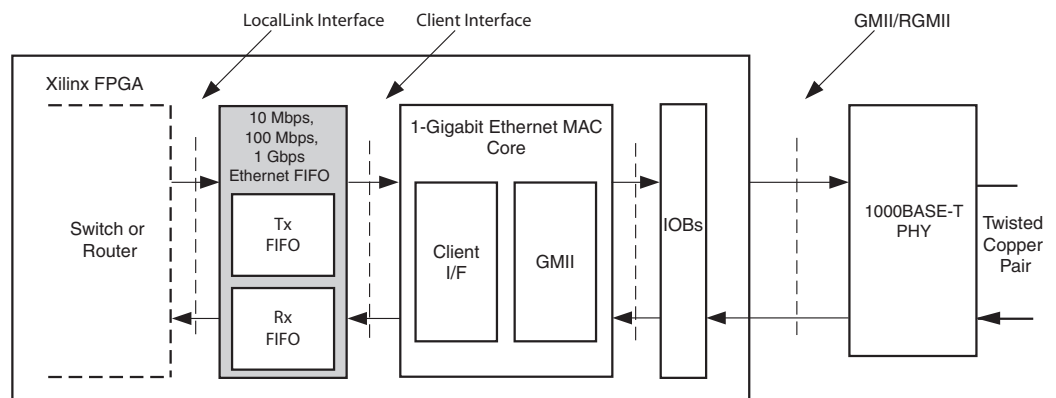


## Using the Client-Side FIFO

The example design provided with the GEMAC core contains a FIFO used on the client-side of the core. The source code for the FIFO is provided, and may be used and adjusted for user applications.

The 10 Mbps/100 Mbps/1 Gbps Ethernet FIFO is designed for use with the GEMAC and Tri-Mode Ethernet MAC (TEMAC) cores. The FIFO directly interfaces to the MAC client interface providing a buffer between the MAC and the user's logic. The FIFO implements a LocalLink user interface, (see "[Overview of LocalLink Interface](#)," on page 123) allowing a direct connection to other LocalLink modules or the user's logic.

The 10 Mbps/100 Mbps/1 Gbps Ethernet FIFO consists of independent transmit and receive FIFOs embedded in a top-level wrapper. [Figure A-1](#) shows how the FIFO fits into a typical implementation.



**Figure A-1: Typical 10 Mbps/100 Mbps/ 1 Gbps Ethernet FIFO Implementation**

## Interfaces

### Transmit FIFO

Table A-1 describes the transmit FIFO client interface. For more information on the MAC client interface, see [“Transmitting Outbound Frames,”](#) on page 43.

**Table A-1: Transmit FIFO Client Interface**

Signal	Direction	Clock Domain	Description
tx_clk	Input	N/A	Transmit clock used by MAC.
tx_reset	Input	tx_clk	Synchronous reset.
tx_enable	Input	tx_clk	Clock enable for tx_clk. Tie to logic 1 when using GEMAC.
tx_data[7:0]	Output	tx_clk	Data presented to MAC for transmission.
tx_data_valid	Output	tx_clk	Valid signal for data.
tx_ack	Input	tx_clk	Ack signal from MAC.
tx_underrun	Output	tx_clk	Underrun signal to MAC.
tx_collision	Input	tx_clk	Collision indication from MAC. Tie to logic 0 when using GEMAC.
tx_retransmit	Input	tx_clk	Retransmit request from MAC. Tie to logic 0 when using GEMAC.

Table A-2 describes the transmit FIFO LocalLink interface. For more information on the LocalLink interface see [“Overview of LocalLink Interface,”](#) on page 123.

**Table A-2: Transmit FIFO LocalLink Interface**

Signal	Direction	Clock Domain	Description
tx_ll_clock	Input	N/A	Write clock for LocalLink interface
tx_ll_reset	Input	tx_ll_clock	Synchronous reset
tx_ll_data_in[7:0]	Input	tx_ll_clock	Write data to be sent to transmitter
tx_ll_sof_in_n	Input	tx_ll_clock	Start of frame indicator
tx_ll_eof_in_n	Input	tx_ll_clock	End of frame indicator
tx_ll_src_rdy_in_n	Input	tx_ll_clock	Source ready indicator
tx_ll_dst_rdy_out_n	Output	tx_ll_clock	Destination ready indicator
tx_fifo_status[3:0]	Output	tx_ll_clock	FIFO memory status
tx_overflow	Output	tx_ll_clock	Overflow signal indicates when a frame has been dropped in the FIFO

## Receive FIFO

Table A-3 describes the receive FIFO client interface. For more information on the MAC client interface, see [“Receiving Inbound Frames,”](#) on page 37.

Table A-3: Receive FIFO Client Interface

Signal	Direction	Clock Domain	Description
rx_clk	Input	N/A	Receive clock used by MAC
rx_reset	Input	rx_clk	Synchronous reset
rx_enable	Input	rx_clk	Clock enable for rx_clk, tie to logic 1 when using GEMAC
rx_data[7:0]	Input	rx_clk	Data received from MAC
rx_data_valid	Input	rx_clk	Valid signal for data
rx_good_frame	Input	rx_clk	Indicates if frame is valid and should be accepted by client
rx_bad_frame	Input	rx_clk	Indicates if frame is invalid and should be dropped by the FIFO
rx_overflow	Output	rx_clk	Overflow signal indicates when a frame has been dropped in the FIFO

Table A-4 describes the receive FIFO LocalLink interface. For more information on the LocalLink interface, see [“Overview of LocalLink Interface,”](#) on page 123.

Table A-4: Receive FIFO LocalLink Interface

Signal	Direction	Clock Domain	Description
rx_ll_clock	Input	N/A	Read clock for LocalLink interface
rx_ll_reset	Input	rx_ll_clock	Synchronous reset
rx_ll_data_out[7:0]	Output	rx_ll_clock	Data read from FIFO
rx_ll_sof_out_n	Output	rx_ll_clock	Start of frame indicator
rx_ll_eof_out_n	Output	rx_ll_clock	End of frame indicator
rx_ll_src_rdy_out_n	Output	rx_ll_clock	Source ready indicator
rx_ll_dst_rdy_in_n	Input	rx_ll_clock	Destination ready indicator
rx_fifo_status[3:0]	Output	rx_ll_clock	FIFO memory status

## Overview of LocalLink Interface

### Data Flow

Data is transferred on the LocalLink interface from source to destination, with the flow being governed by the four active low control signals `sof_n`, `eof_n`, `src_rdy_n`, and

`dst_rdy_n`. The flow of data is controlled by the `src_rdy_n` and `dst_rdy_n` signals. Only when these signals are asserted simultaneously is data transferred from source to destination. The individual packet boundaries are marked by the `sof_n` and `eof_n` signals. For more information on the LocalLink interface see to Xilinx Application Note [XAPP691](http://direct.xilinx.com/bvdocs/appnotes/xapp691.pdf), “Parameterizable LocalLink FIFO” available at [direct.xilinx.com/bvdocs/appnotes/xapp691.pdf](http://direct.xilinx.com/bvdocs/appnotes/xapp691.pdf).

Figure A-2 shows the transfer of an 8-byte frame.

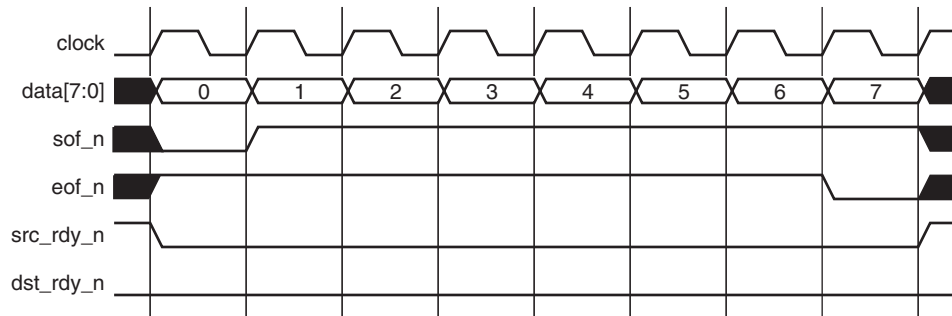


Figure A-2: Frame Transfer across LocalLink Interface

Figure A-3 illustrates frame transfer of a 5-byte frame, where both the `src_rdy_n` and `dst_rdy_n` signals are used to control the flow of data across the interface.

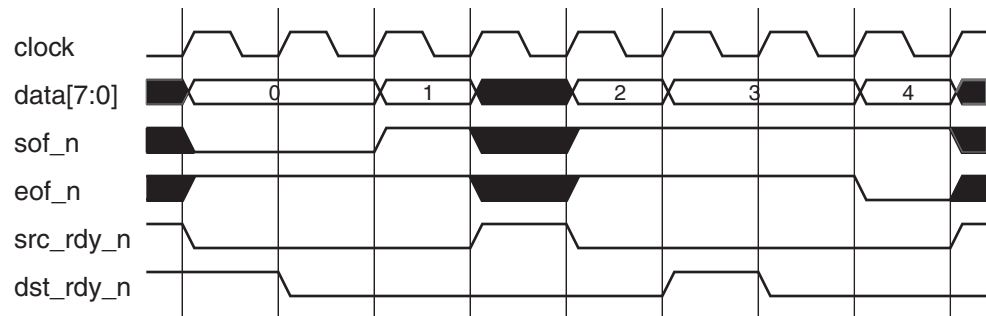


Figure A-3: Frame Transfer with Flow Control

## Functional Operation

### Clock Requirements

The FIFO is designed to work with `rx_clk` and `tx_clk` running at MAC clock speeds up to 125 MHz. The `rx_ll_clock` should be no slower than the `rx_clk`. The `tx_ll_clock` should be no slower than the clock on the transmitter client interface divided by 2. For this reason, it is suggested that the `rx_ll_clock` and `tx_ll_clock` are always 125 MHz or faster.

### Receive FIFO

The receive FIFO is built around two Dual Port Block RAMs giving a memory capacity of 4096 bytes.

The receive FIFO takes data from the client interface of the GEMAC core and converts it into LocalLink format. See “Receiving Inbound Frames,” on page 37 for a description of the GEMAC receive client interface. If the frame is marked as good by the GEMAC, that frame is presented on the LocalLink interface for reading by the user. If the frame is marked as bad, it is dropped by the FIFO.

If the receive FIFO memory overflows, the frame currently being received is dropped, regardless of whether it is a good or bad frame. The signal `rx_overflow` is then asserted. Situations in which the memory may overflow are:

- The FIFO may overflow if the user is reading data from the FIFO at a slower data rate than data is being written in from the MAC receiver.
- The FIFO size of 4096 bytes limits the size of the frames that it can store without error. If a frame is larger than 4000 bytes, the FIFO may overflow and data will be lost. For this reason, it is recommended that the FIFO not be used with the GEMAC in jumbo frame mode for frames larger than 4000 bytes.

## Transmit FIFO

The transmit FIFO is built around two Dual Port block RAMs, giving a total memory capacity of 4096 bytes of frame data.

The transmit FIFO accepts frames in LocalLink format and stores them in block RAM for transmission via the GEMAC. When a full frame has been written into the transmit FIFO, the FIFO will present data to the MAC transmitter. On receiving the `tx_ack` signal from the MAC, the rest of the frame is transmitted to the MAC.

### VHDL

The generic `FULL_DUPLEX_ONLY` is provided to allow the removal of logic and performance constraints necessary for half-duplex operation when using with the Xilinx Tri-Mode Ethernet MAC core. This generic can always be set to true when the FIFO is used with the GEMAC.

### Verilog

The compiler directive `FULL_DUPLEX_ONLY` is defined to allow for removal of logic and performance constraints that are necessary only in half-duplex operation, that is, when using with the Tri-Mode Ethernet MAC core. This directive can always be defined when the FIFO is used with the GEMAC.

The FIFO has two signal inputs specific to half-duplex operation, `tx_collision` and `tx_retransmit`. These signals are provided to make the FIFO compatible with both the 1-Gigabit Ethernet MAC and Tri-Mode Ethernet MAC cores, and should be tied to logic 0 when using the FIFO with the GEMAC core.

If the FIFO memory fills up, the `dst_rdy_out_n` signal is used to halt the LocalLink interface writing in data until space becomes available in the FIFO. If the FIFO memory fills up but no frames are available for transmission (for example, if a frame larger than 4000 bytes is written into the FIFO), the FIFO may assert the `tx_overflow` signal and continue to accept the rest of the frame from the user. The overflow frame will be dropped by the FIFO. This ensures that the LocalLink interface does not lock up. For this reason, it is recommended that the FIFO not be used with the GEMAC in jumbo frame mode for frames larger than 4000 bytes.

## Expanding Maximum Frame Size

The transmit FIFO size is optimized to allow line rate transmission of maximum size Ethernet frames at 1518 bytes in half-duplex operation.

When using the FIFO in full-duplex operation, the full block RAM capacity can be utilized at all times. As a whole frame must be stored in the FIFO block RAM before being presented to the MAC transmitter, the maximum size frame that can be handled is determined by the memory capacity of the FIFO (in this case 4000 bytes).

Both transmit and receive FIFO sizes can be expanded by the user to handle larger frame sizes. This can be done by instantiating further block RAMs into the FIFO design, expanding the block RAM address signals, and adding the necessary control signals. The HDL source files provide guidance in the comments on how to achieve this.

## User Interface Data Width Conversion

Conversion of the user interface 8 bit data path to a 16, 32, 64 or 128 bit data path can be made by connecting the LocalLink interface directly to the Parameterizable LocalLink FIFO, Xilinx Application Note [XAPP691, Parameterizable LocalLink FIFO](http://direct.xilinx.com/bvdocs/appnotes/xapp691.pdf) found at [direct.xilinx.com/bvdocs/appnotes/xapp691.pdf](http://direct.xilinx.com/bvdocs/appnotes/xapp691.pdf).

# *Core Verification, Compliance, and Interoperability*

---

The GEMAC core has been verified with extensive simulation and hardware testing.

## **Verification by Simulation**

A highly parameterizable transaction-based test bench (not part of the core deliverables) was used to test the core. Tests include:

- Register Access
- MDIO Access
- Frame Transmission and error handling
- Frame Reception and error handling
- Address Filtering

## **Hardware Verification**

The GEMAC core has been tested in a variety of hardware test platforms at Xilinx to include a variety of parameterizations, including the following.

The core has been tested with the Ethernet 1000BASE-X PCS/PMA or SGMII core from Xilinx. This follows the architecture shown in [Figure 11-2, page 109](#). A test platform was built around these cores, including a back-end FIFO capable of performing a simple ping function, and a test pattern generator. Software running on the embedded PowerPC™ was used to provide access to all configuration, status and statistical counter registers. Version 3.0 of this core was taken to the University of New Hampshire Interoperability Lab (UNH IOL) where conformance and interoperability testing was performed.

The core has been tested with an external 1000BASE-T PHY device. The MAC was connected to the external PHY device using GMII, RGMII, and SGMII (in conjunction with the Ethernet 1000BASE-X PCS/PMA or SGMII core).

The core has been tested with the Ethernet Statistics core from Xilinx, following the architecture show in [Figure 11-5, page 114](#).





# Calculating DCM Phase-Shifting

---

## DCM Phase-Shifting

A DCM is used in the receiver clock path to meet the input setup and hold requirements when using the core with an RGMII (see [“Implementing External RGMII,” on page 60](#)) and with an external GMII implementation in Spartan-3, Spartan-3E, and Spartan-3A devices (see [“Spartan-3, Spartan-3E and Spartan-3A Devices,” on page 58.](#))

In these cases, a fixed phase-shift offset is applied to the receiver clock DCM to skew the clock. This performs static alignment by using the receiver clock DCM to shift the internal version of the receiver clock such that its edges are centered on the data eye at the IOB DDR flip-flops. The ability to shift the internal clock in small increments is critical for sampling high-speed source synchronous signals such as RGMII. For statically aligned systems, the DCM output clock phase offset (as set by the phase shift value) is a critical part of the system, as is the requirement that the PCB is designed with precise delay and impedance-matching for all the GMII/RGMII receiver data bus and control signals.

You must determine the best DCM setting (phase-shift) to ensure that the target system has the maximum system margin to perform across voltage, temperature, and process (multiple chips) variations. Testing the system to determine the best DCM phase-shift setting has the added advantage of providing a benchmark of the system margin based on the UI (unit interval or bit time).

System margin is defined as the following:

$$\text{System Margin (ps)} = \text{UI(ps)} * (\text{working phase-shift range}/128)$$

## Finding the Ideal Phase-Shift

Xilinx cannot recommend a singular phase-shift value that is effective across all hardware platforms, and does not recommend attempting to determine the phase-shift setting empirically. In addition to the clock-to-data phase relationship, other factors such as package flight time (package skew) and clock routing delays (internal to the device) affect the clock-to-data relationship at the sample point (in the IOB) and are difficult to characterize.

Xilinx recommends extensive investigation of the phase-shift setting during hardware integration and debugging. The phase-shift settings provided in the example design UCF are placeholders, and work successfully in back-annotated simulation of the example design.

Perform a complete sweep of phase-shift settings during your initial system test. Use only positive (0 to 255) phase-shift settings, and use a test range that covers a range of no less than 128, corresponding to a total 180 degrees of clock offset. This does not imply that 128 phase-shift values must be tested; increments of 4 (52, 56, 60, etc.) correspond to roughly

one DCM tap, and consequently provide an appropriate step size. Additionally, it is not necessary to characterize areas outside the working phase-shift range.

At the edge of the operating phase-shift range, system behavior changes dramatically. In eight phase-shift settings or less, the system can transition from no errors to exhibiting errors. Checking the operational edge at a step size of two (on more than one board) refines the typical operational phase-shift range. Once the range is determined, choose the average of the high and low working phase-shift values as the default. During the production test, Xilinx recommends that you re-examine the working range at corner case operating conditions to determine whether any final adjustments to the final phase-shift setting are needed.

You can use the FPGA Editor to generate the required test file set instead of resorting to multiple PAR runs. Performing the test on design files that differ only in phase-shift setting prevents other variables from affecting the test results. FPGA Editor operations can even be scripted further, reducing the effort needed to perform this characterization.

## *Core Latency*

---

### **Transmit Path Latency**

As measured from a data octet accepted on `tx_data[7:0]` of the transmitter client-side interface, until that data octet appears on `gmii_txd[7:0]` of the physical side GMII style interface, the latency through the core in the transmit direction is 9 clock periods of `gtx_clk`.

### **Receive Path Latency**

As measured from a data octet accepted on `gmii_rxd[7:0]` of the physical side GMII style interface, until that data octet appears on `rx_data[7:0]` of the receiver client-side interface, the latency through the core in the receive direction is 9 clock periods of `gmii_rx_clk`.

